



Redpanda 21.10.1



- A Jepsen Report -

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Distsys Safety Tester

In a nutshell

3 liveness (crash, delayed message)

7 safety (dup, last write, split-brain, aborted read, G1c)

2 ambiguous (G0, G1c, indefinite error code)

Redpanda

Fundamentals

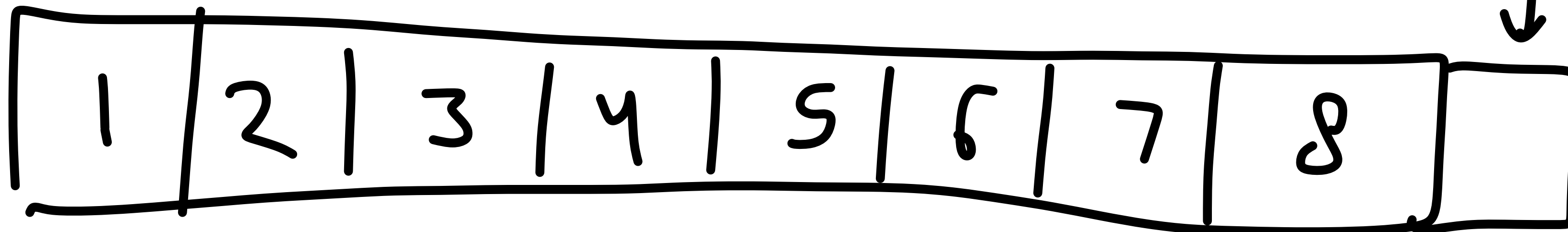
Distributed Log, aka kafka

- API compatible
- Aiming for better performance
- No Zookeeper

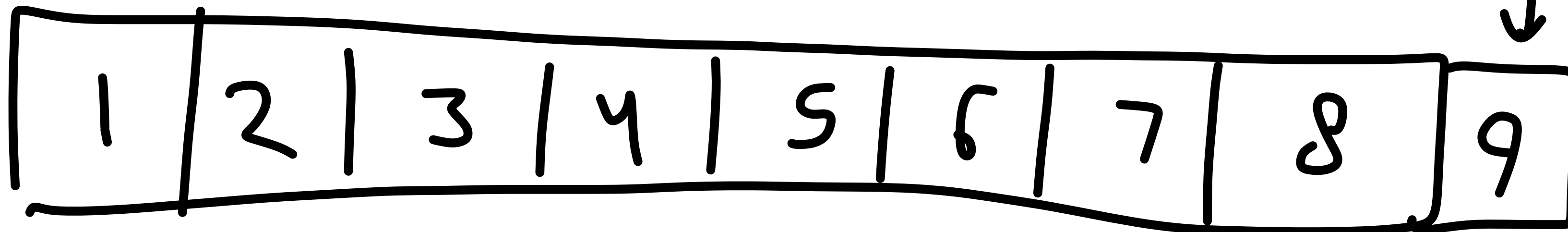
A Partition

1	2	3	4	5	6	7	8	
---	---	---	---	---	---	---	---	--

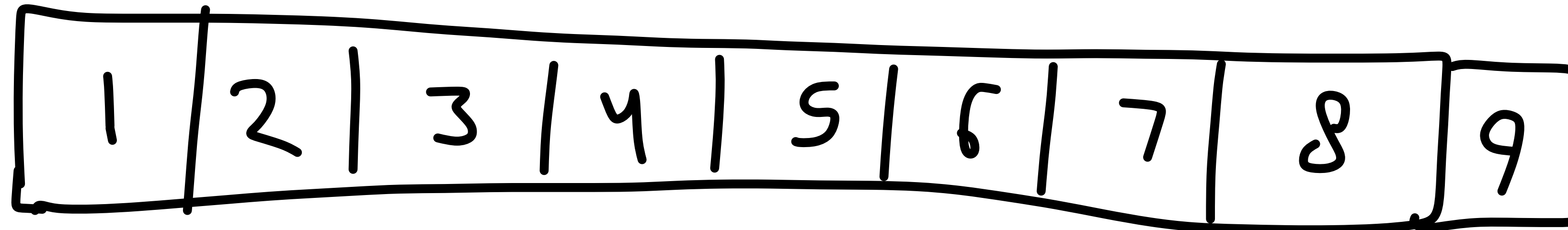
Producer.send(9)



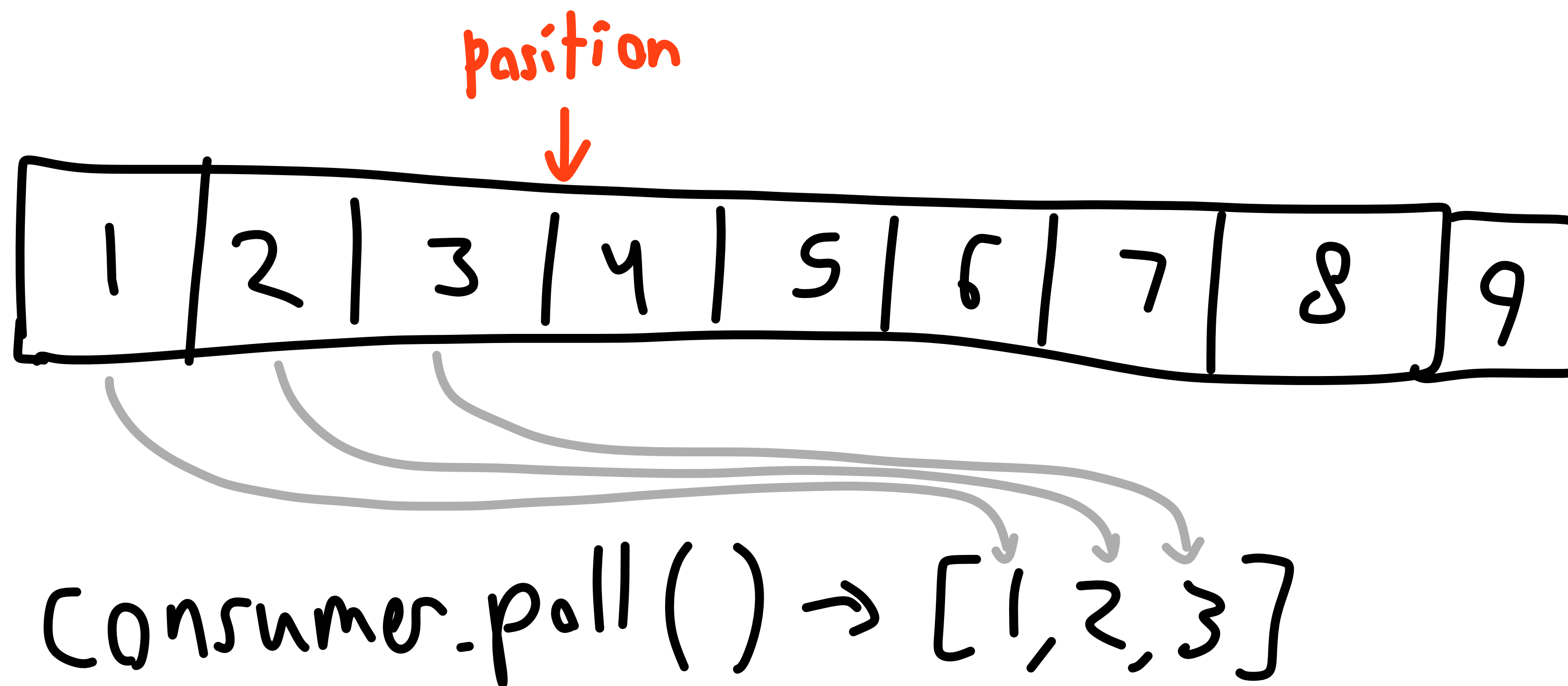
Producer.send(9)

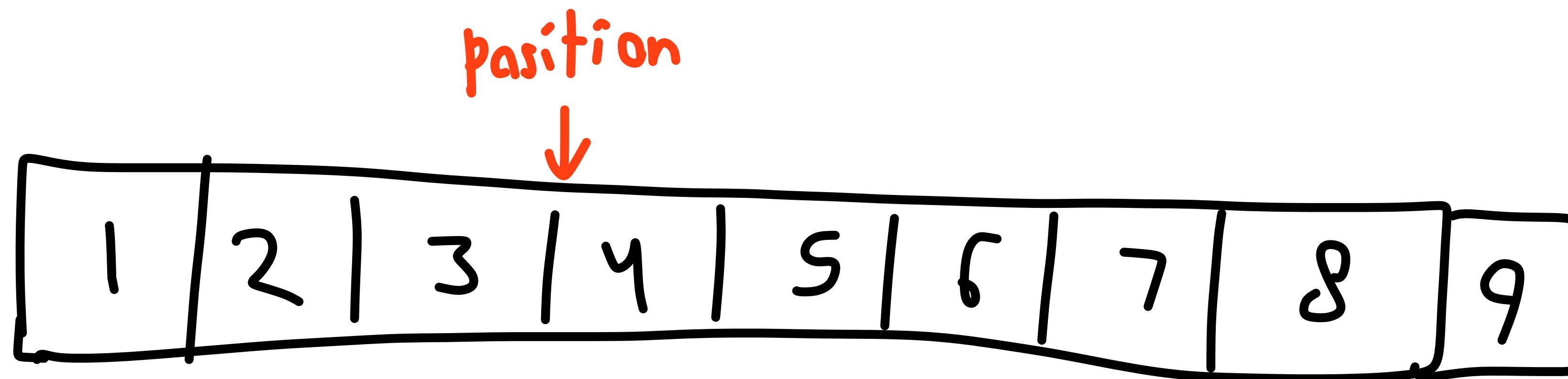


position



Consumer.poll()

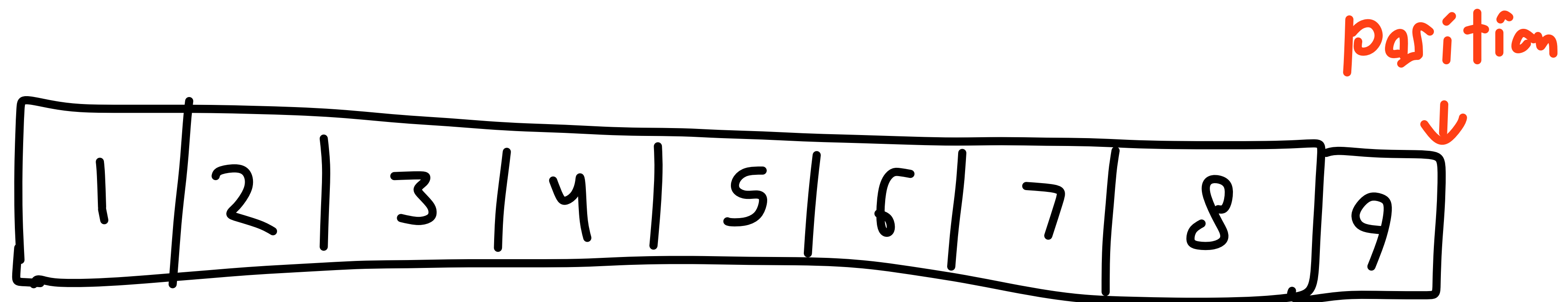




Consumer.poll()

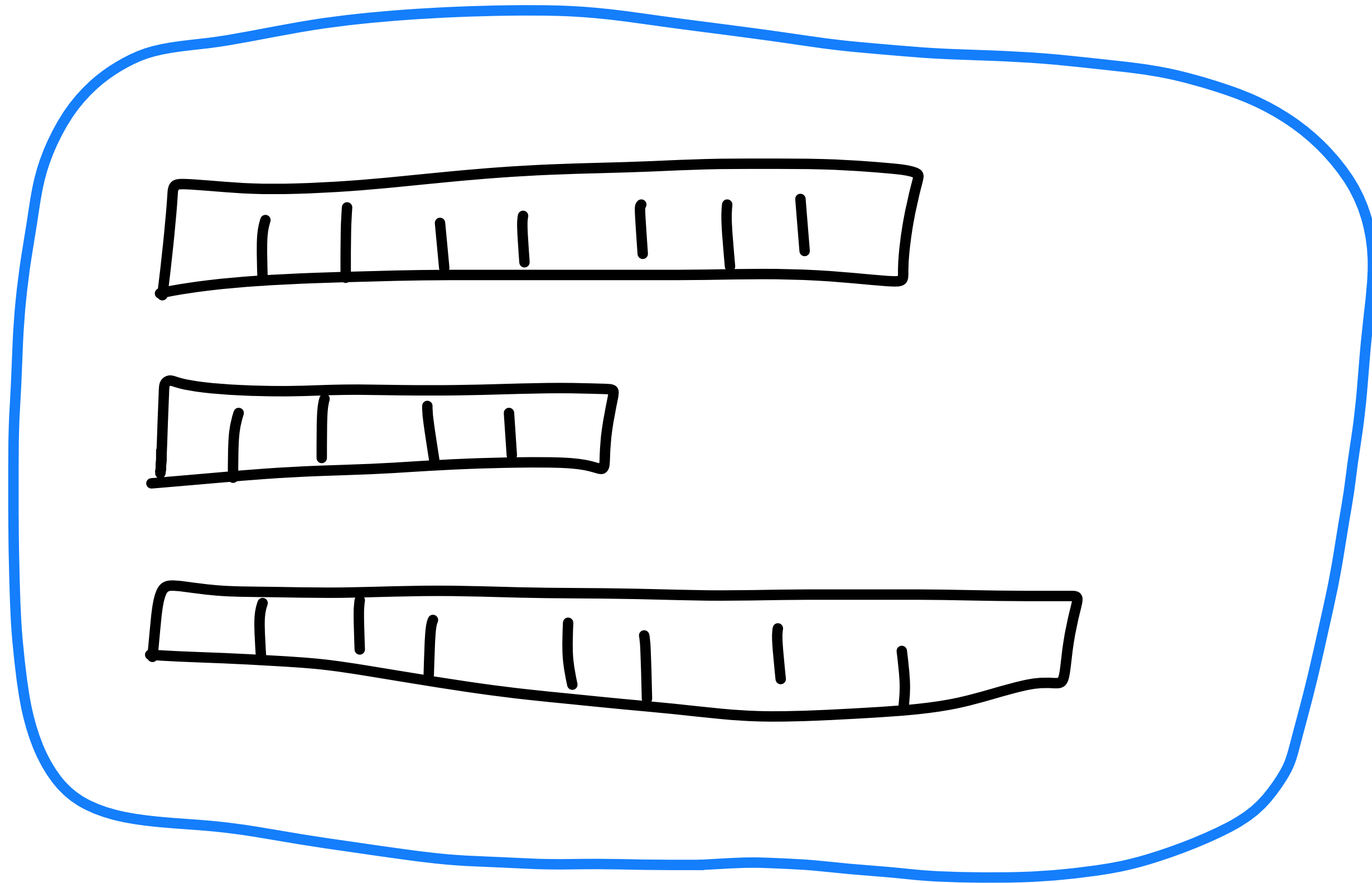


Consumer.poll() → [4, 5, 6, 7, 8, 9]



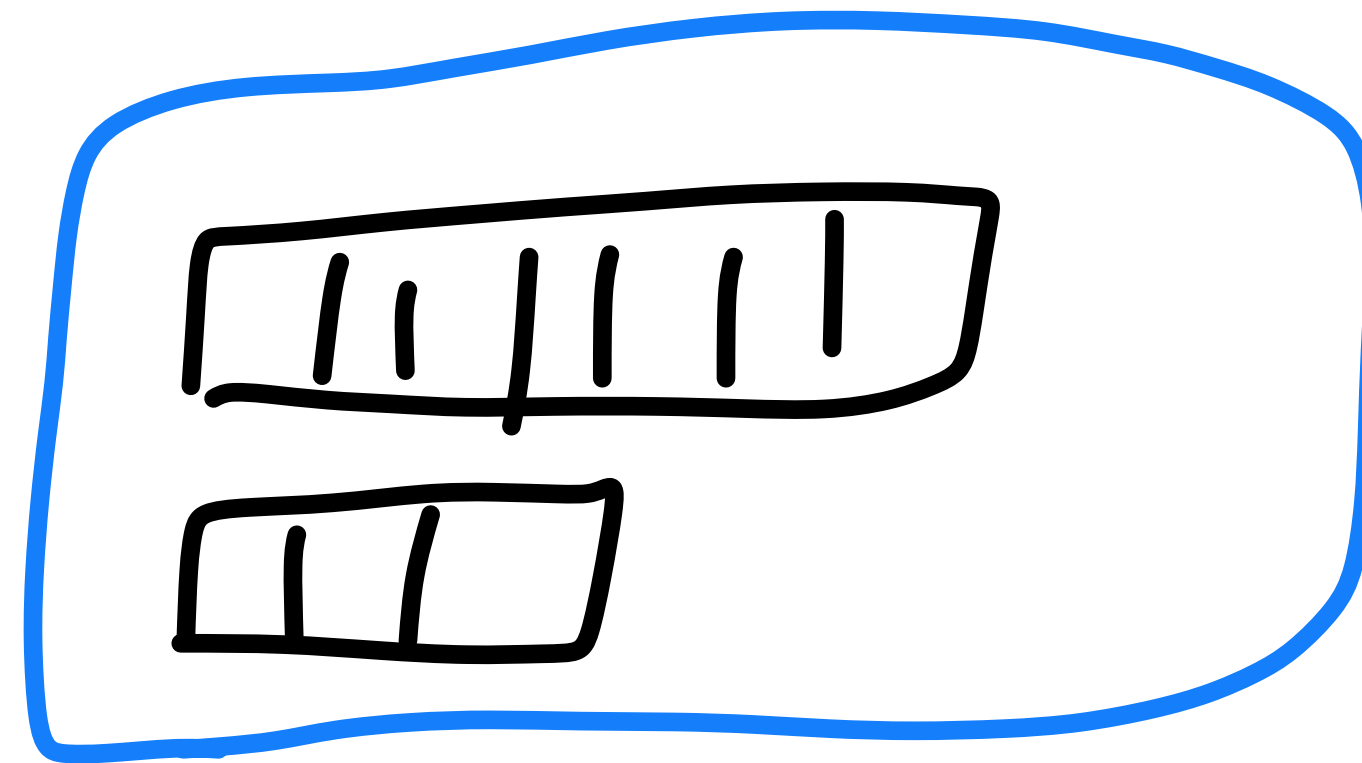
Consumer.poll() → []

A Topic

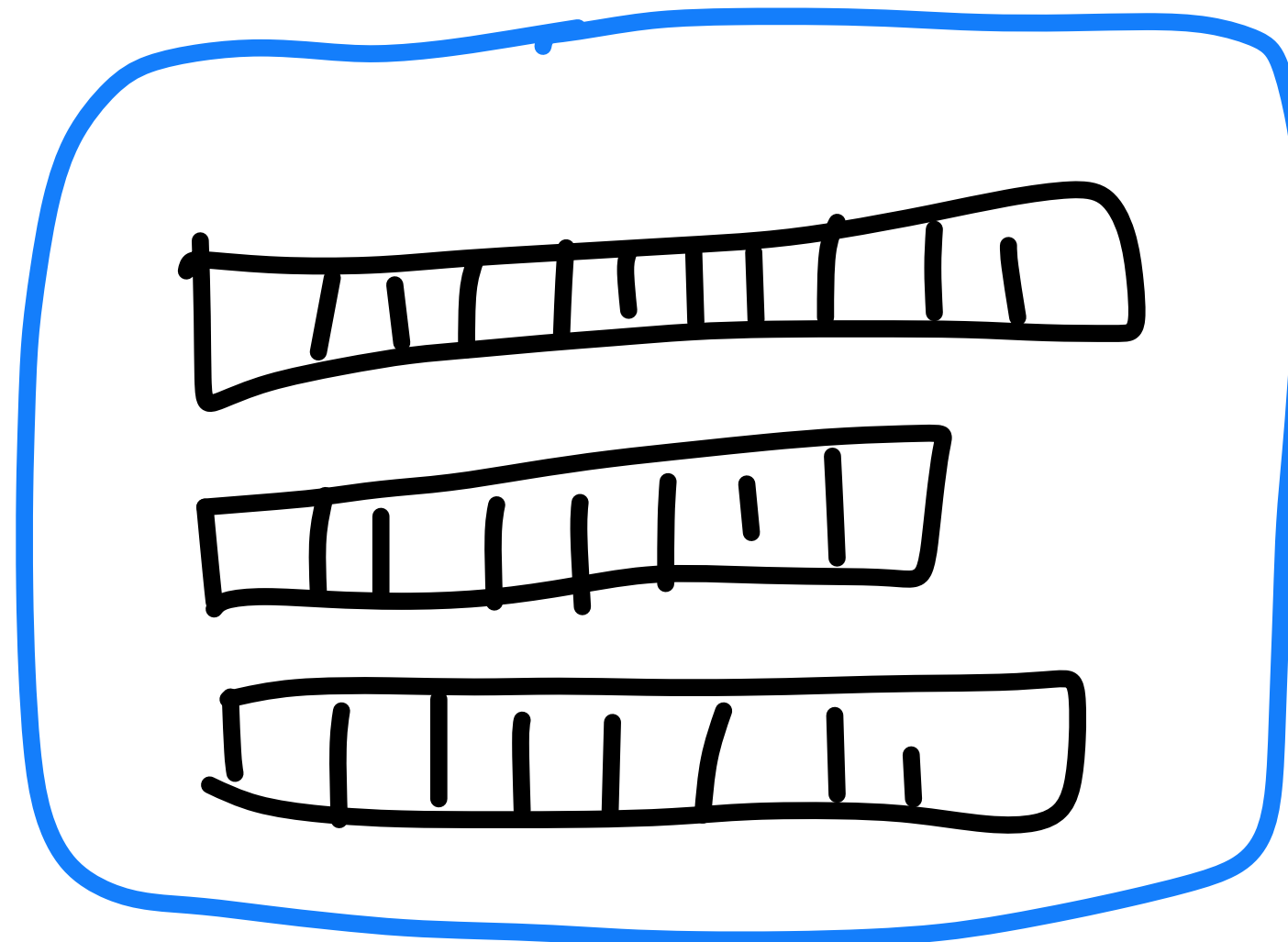


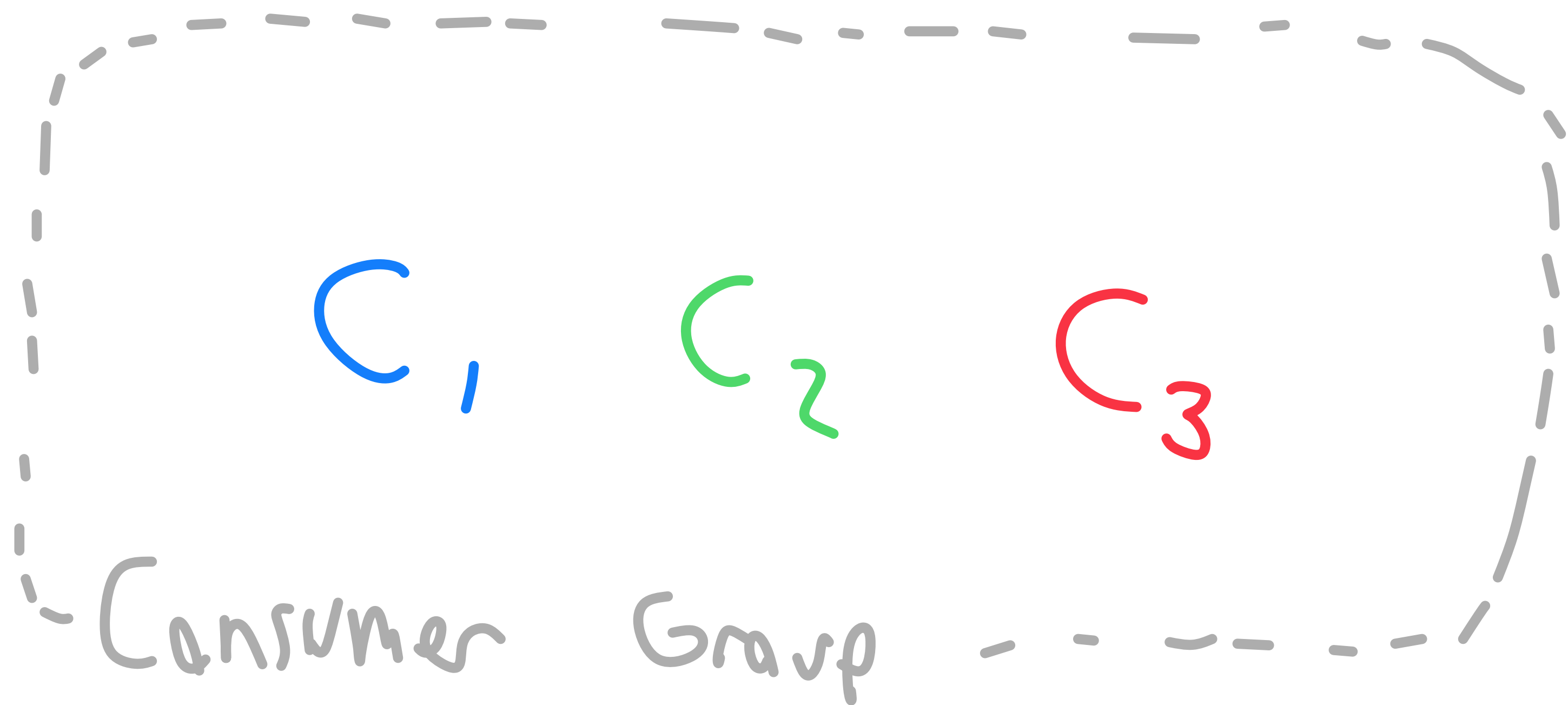
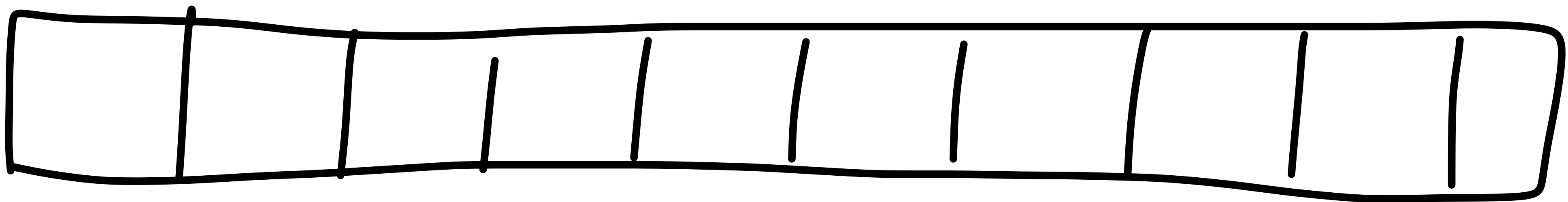
Redpanda

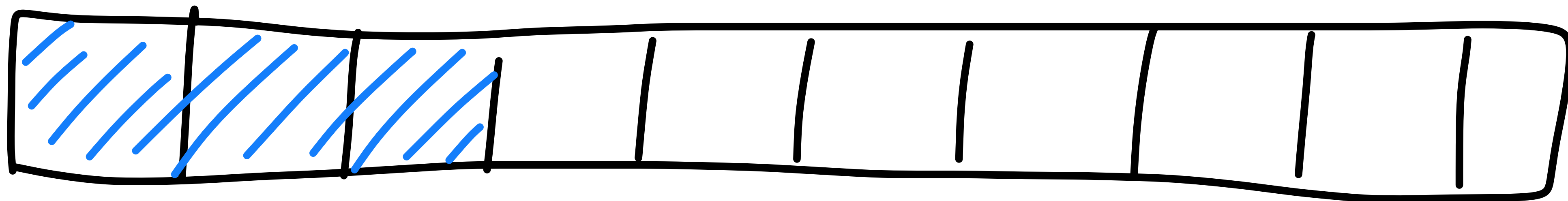
rainfall →



harvests →







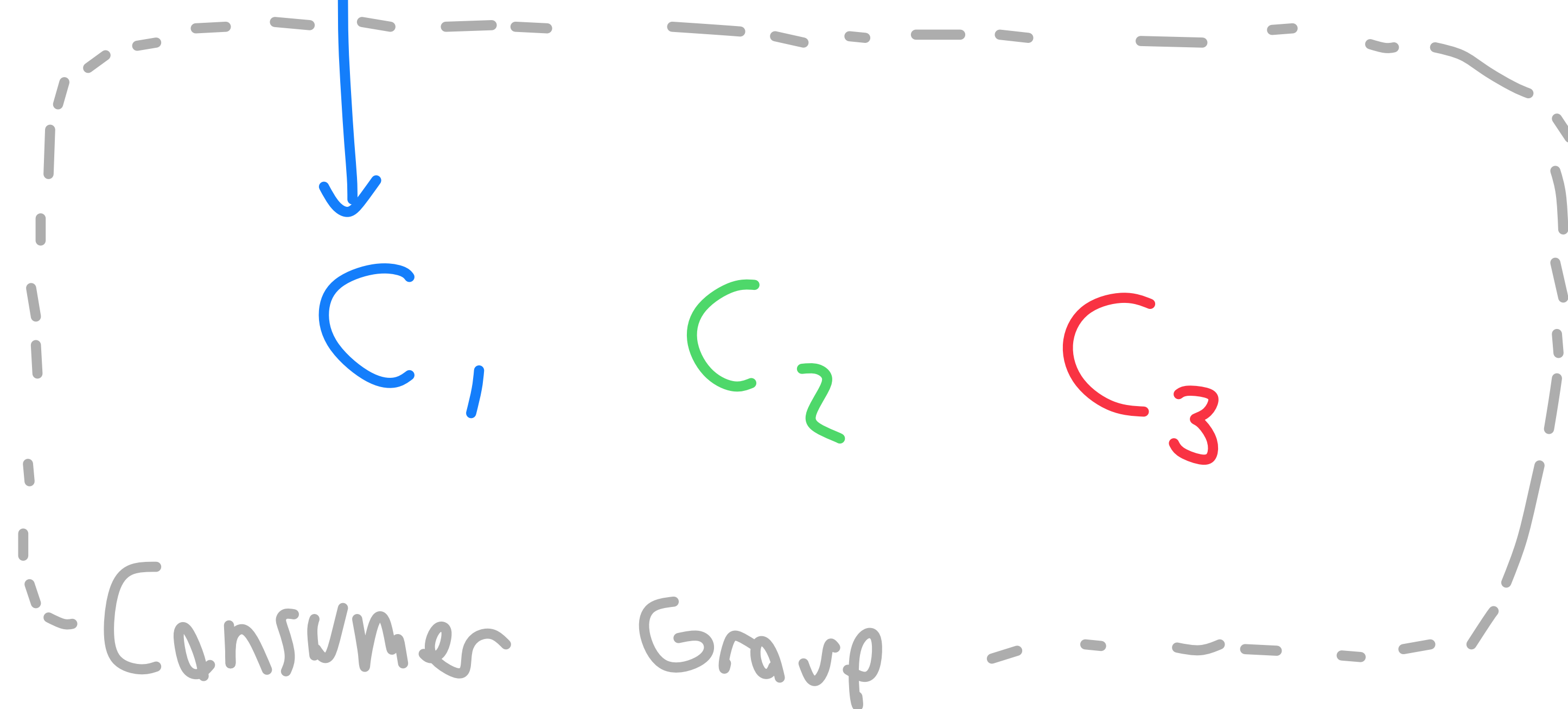
pull

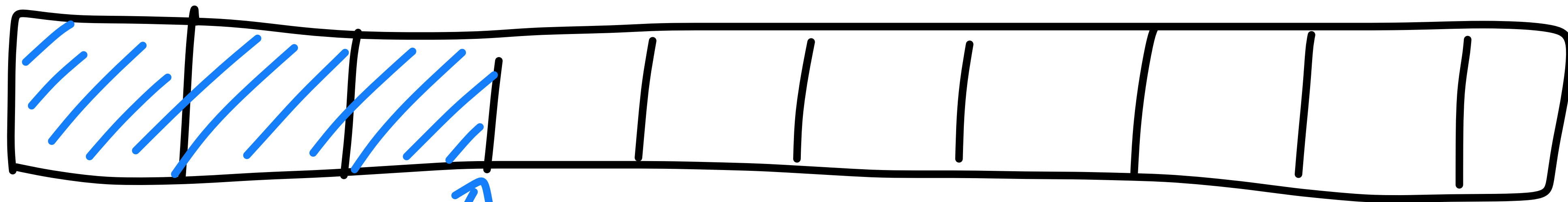


C_1

C_2

C_3





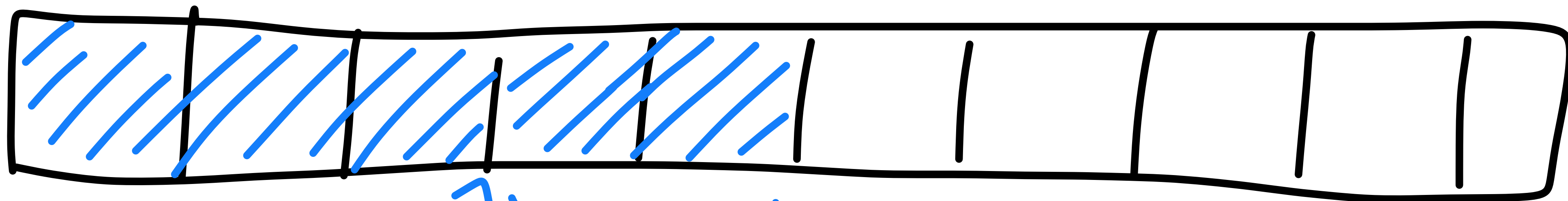
Commit

C_1

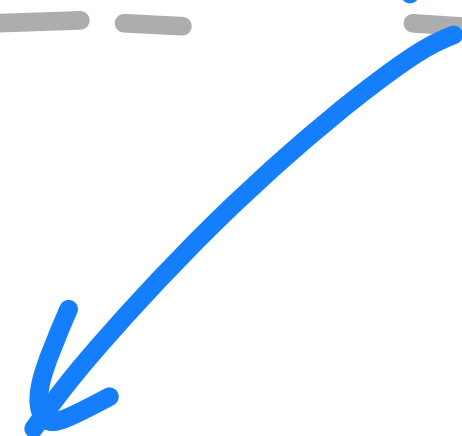
C_2

C_3

Consumer Group



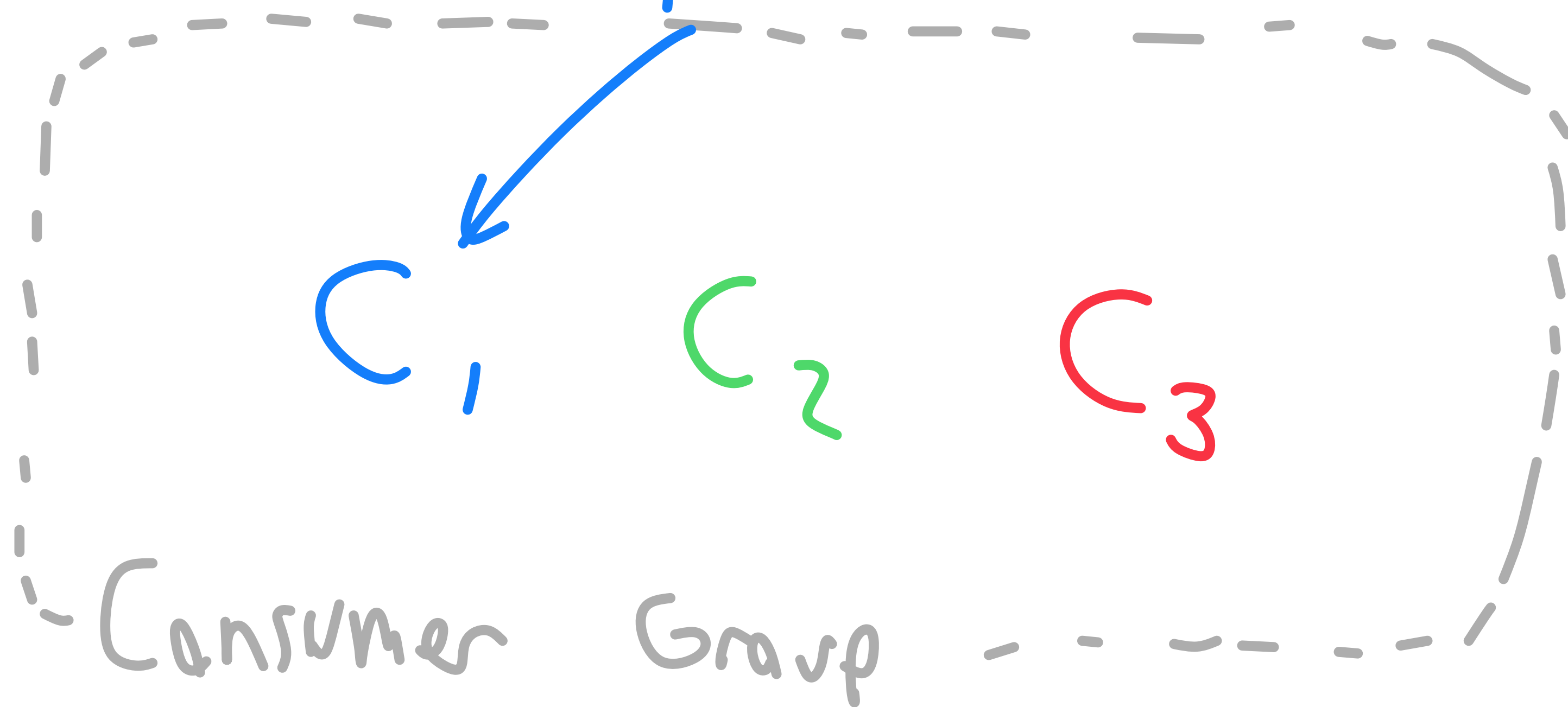
poll

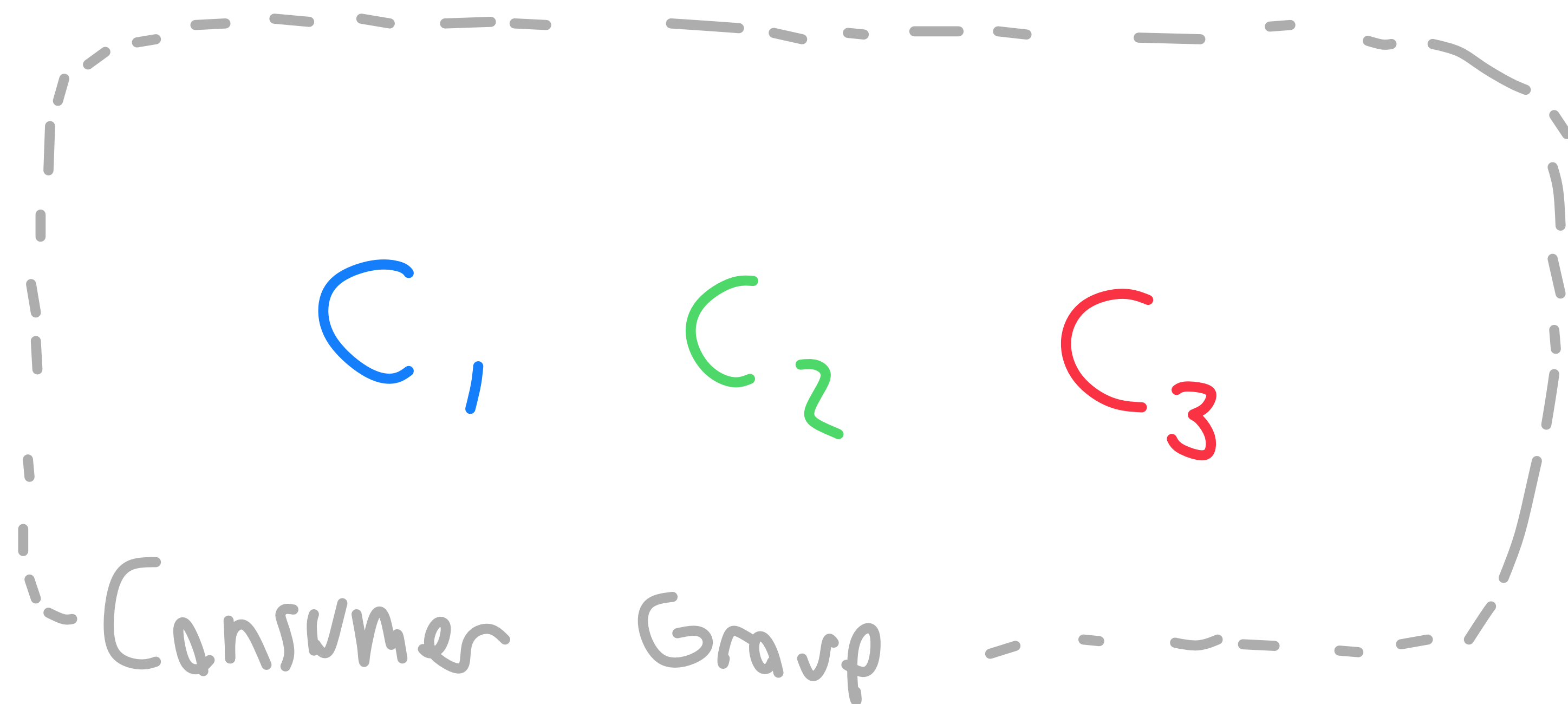
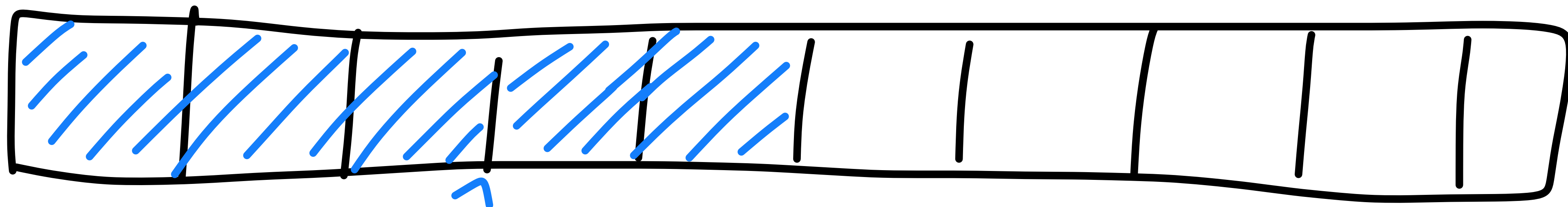


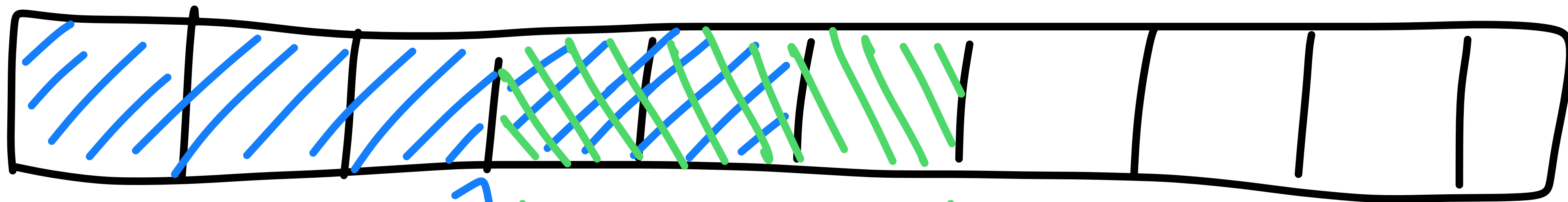
C_1

C_2

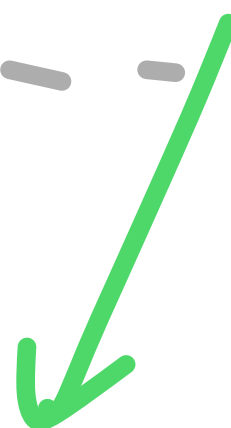
C_3







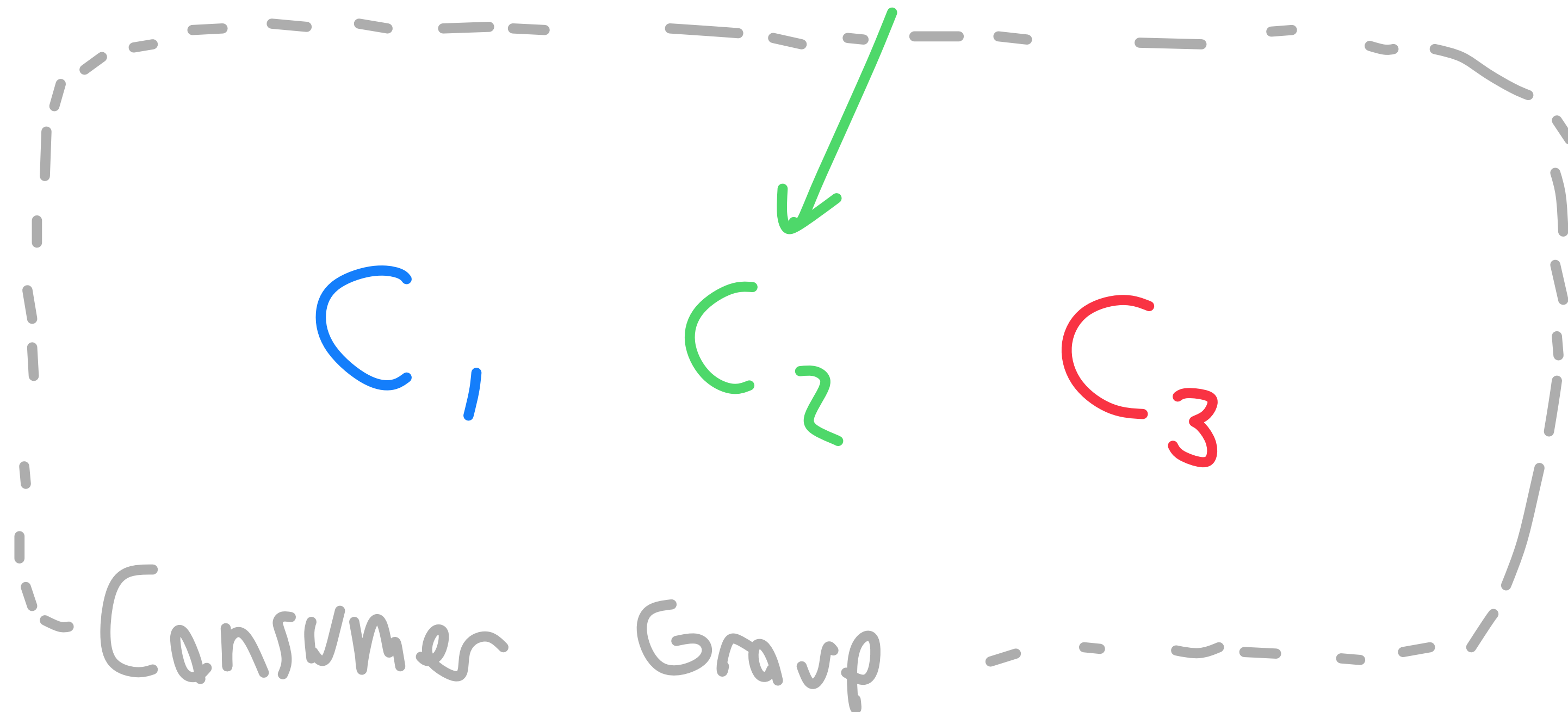
poll



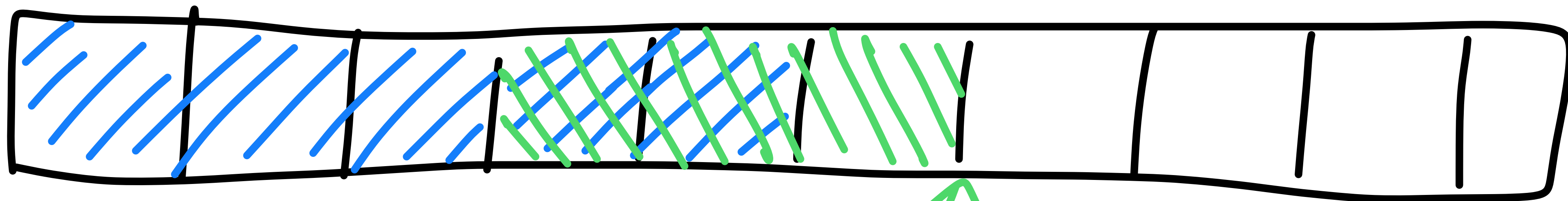
C₁

C₂

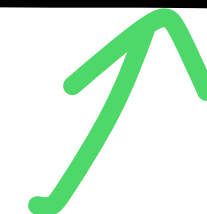
C₃



Consumer Group



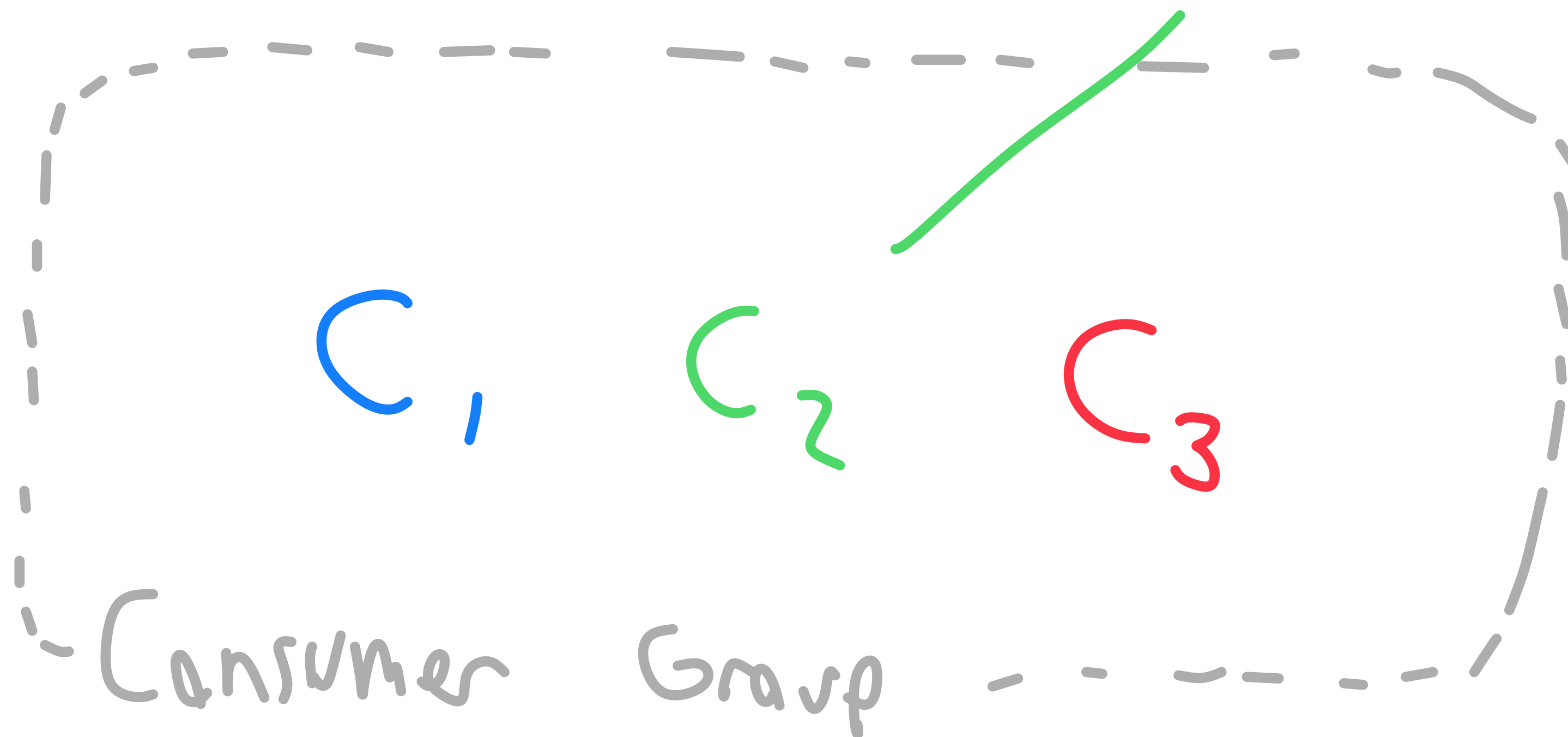
Commit



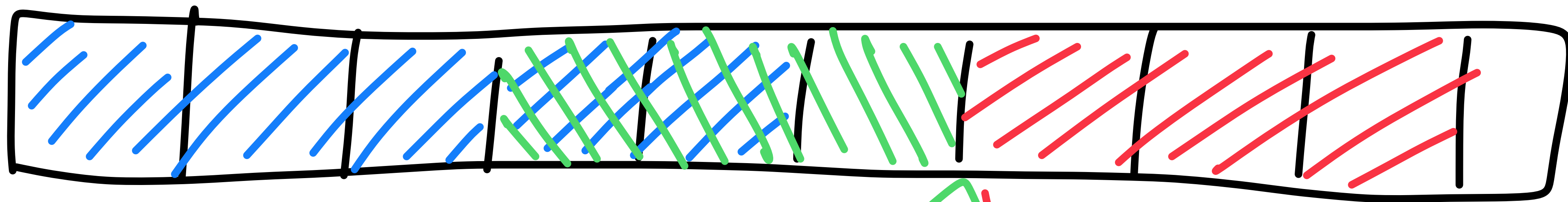
C₁

C₂

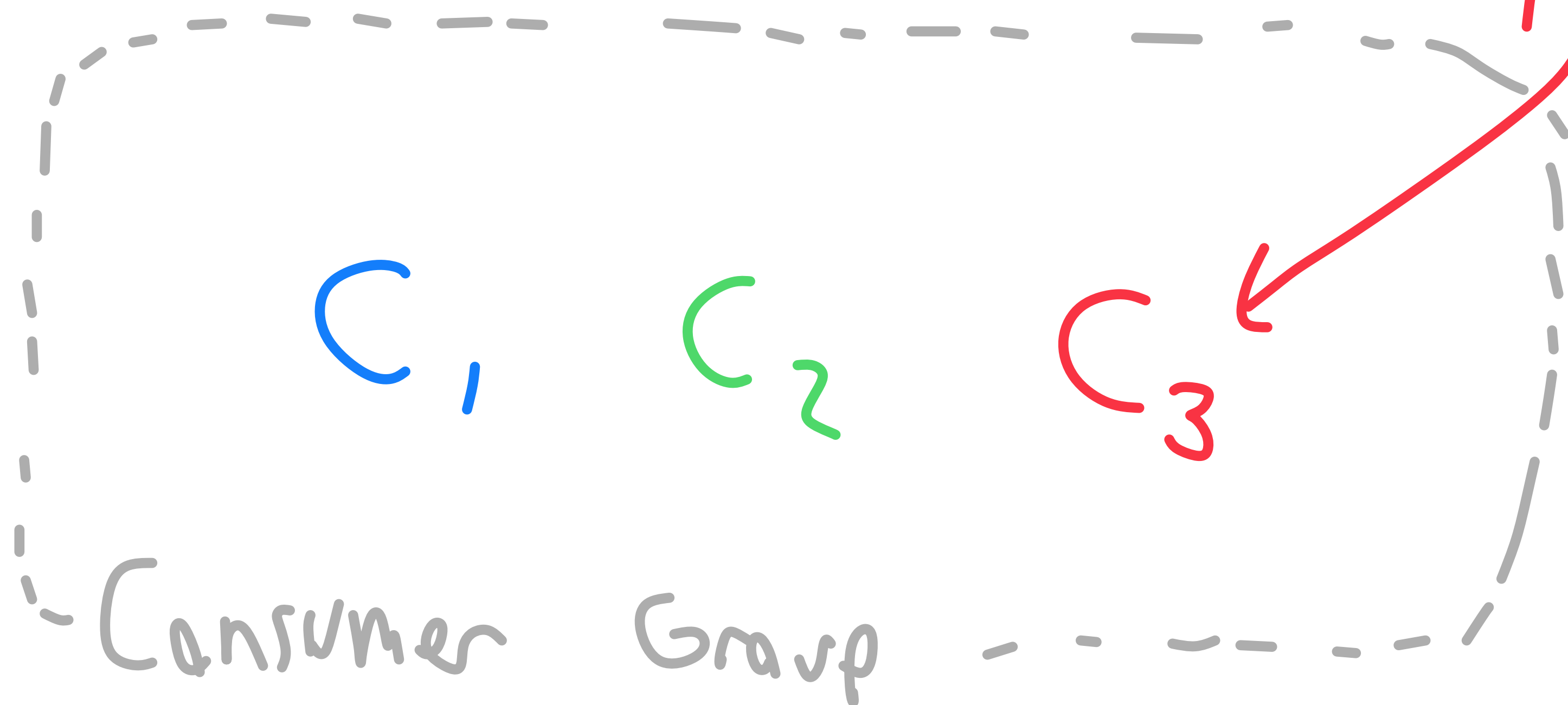
C₃



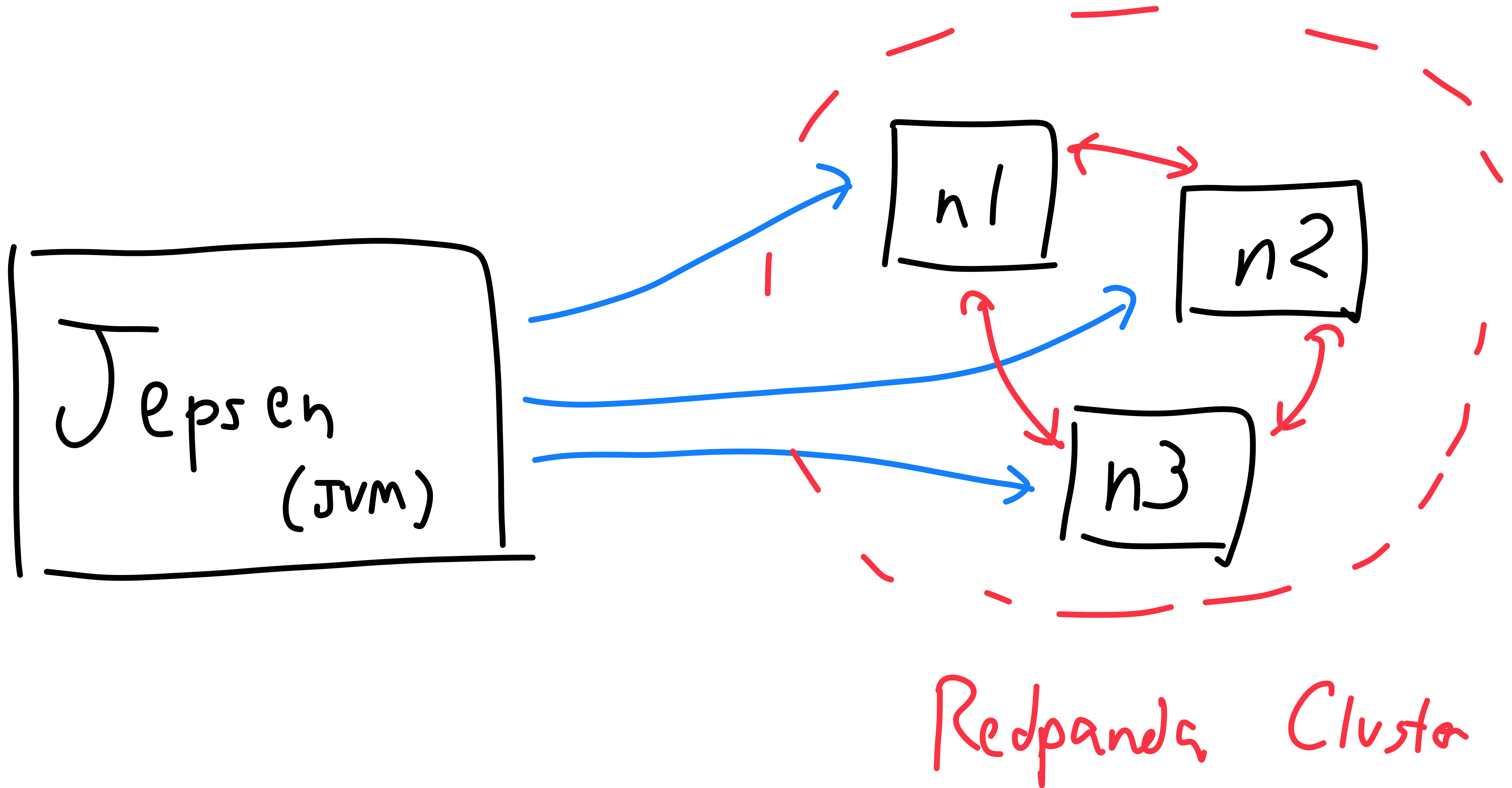
Consumer Group

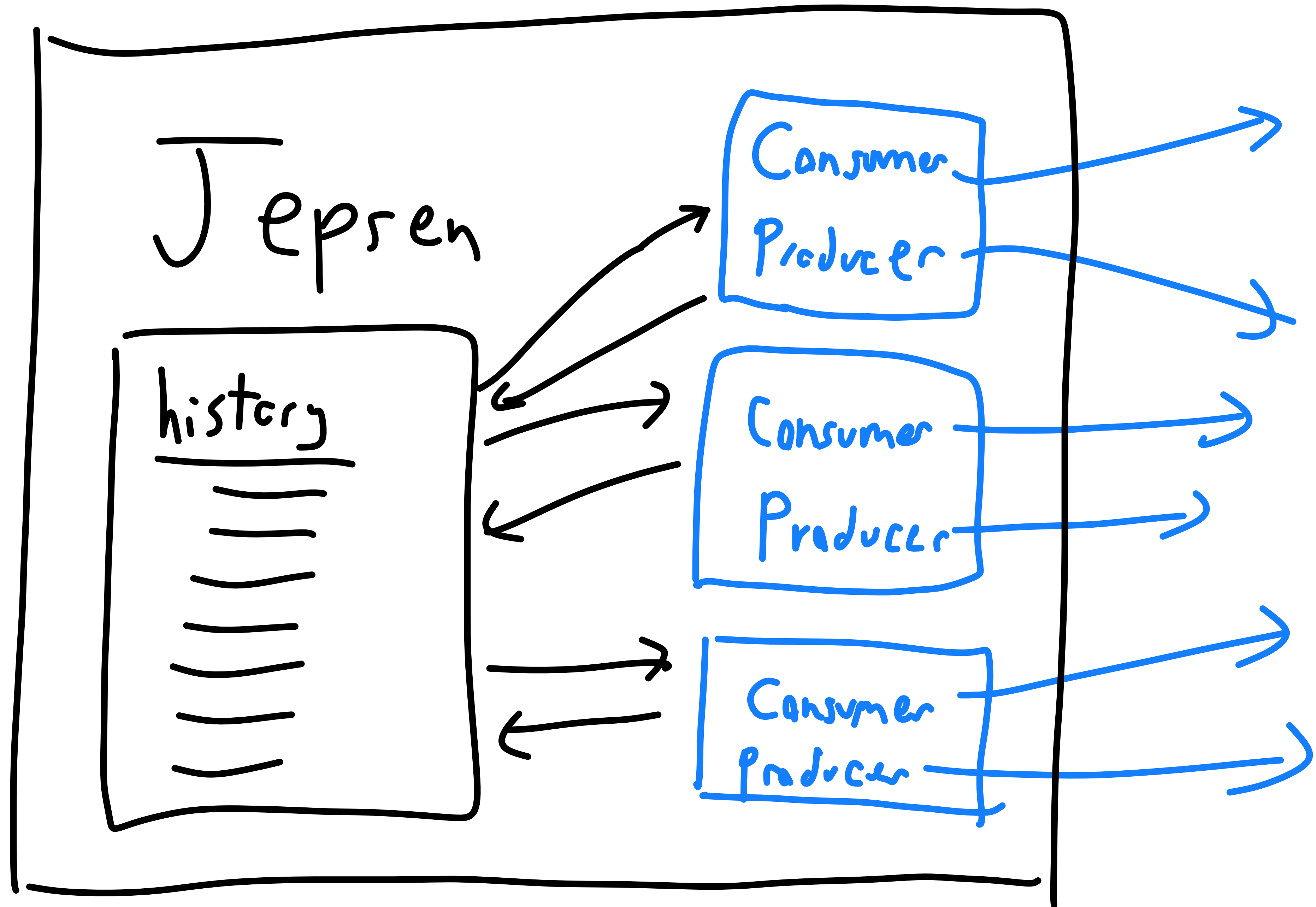


part



Test Design





Queue

History

invoke subscribe [k]

ok subscribe [k]

History

invoke subscribe $[k]$

ok subscribe $[k]$

invoke poll $[poll]$

ok poll $[poll \{k, [[0:a] [1:b]]\}]$

History

invoke subscribe $[k]$

ok subscribe $[k]$

invoke poll $[poll]$

ok poll $[poll \ \{k_1 \ [[0:a] \ [1:b]]\}]$

invoke send $[send \ k_2 : c]$

ok send $[send \ k_2 \ [3 : c]]$

History

ok txn $\left[\left[\text{poll } \{ k_1, [[0:a] [1:b]] \} \right], \right.$
 $\left. \left[\text{send } k_2 [3:c] \right] \right]$

Checker

fail $\left[\text{send } x \text{ } a \right]$

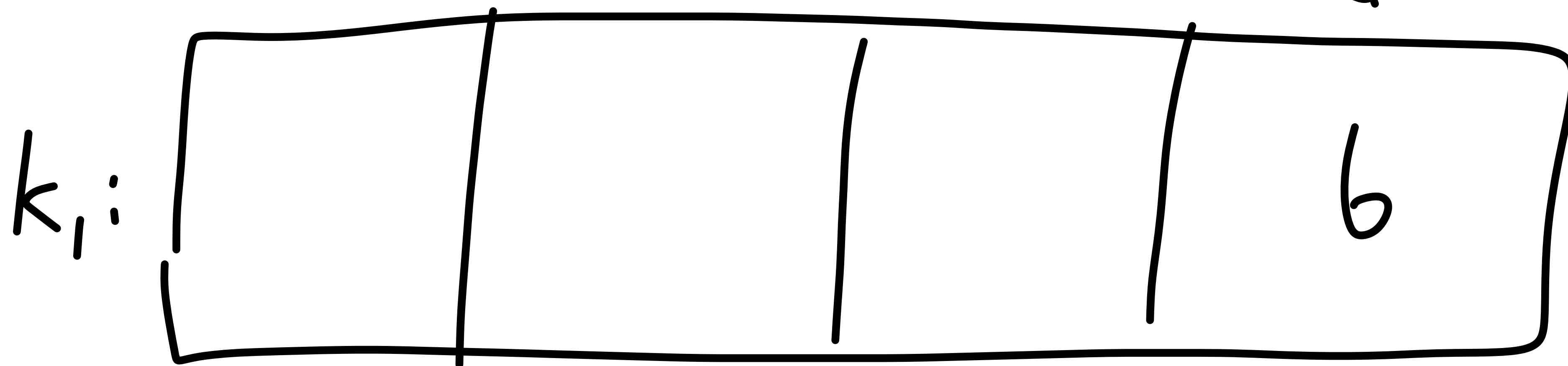
ok $\left[\text{poll } \{ x \left[\exists, a \right] \} \right]$

fail $\left[\text{send } x \textcircled{a} \right]$ ok $\left[\text{poll } \{ x \left[\exists, \textcircled{a} \right] \} \right]$

Aborted Read!

(G1_a)

$[send\ k, [3, 6]]$



"version order"

$[\text{poll } \{ k_1, [[0 \ a] \ [2 \ b] \ [3 \ c]] \}]$

$k_1:$

a		b	b / c
---	--	---	-------

Ok! We might not observe all



$k_1:$

a		b	b / c
---	--	---	-------

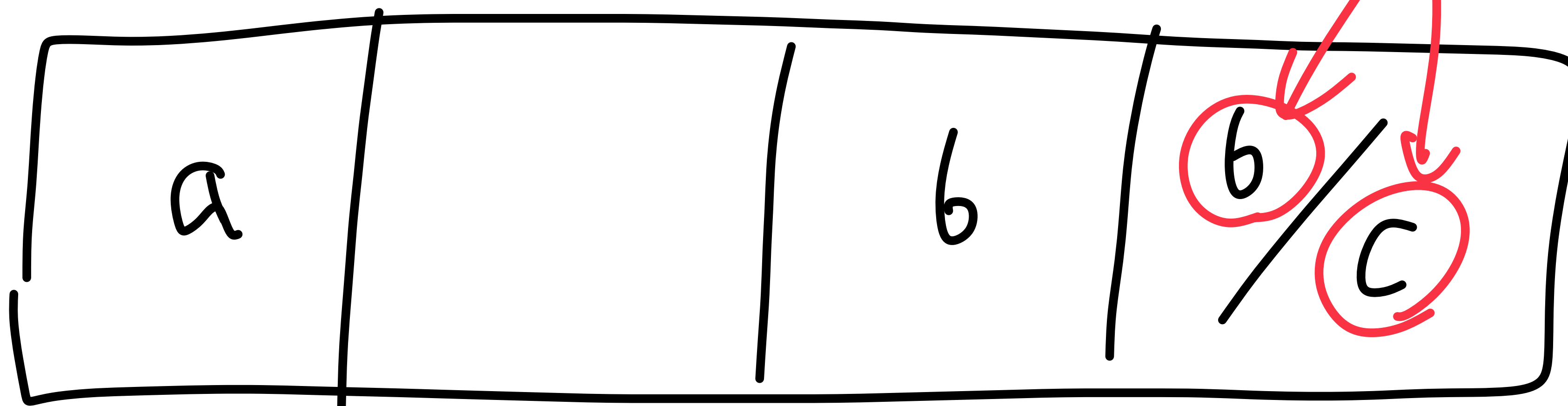
$k_1:$

a		b	b/c
---	--	---	-----

duplicate!

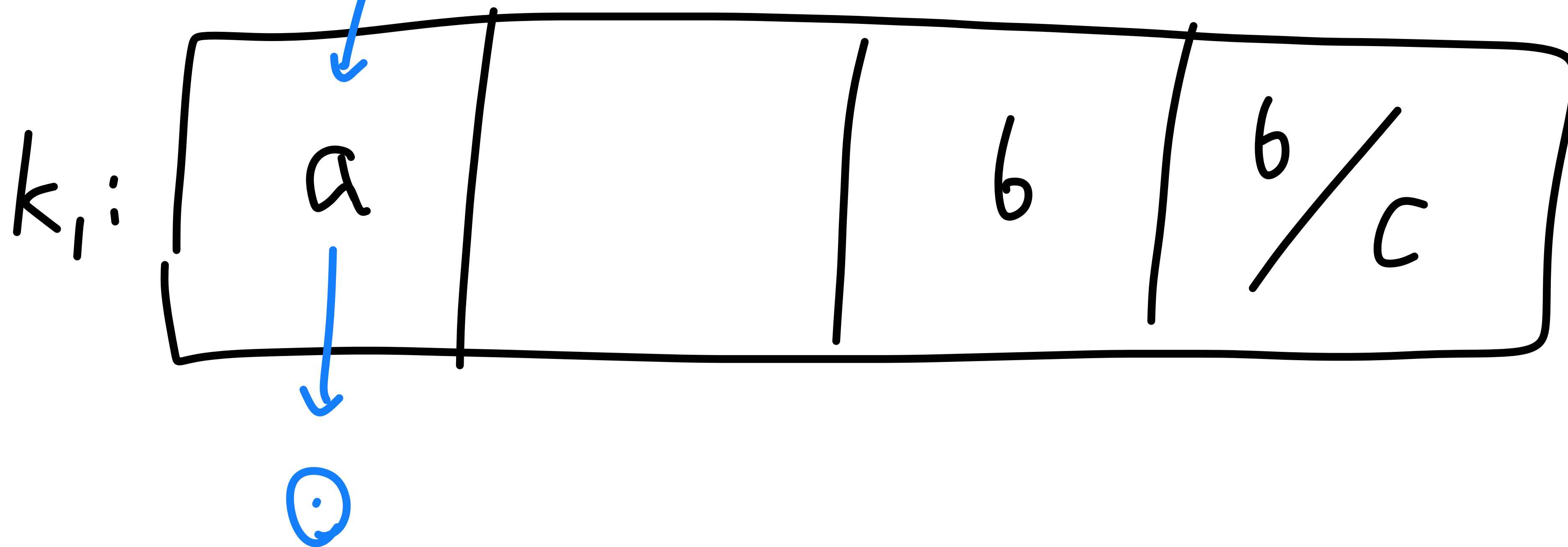


$k_1:$

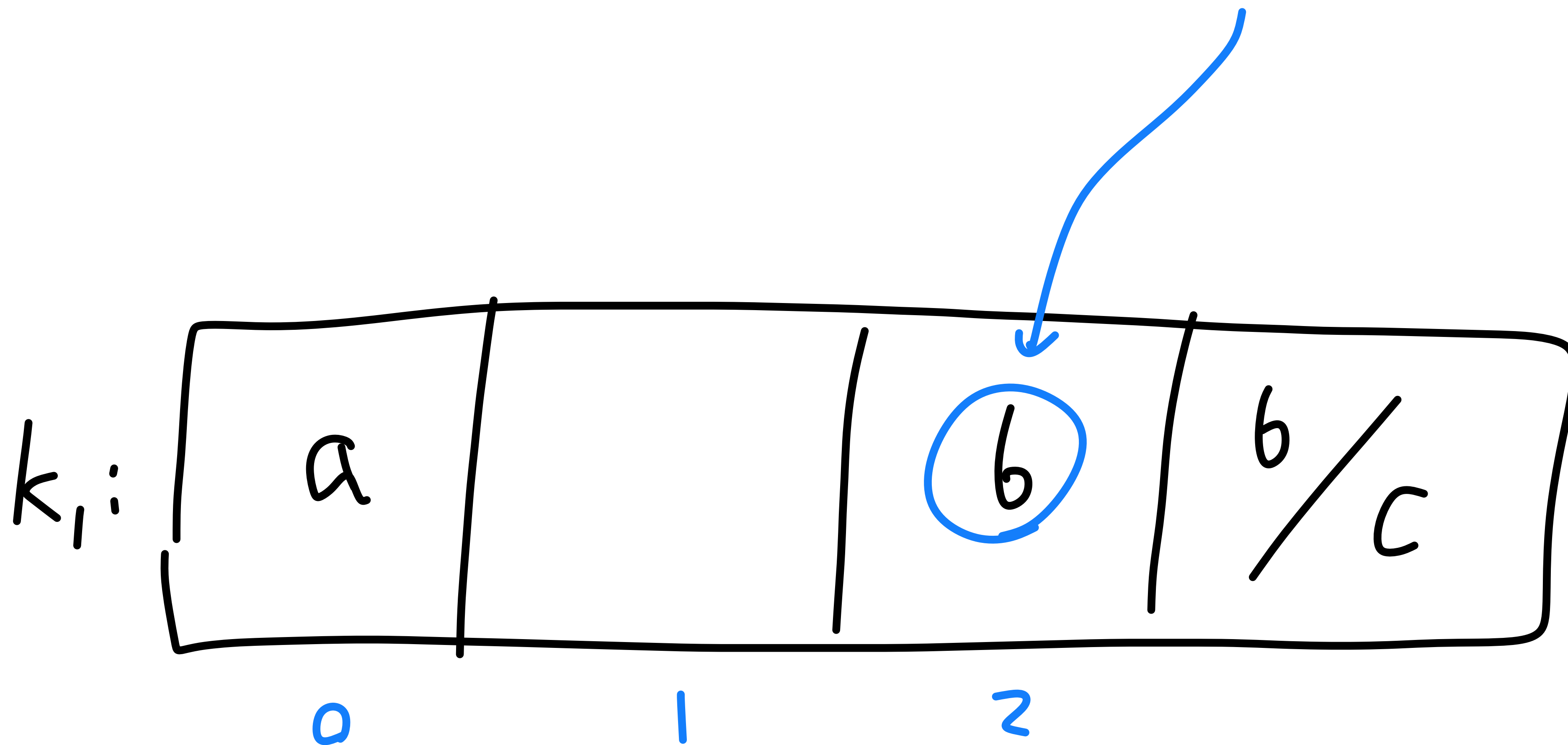


inconsistent offset!

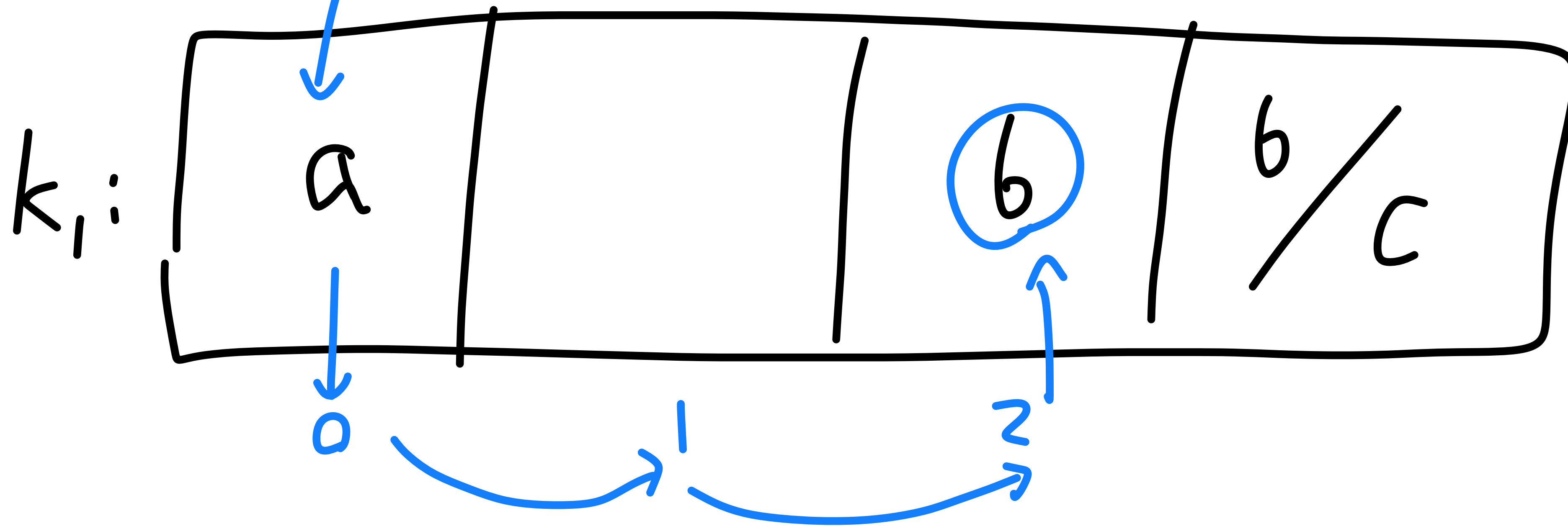
"What offset is a ?"



"What value got offset 2?"



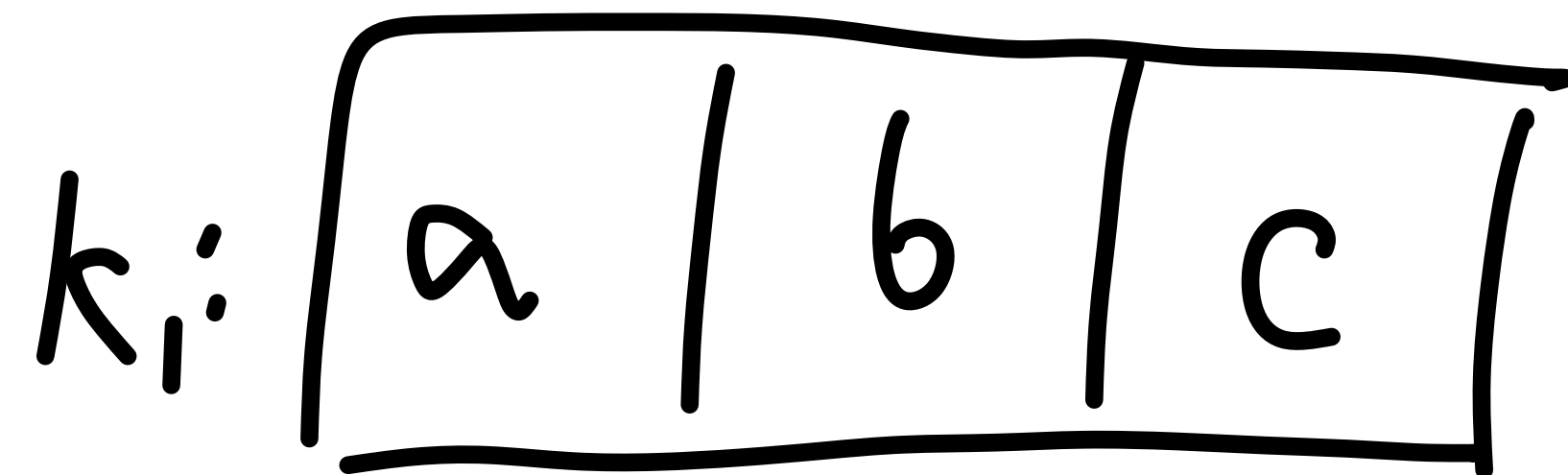
"Which value came just after a?"



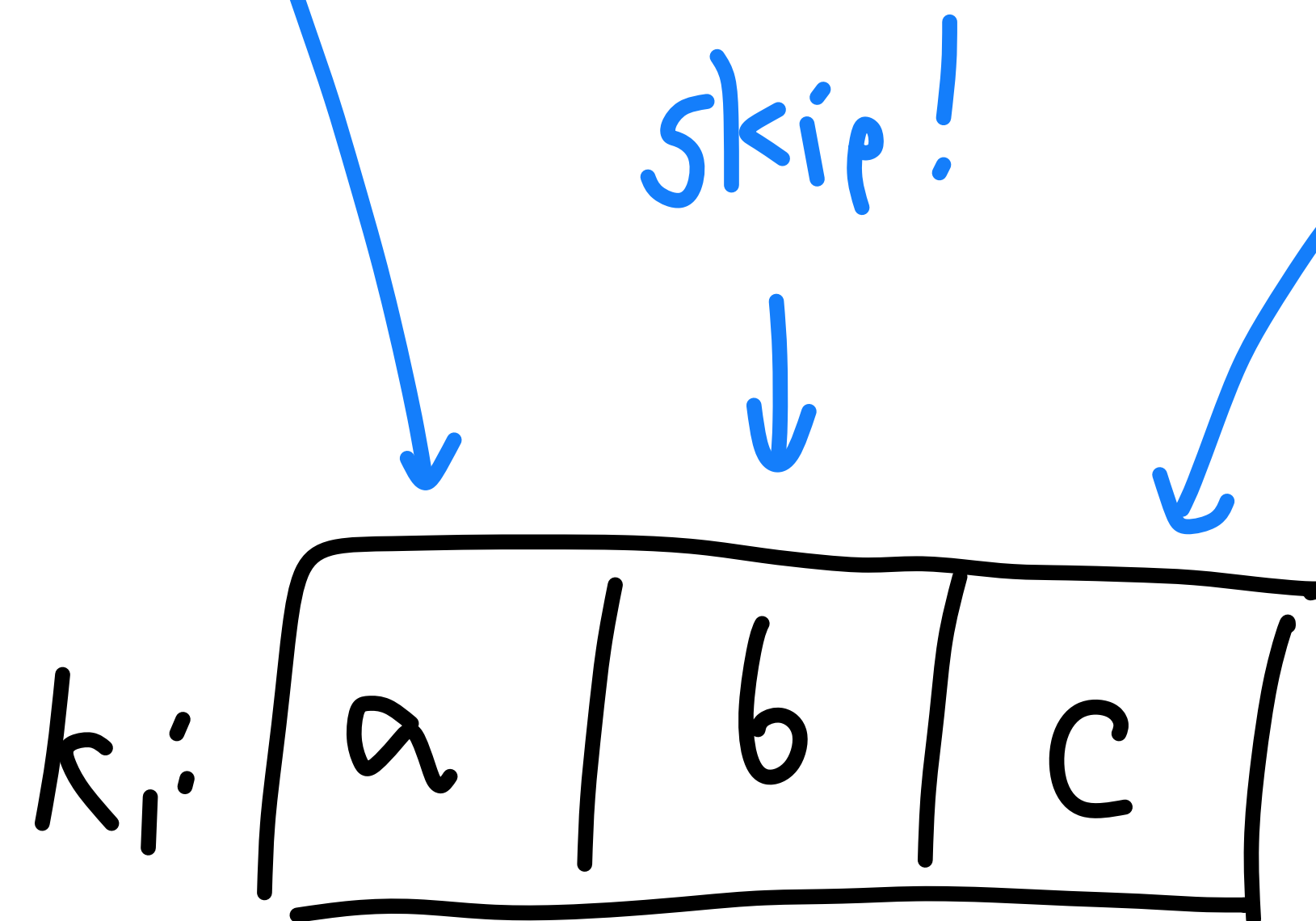
k_i :

a	b	c
---	---	---

$[\text{par} \{ k_i : [0, a] \}], [\text{par} \{ k_i : [2, c] \}]$



$[\text{poll } \{k_i : [0, a]\}], [\text{poll } \{k_i : [2, c]\}]$



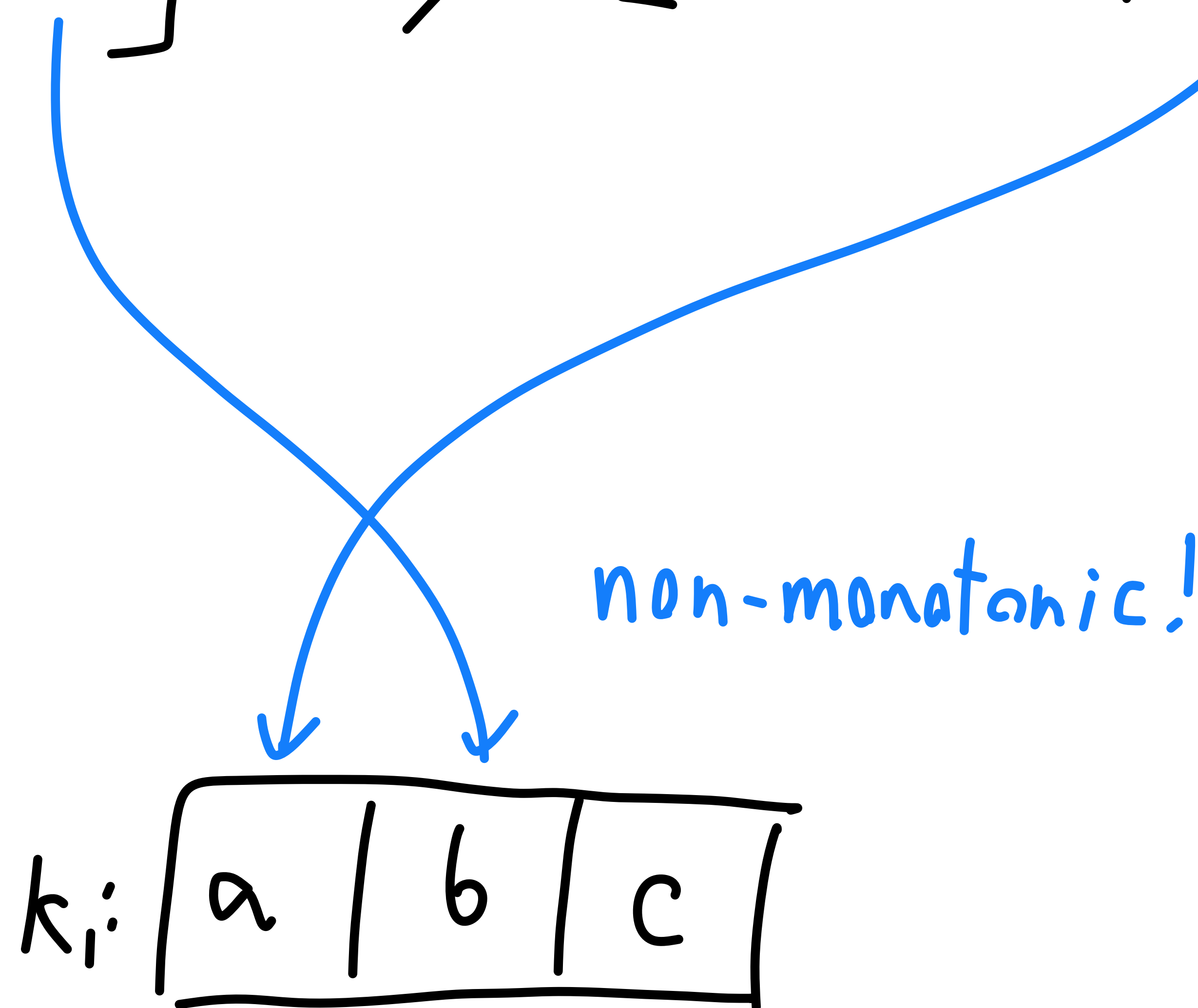
$[\text{poll } \{ k_i : [[1, b]] \}] , [\text{poll } \{ k_i : [[0, a]] \}]$

non-monotonic!

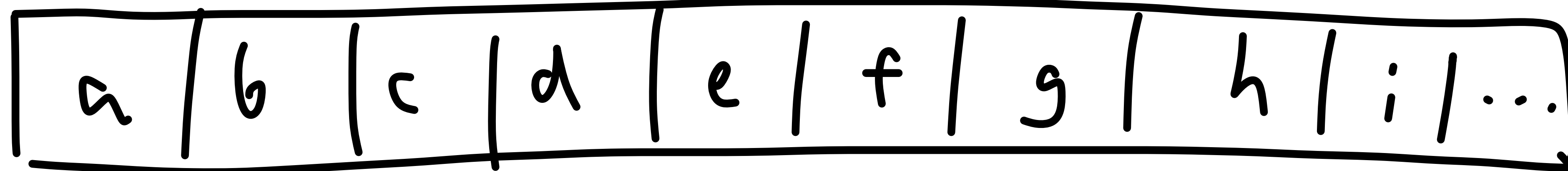
k_i :

a	b	c
---	---	---

$[\text{send } k, b]$, $[\text{send } k, a]$



a	b	c	d	e	f	g	h	i	...
---	---	---	---	---	---	---	---	---	-----



///

////////

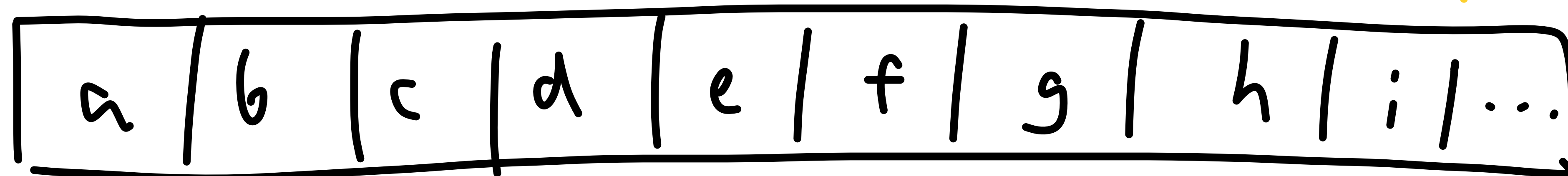
polled by C₁

////////

////////

polled by C₂

Unseen



///

////////

polled by C₁

////////

////////

polled by C₂

Lost!

Unseen



///

////////

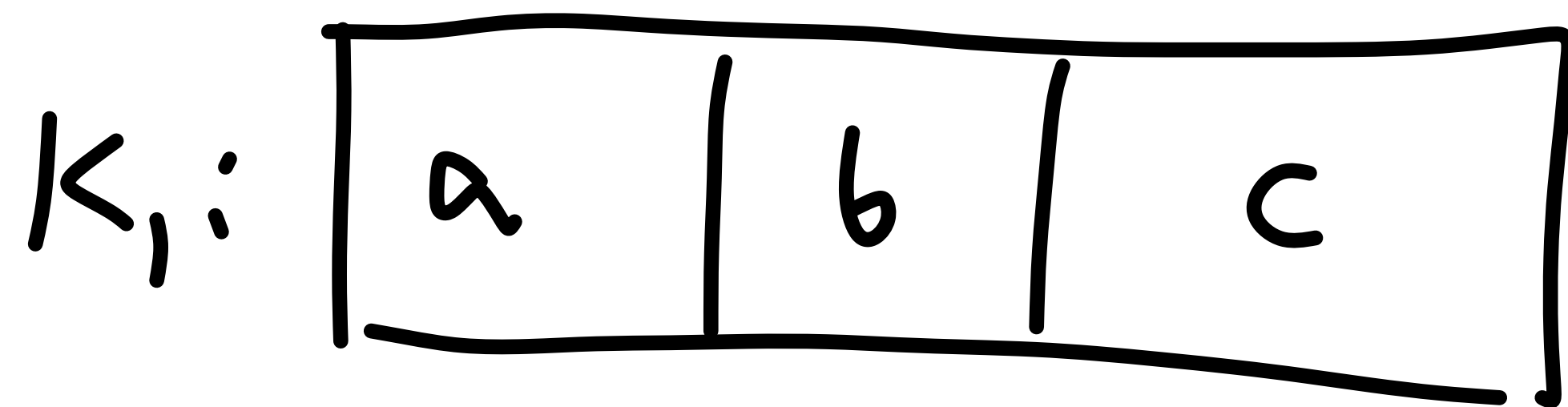
polled by C₁

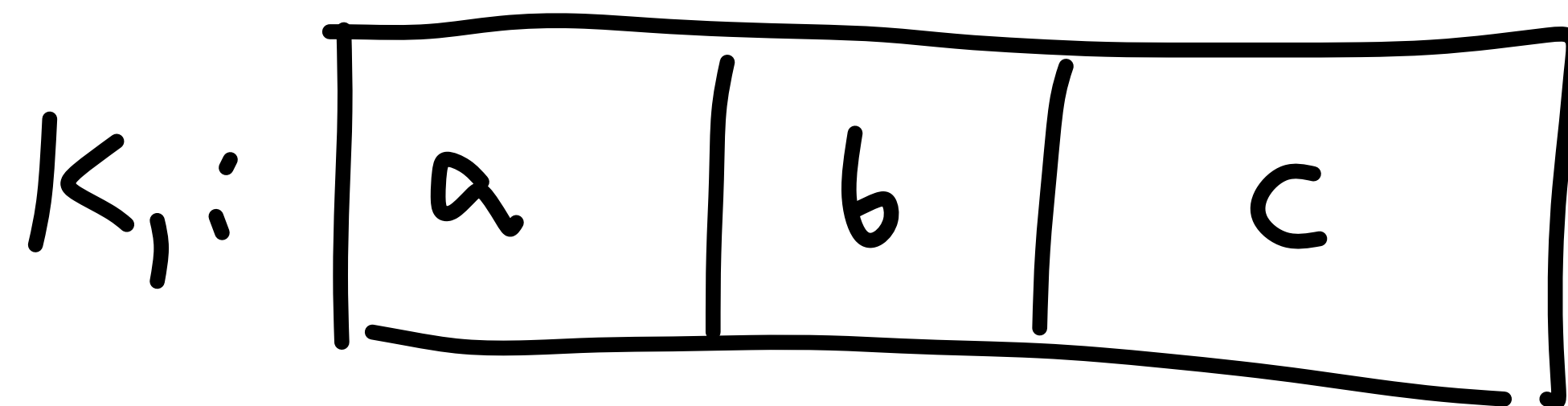
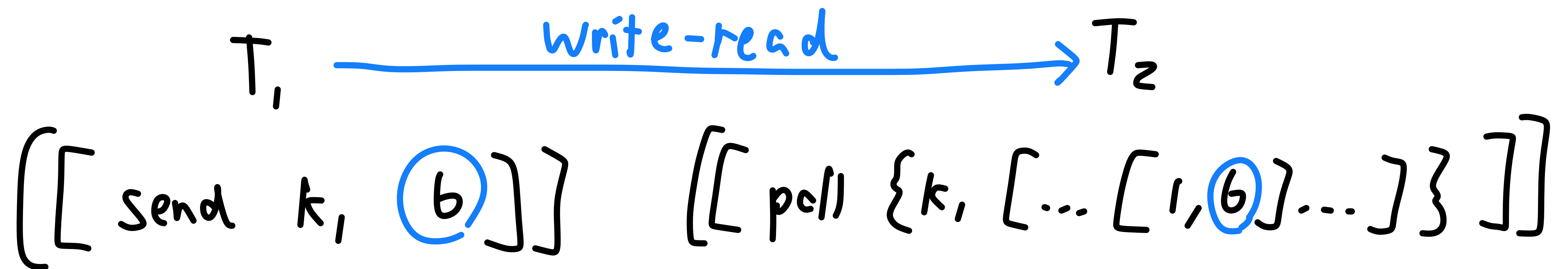
////////

////////

polled by C₂

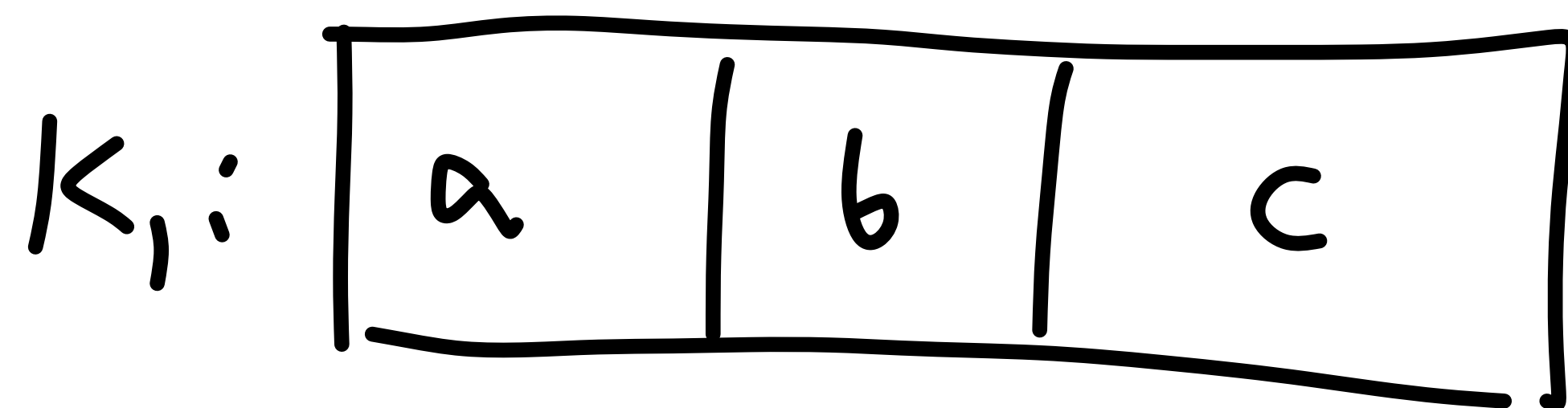
T_1
 T_2
 $\left[\left[\text{send } k, b \right] \right] \quad \left[\left[\text{poll } \{k, [\dots [1, b] \dots] \} \right] \right]$





T_1
[send k, b]

T_2
[[send k, c]]



write - write

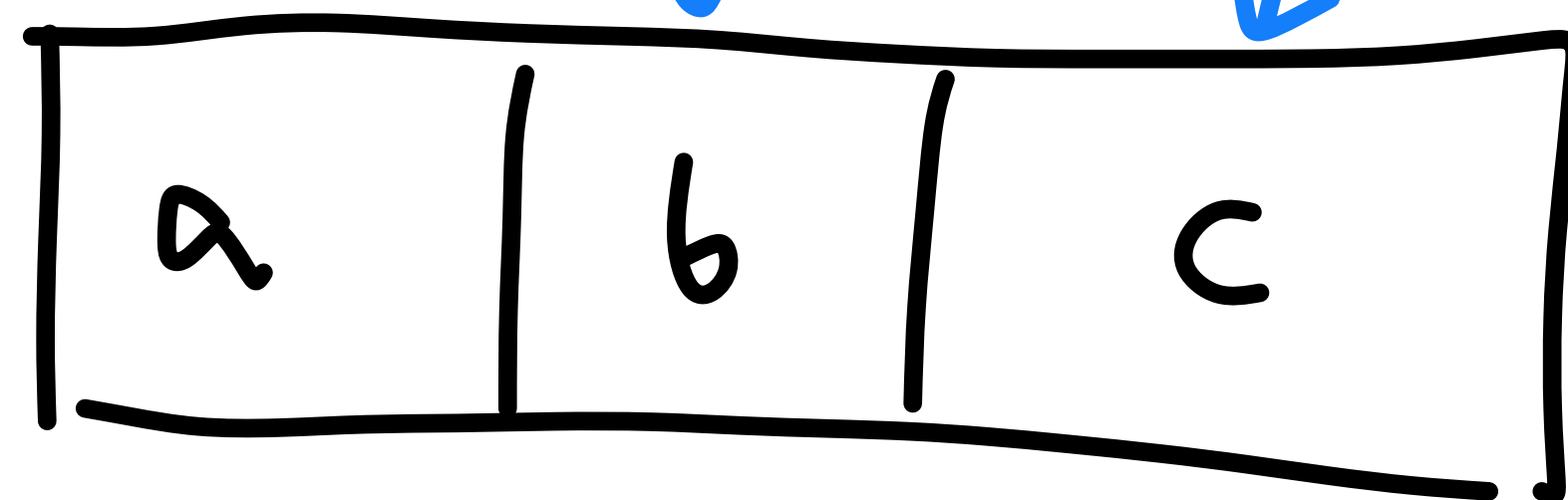
T_1

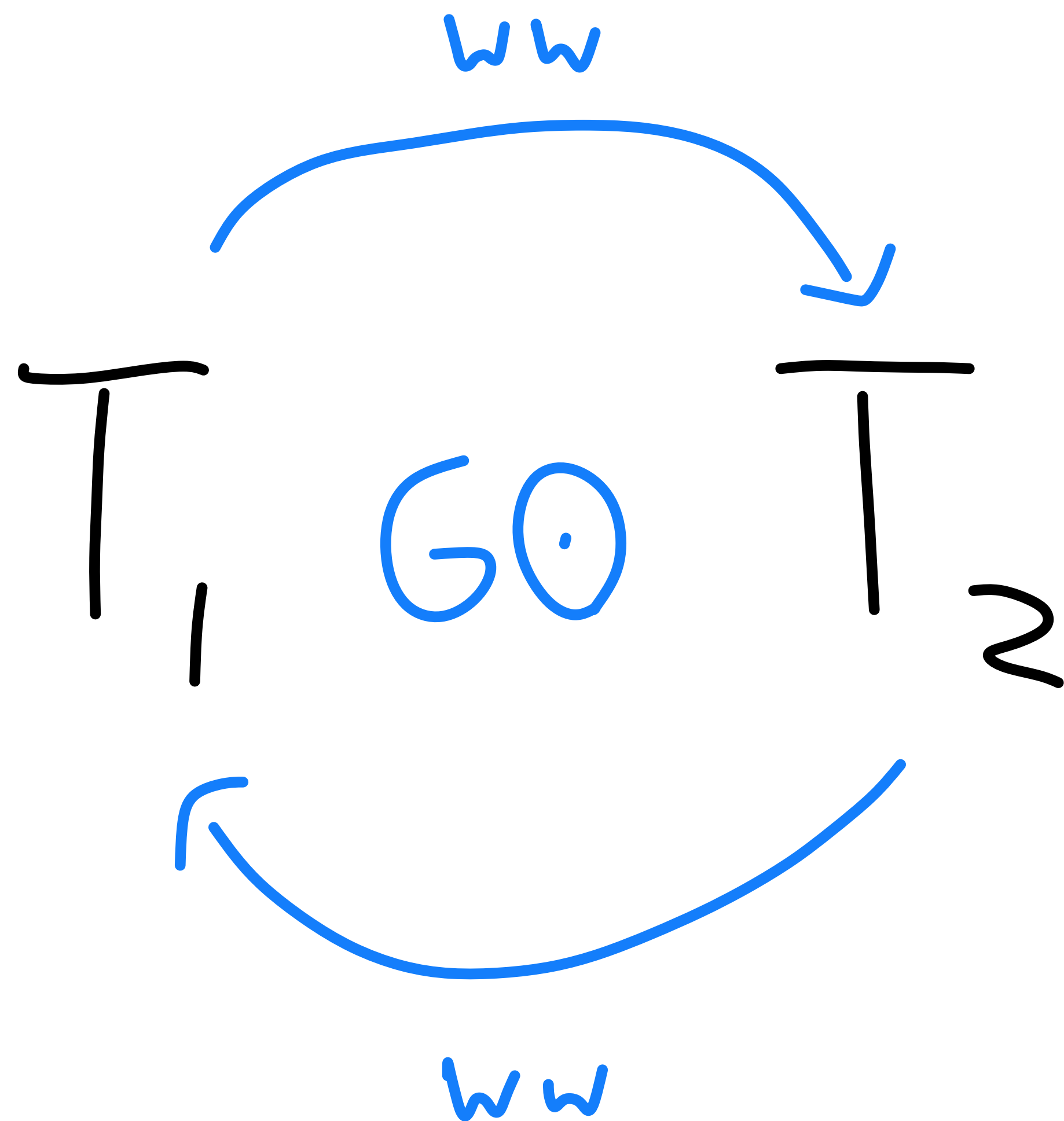
T_2

$[\text{send } k, b]$

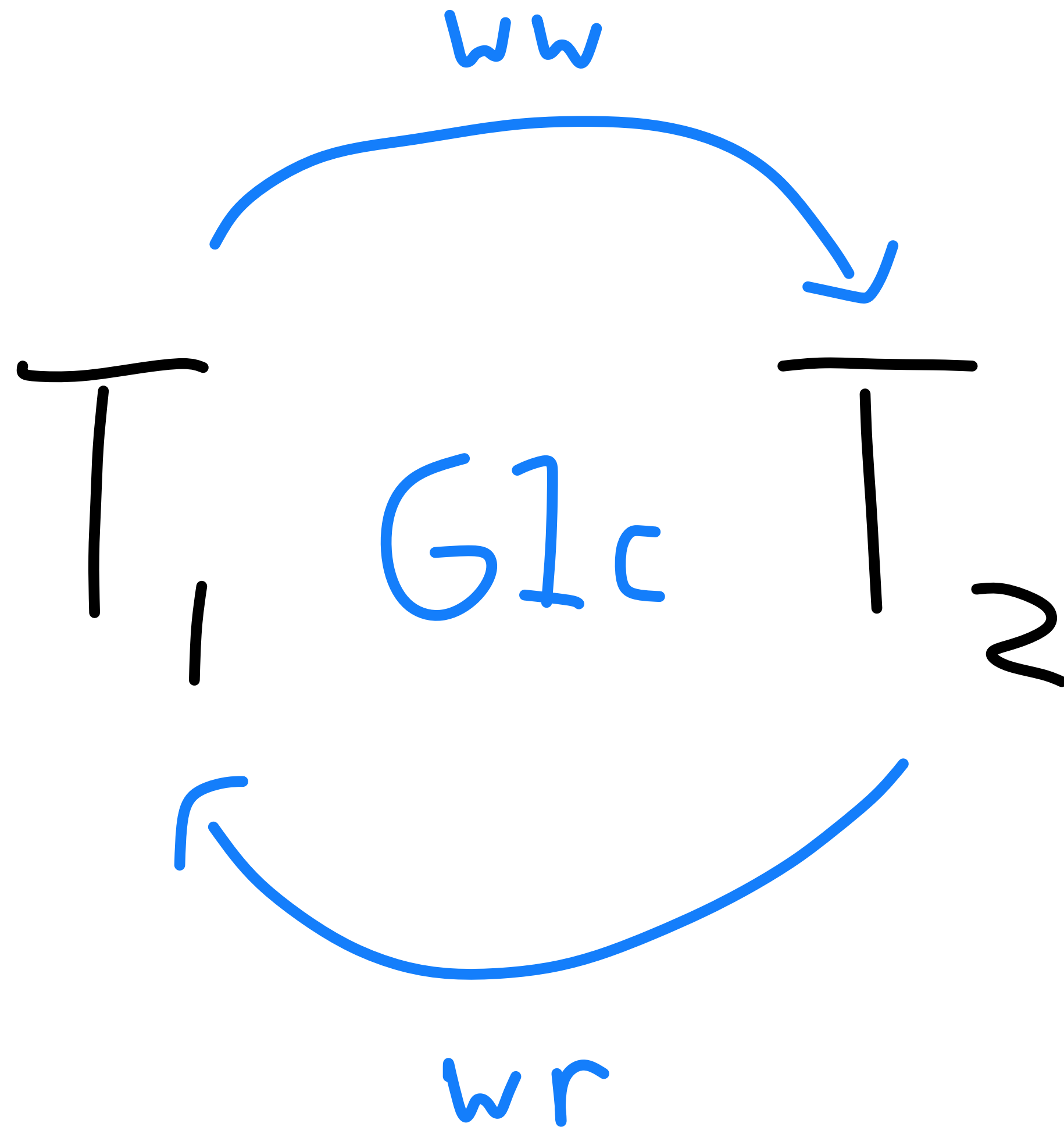
$[\text{send } k, c]$

$K_1:$





Write
Cycle



Circular
Information
Flow

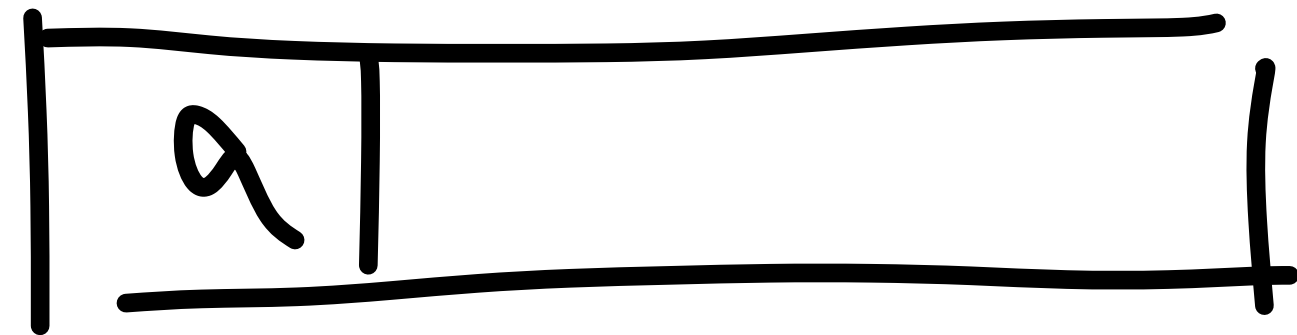
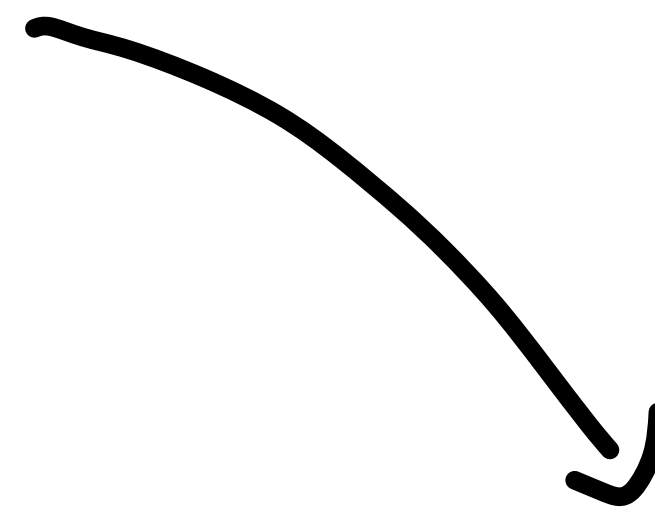
Faults

- Crash
- Pause
- Clock skew
- Network Partition
- Membership change

Results

#1 Duplicate Writers by Default

producer.send(a)



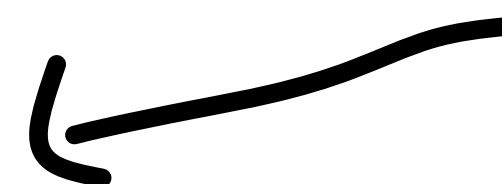
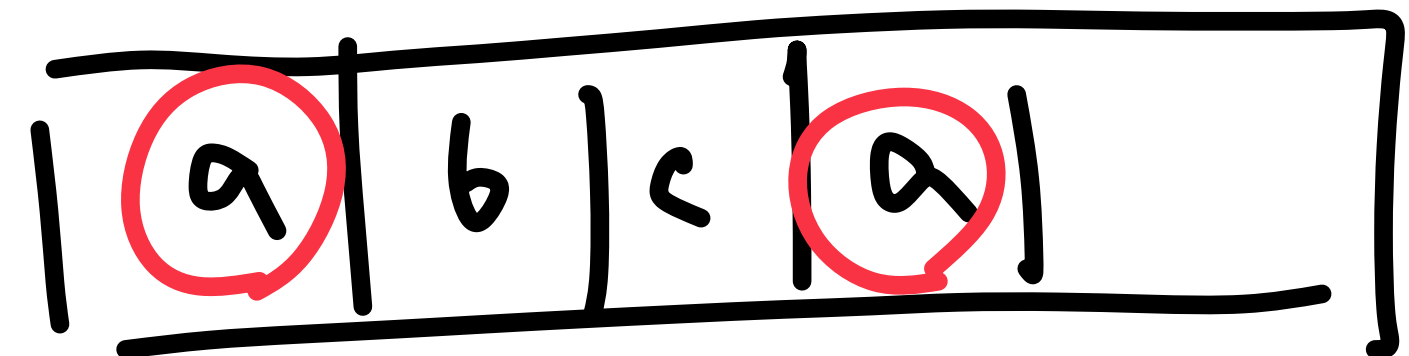
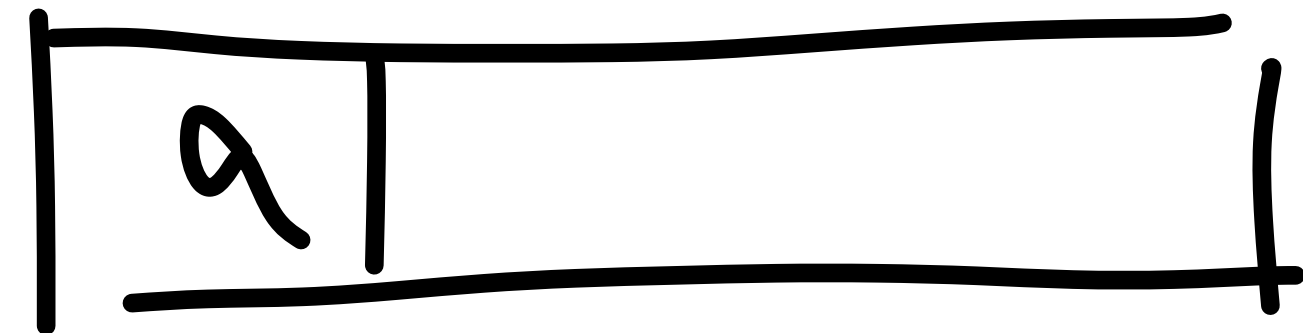
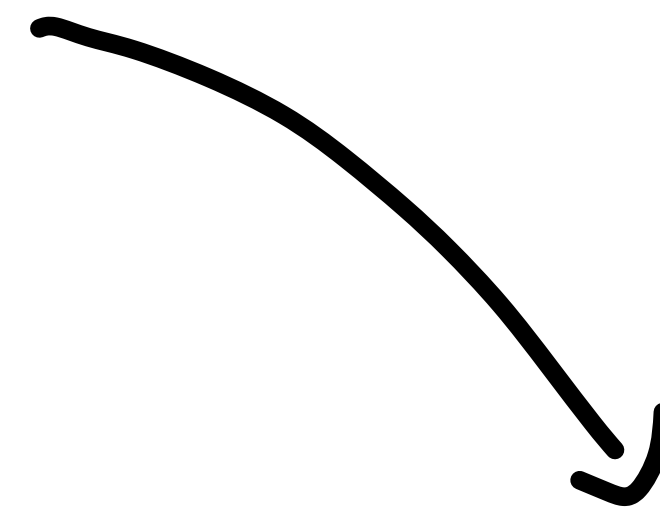
#1 Duplicate Writers by Default

producer.send(a)

⋮

(retry)

ok!



#1 Duplicate Writers by Default

`enable.idempotence = true`

(default in 3.0.0 client)

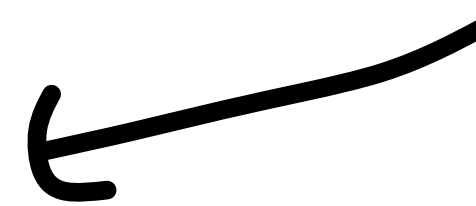
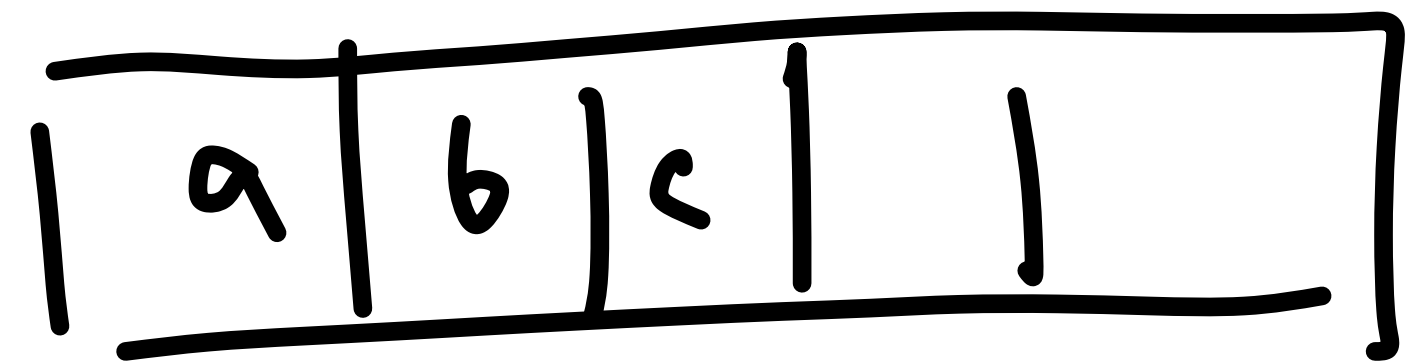
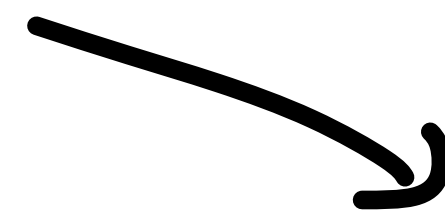
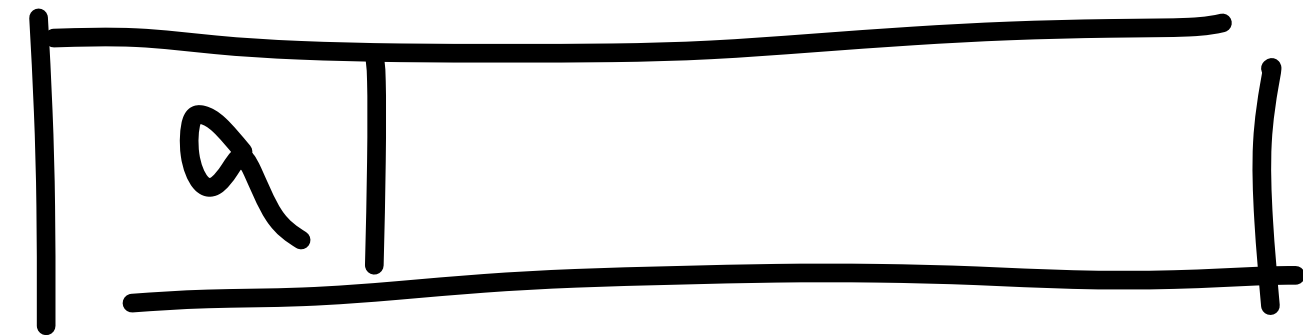
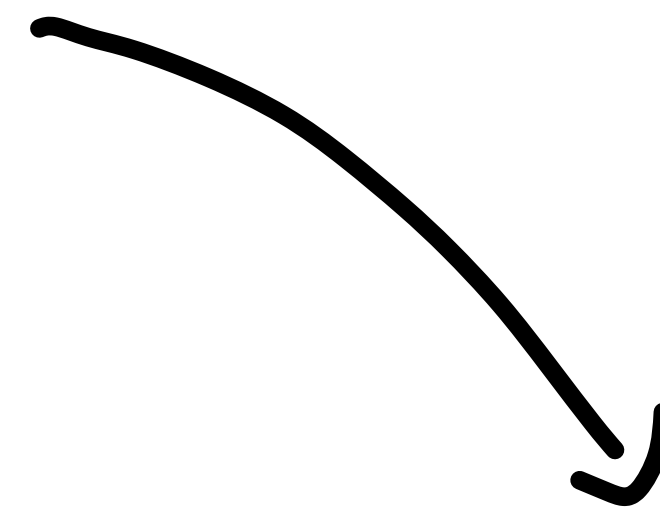
#1 Duplicate Writers by Default

producer.send(a)

⋮

(retry)

ok!



no dup!

#1 Duplicate Writers by Default

Except... we still saw dups

w/any crash, pause, partition...

#1 Duplicate Writers by Default

Partly a client bug!

Redpanda server needs:

- enable_idempotence true
- id_allocator_replication 3

#1 Duplicate Writers by Default

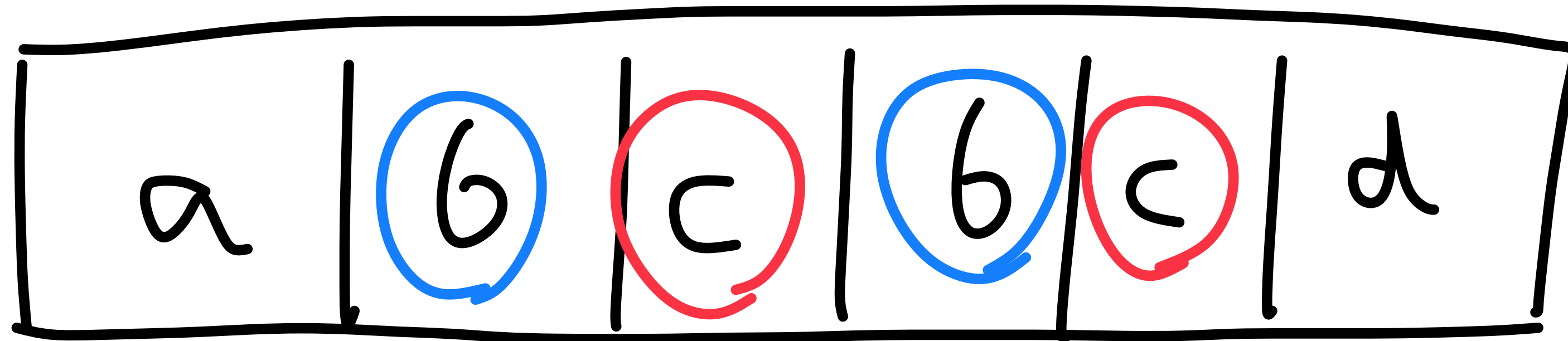
Should be fixed by enabling

idempotence by default

Planned for 22.1.1

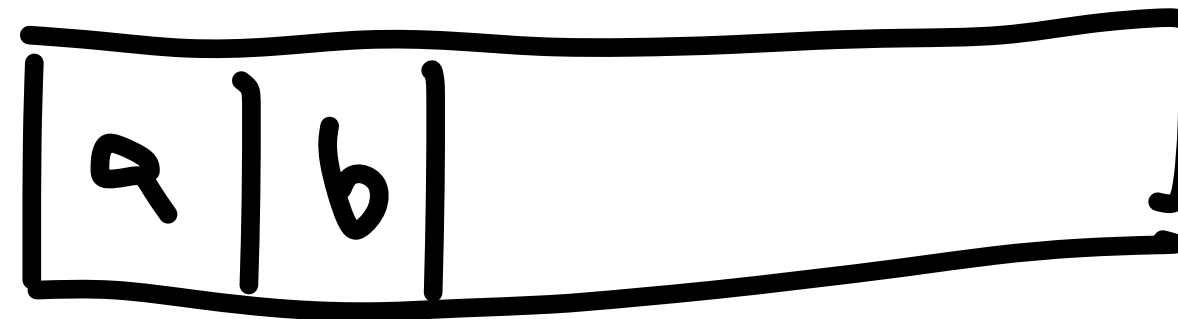
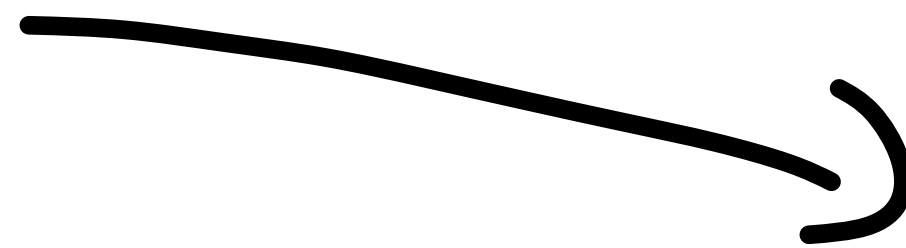
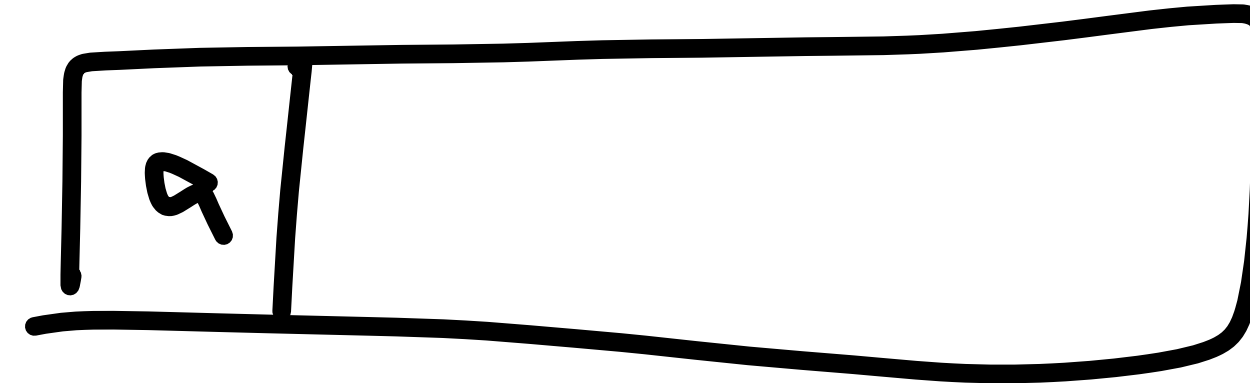
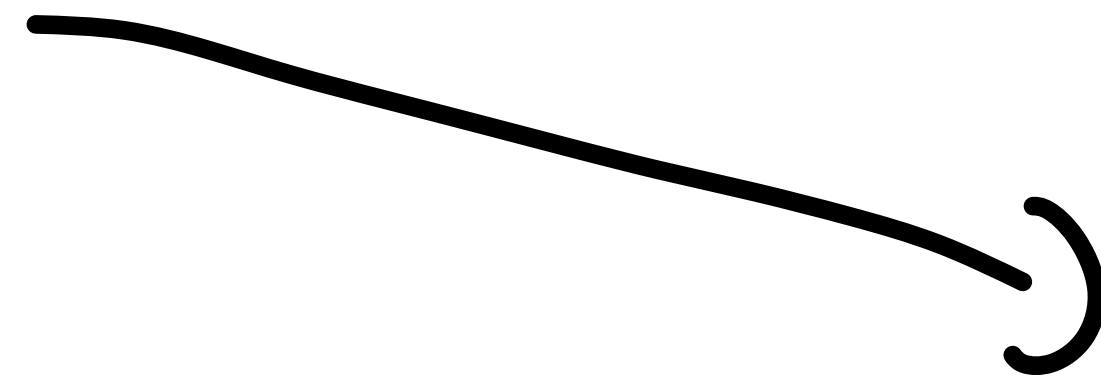
#3039 Duplicate Writers Not by Default

We still saw duplicate messages



#3039 Duplicate Writers Not by Default

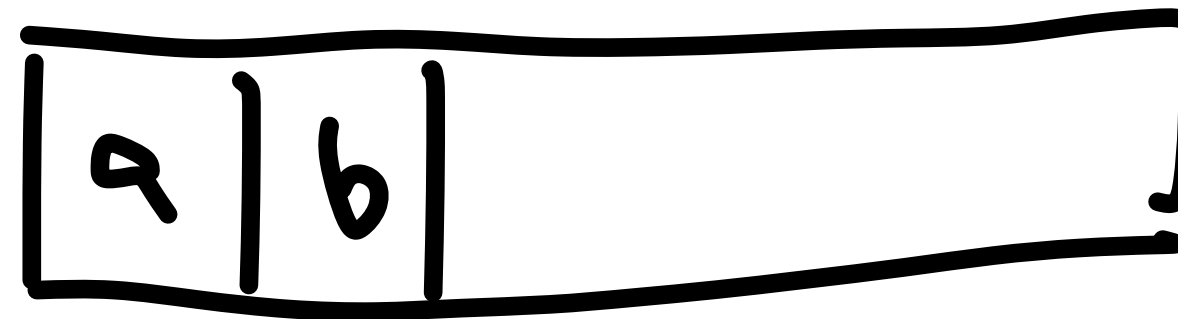
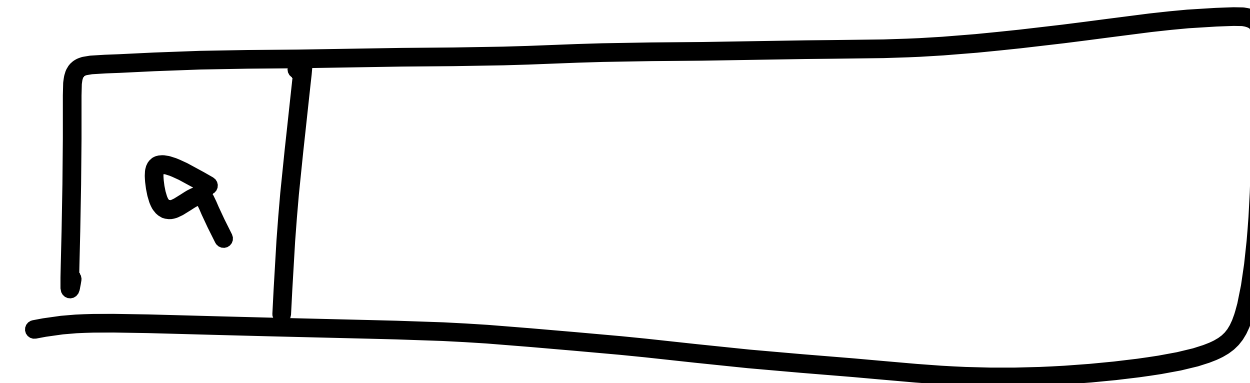
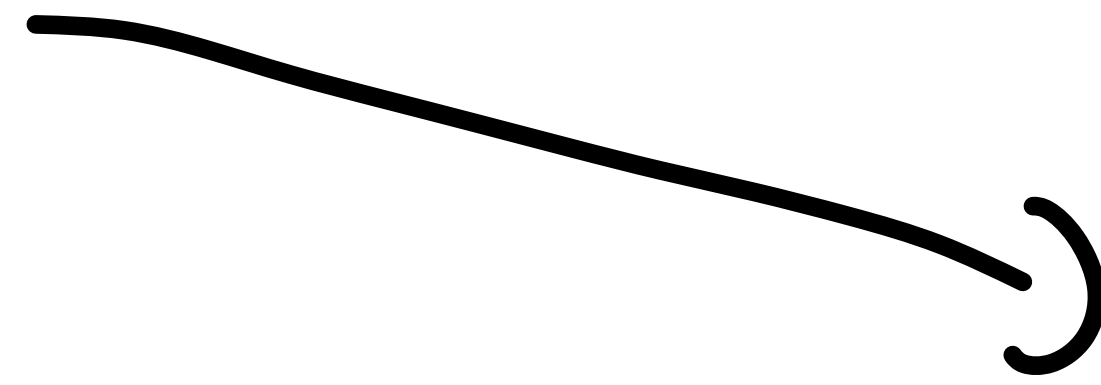
send(a)
⋮
(retry)



Out Of Order Sequence Exception

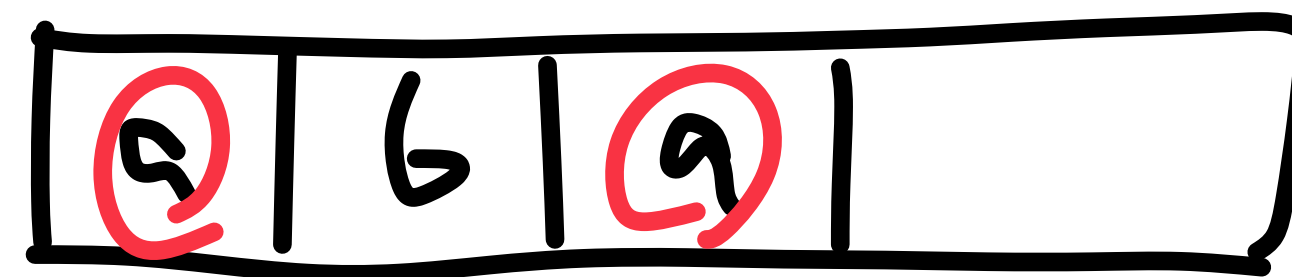
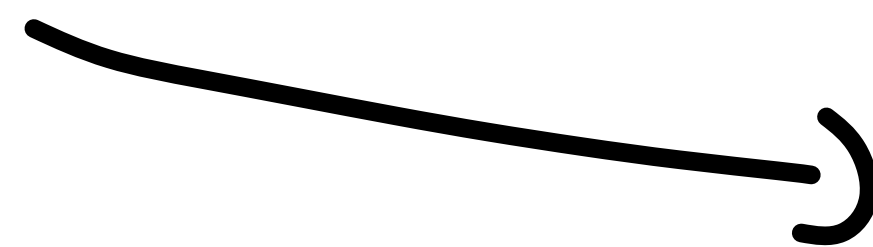
#3039 Duplicate Writers Not by Default

send(a)
⋮
(retry)



Out Of Order Sequence Exception

(retry)



#3039 Duplicate Writes Not by Default

— 21.10.2 : Fixed by adding deduplication

— 21.10.3 : Performance improvements

#3335 Assert Failure Deallocating Partitions

With membership changes...

Assert failure... allocated partitions >

allocation capacity

Still investigating

3336 Assert failure in response.partition_index

w/processor crashed...

Assert failure: ... Response & current

partition ids have to be the same.

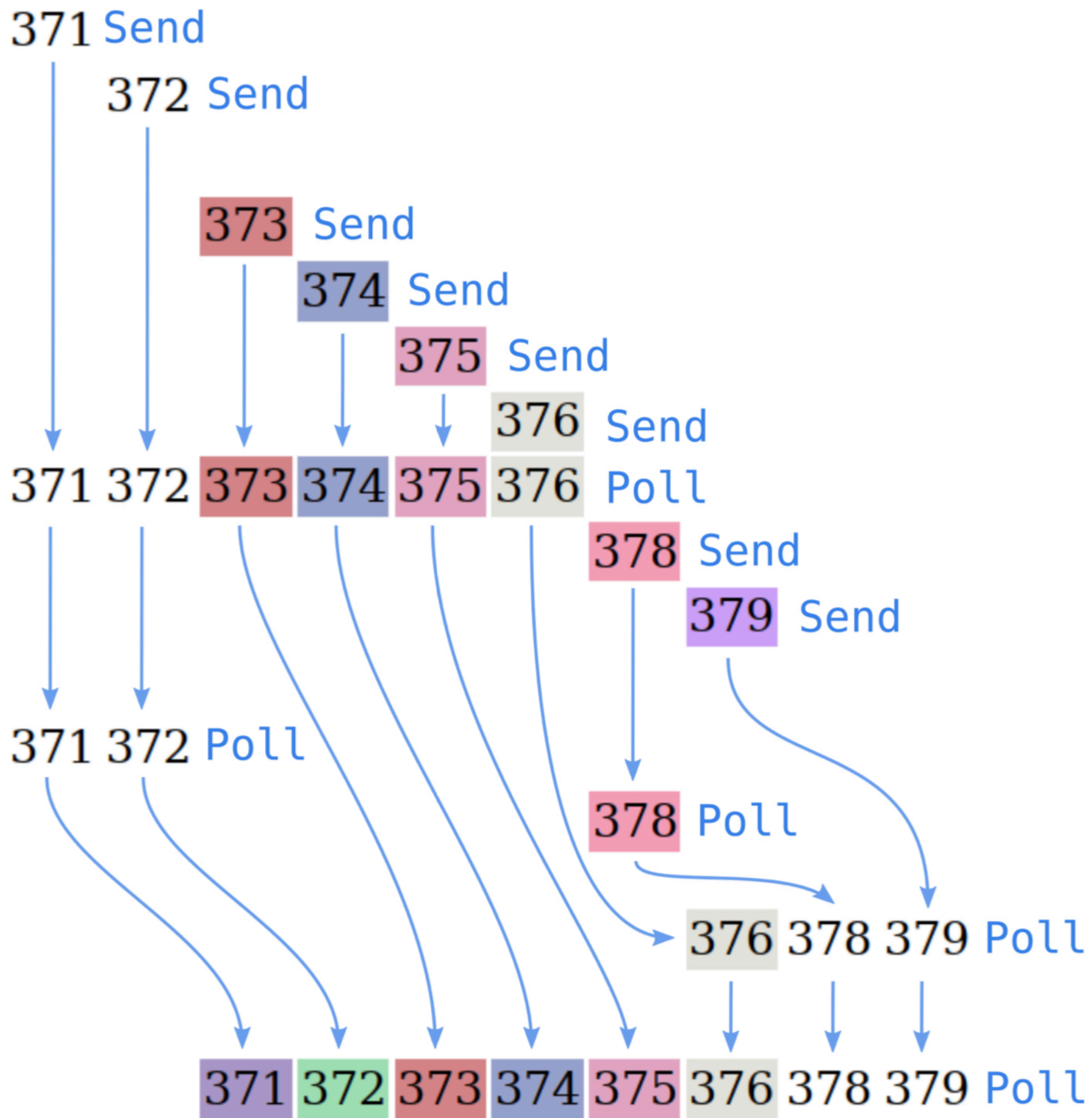
Current response 0, update 1

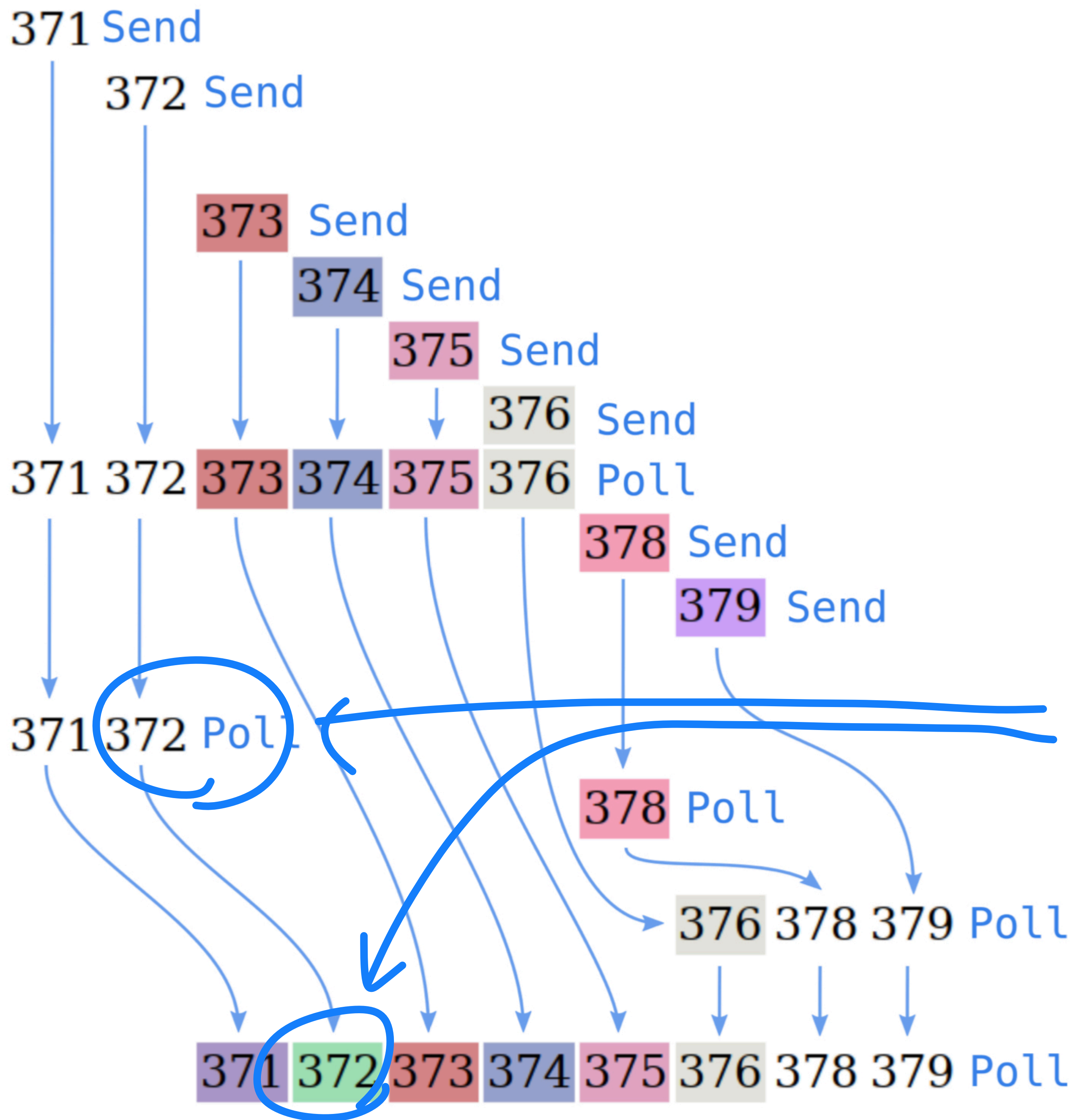
Patched → 22.1.1

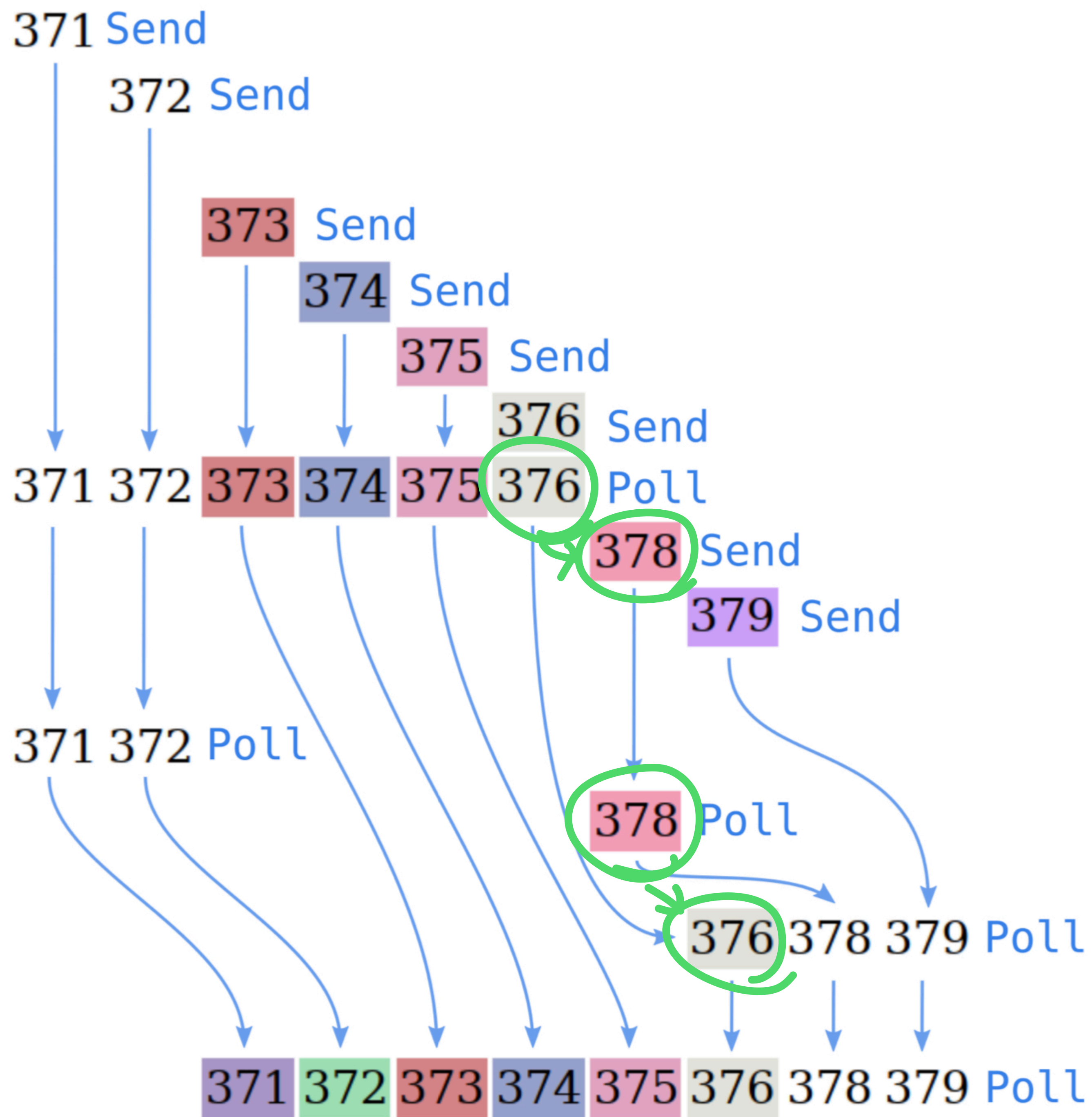
#3003 Incansistent Offsets

w/ pro cers Crashes or

network partitions...







Order inversion

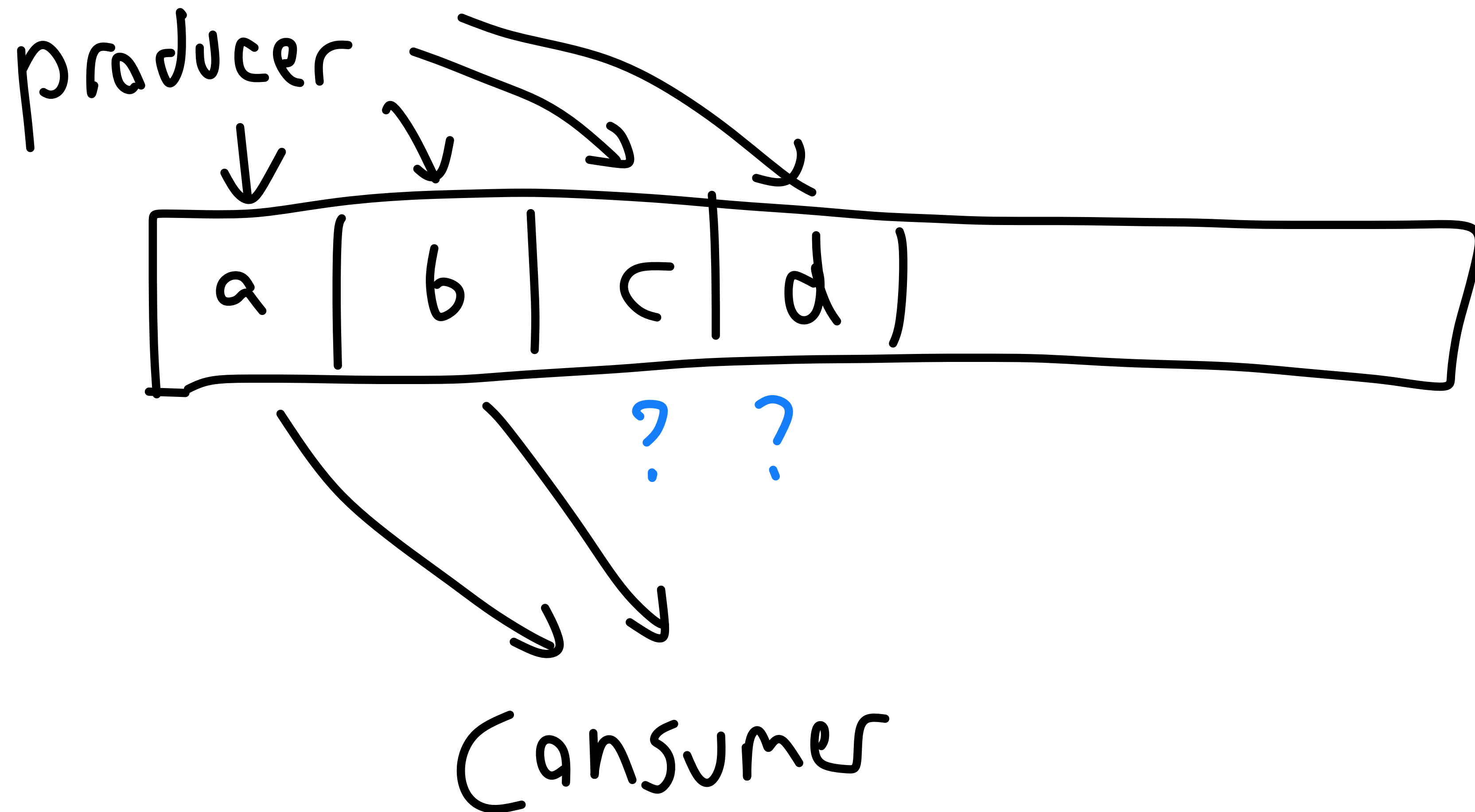
#3003 Incansistent Offsets

Caused by Raft state machine
applying uncommitted log entries

→ Fixed in 21.10.3 by waiting for
commit pointer to advance first

#6 Last / Stale Messages

w/ process pause



#6 Last/Stale Messages

Consumer.poll()

→ OK! Messages! []

Consumer.poll()

→ OK! Messages! []

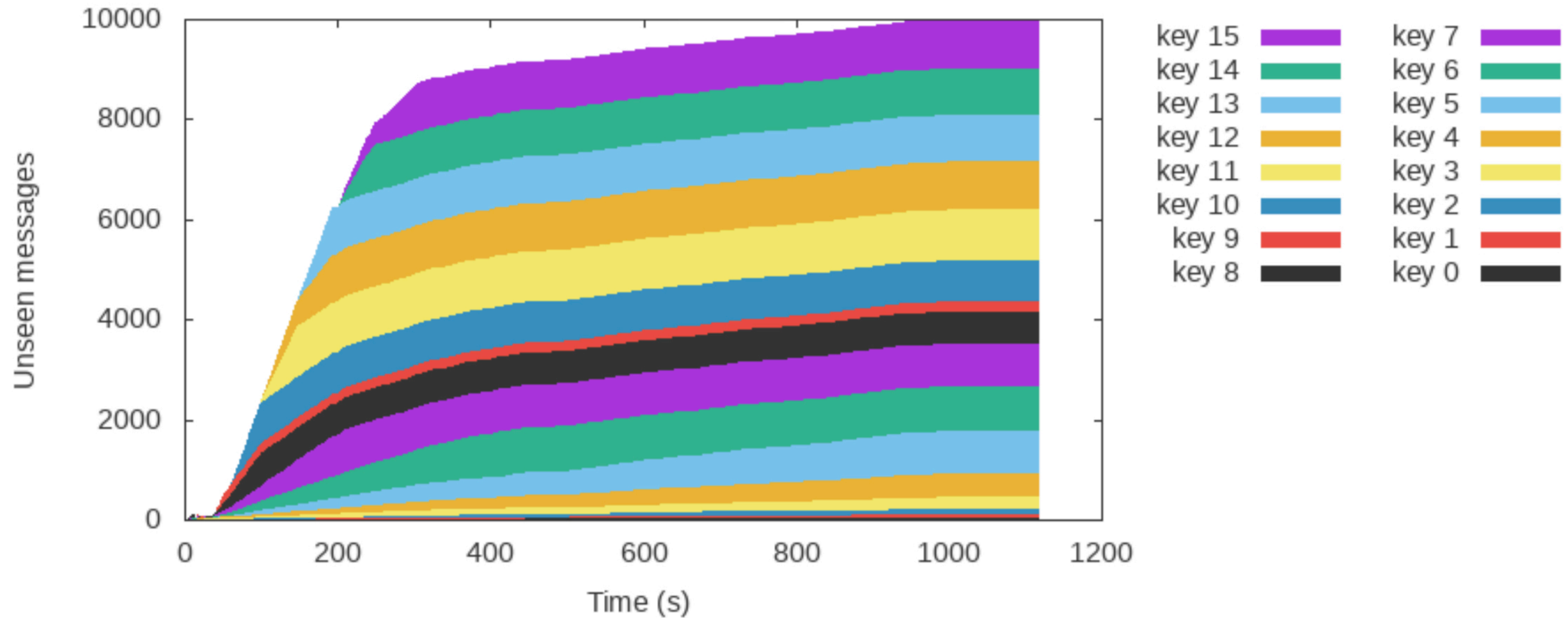
Consumer.poll()

→ OK! Messages! []



#6 Last / Stale Messager

21.11.2 queue assign acks=all retries=0 aor=earliest membership,partition unseen



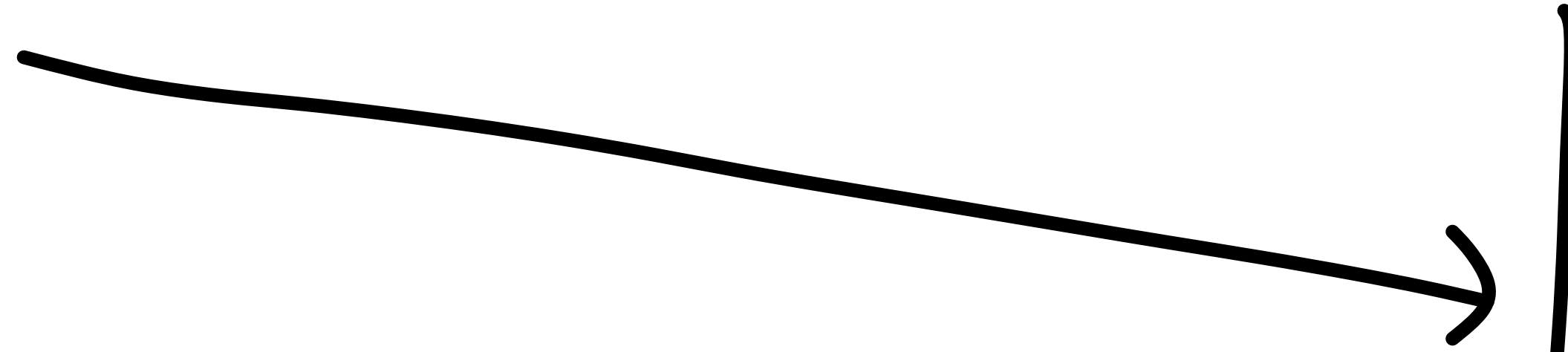
#6 Last/Stale Messages

Still on disk... just not being
delivered...

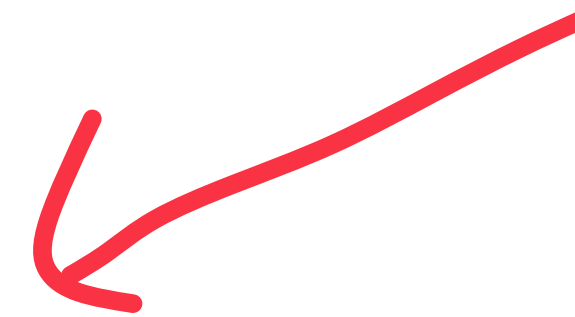
Redpanda still investigating!

#7 Aborted Read — Not Leader Or follower Exception

send(r)



Not Leader Or follower Exception



#7 Aborted Read — Not Leader Or Follower Exception

"This failed, right?"

"I think so..."

#7 Aborted Read — Not Leader Or follower Exception

Broker returns this error if a request could not be processed because the broker is not the leader or follower for a topic partition. This could be a transient exception during leader elections and reassignments. For Produce and other requests which are intended only for the leader, this exception indicates that the broker is not the current leader.

#7 Aborted Read — Not Leader Or follower Exception

Kafka team:

"Actually, it's indeterminate!

Could have succeeded!"

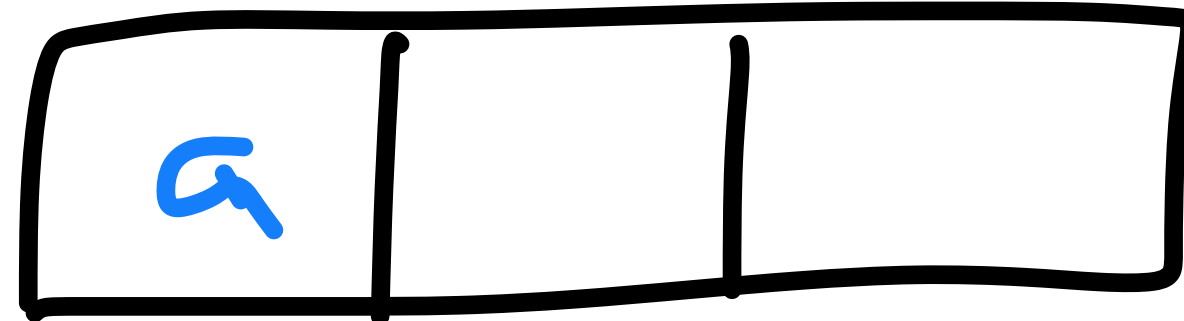
#7 Aborted Read — Not Leader Or follower Exception

Just needs docs!

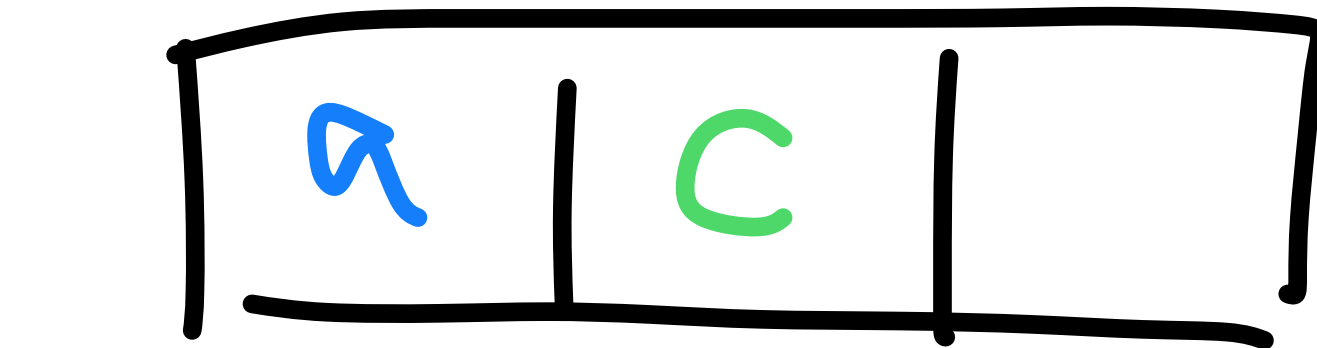
#8 Write Cycles

T_1 : send(a)

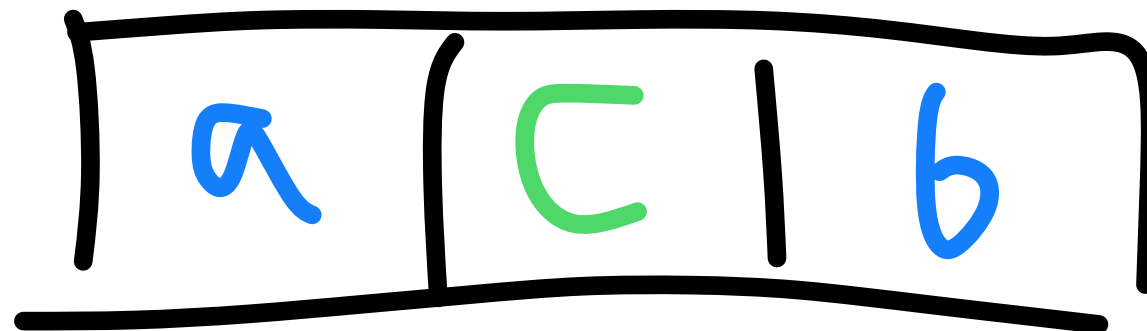
ak!



send(b)



ak!



Commit

ak!

T_2 :

send(c)

ak!

Commit

ak!

#8 Write Cycles

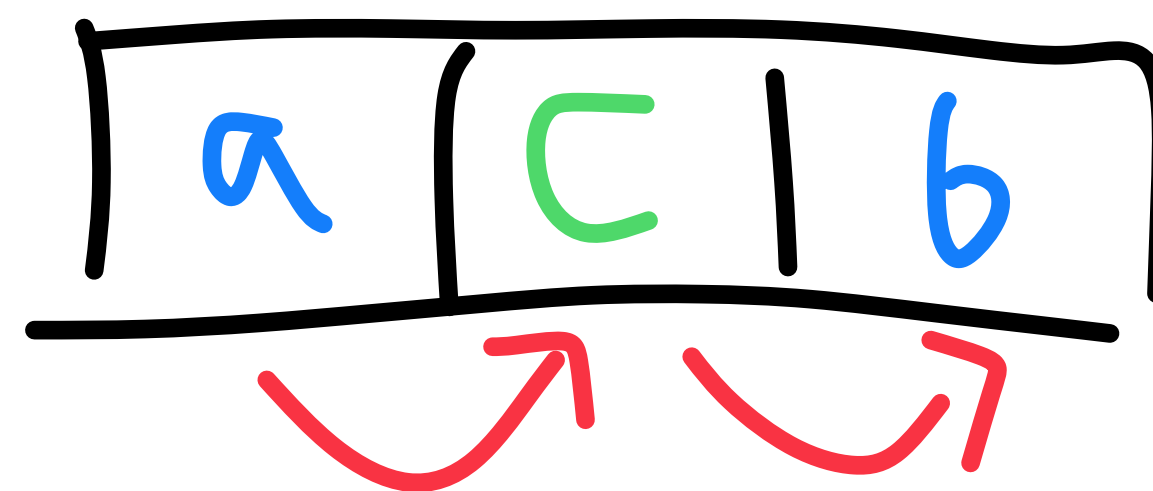
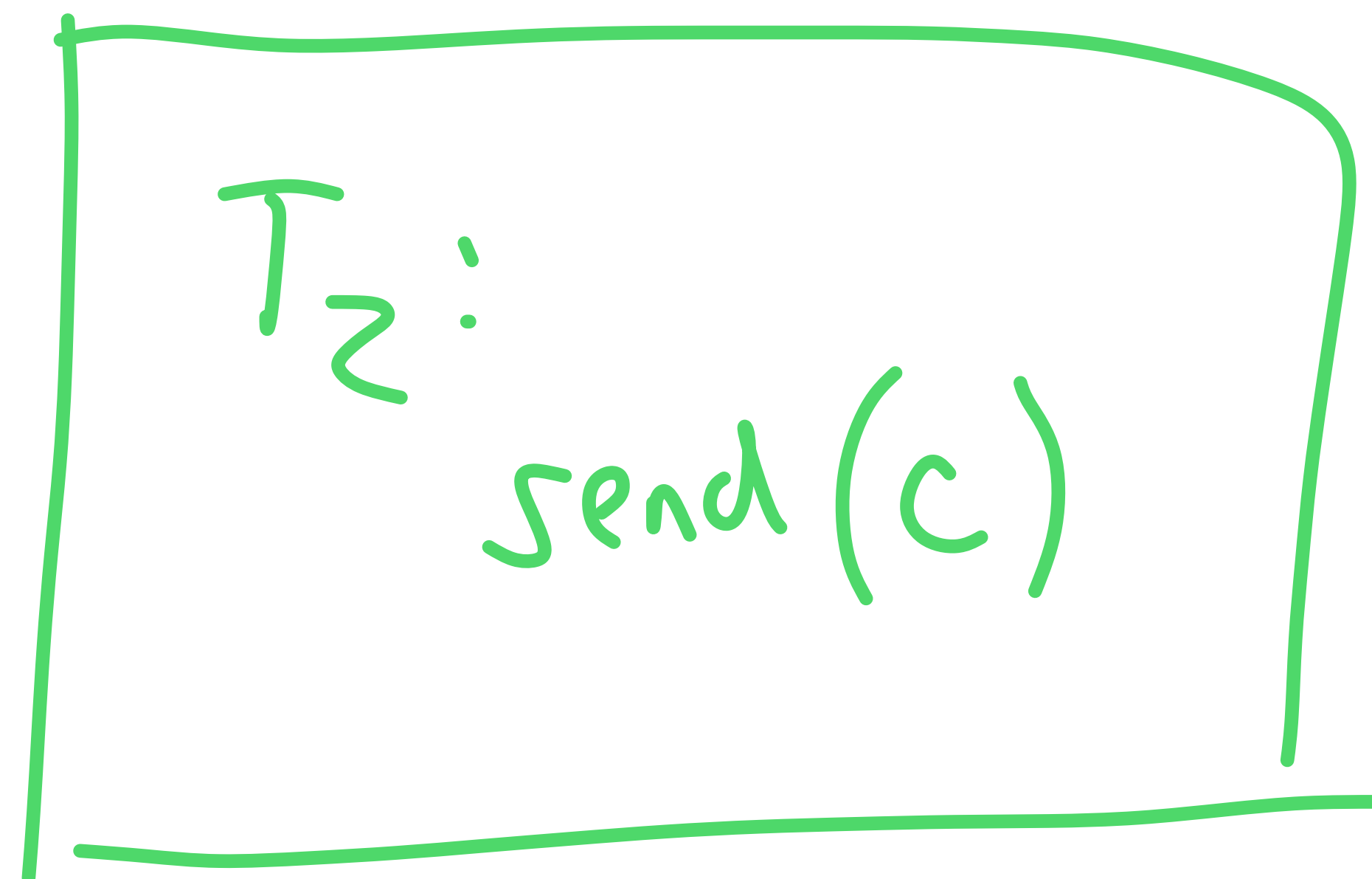
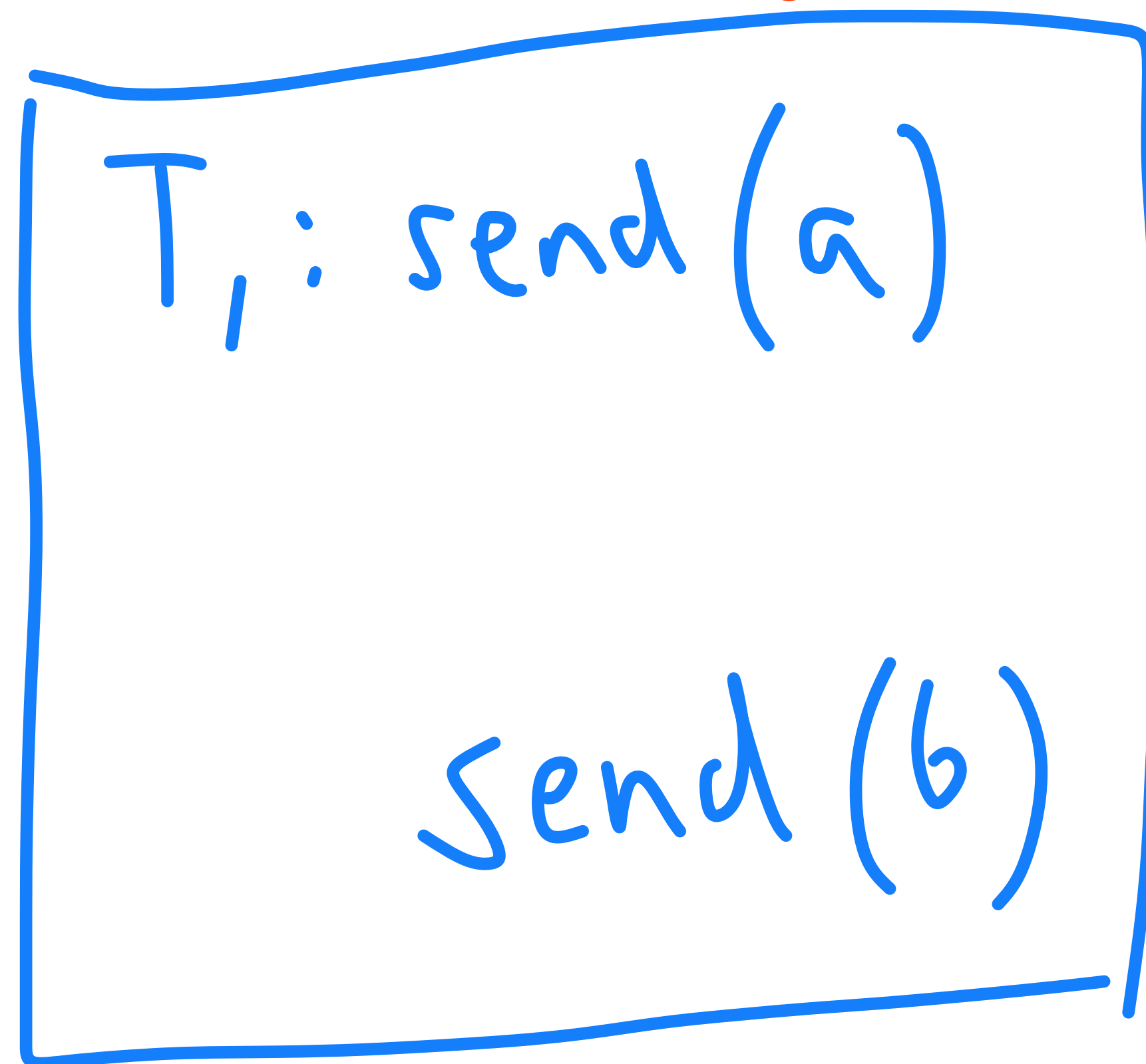
T_1 : send(a)

send(b)

T_2 :
send(c)

a	c	b
---	---	---

#8 Write Cycles



#8 Write Cycles

GO happens constantly in healthy

Redpanda & Kafka clusters...

Forbidden by Adya read-uncommitted!

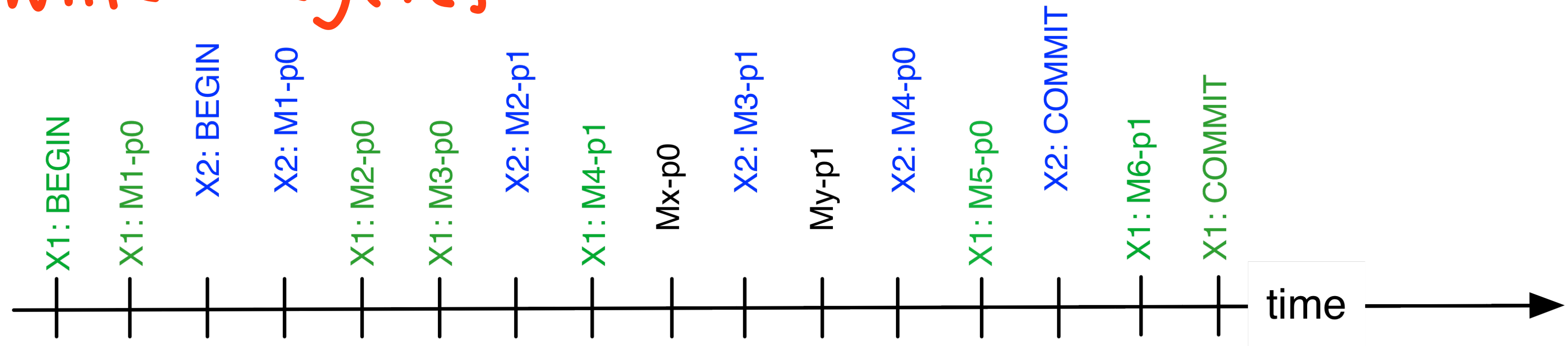
#8 Write Cycles

GO happens constantly in healthy

Redpanda & Kafka clusters...

Forbidden by Adya read-uncommitted!

#8 Write Cycles



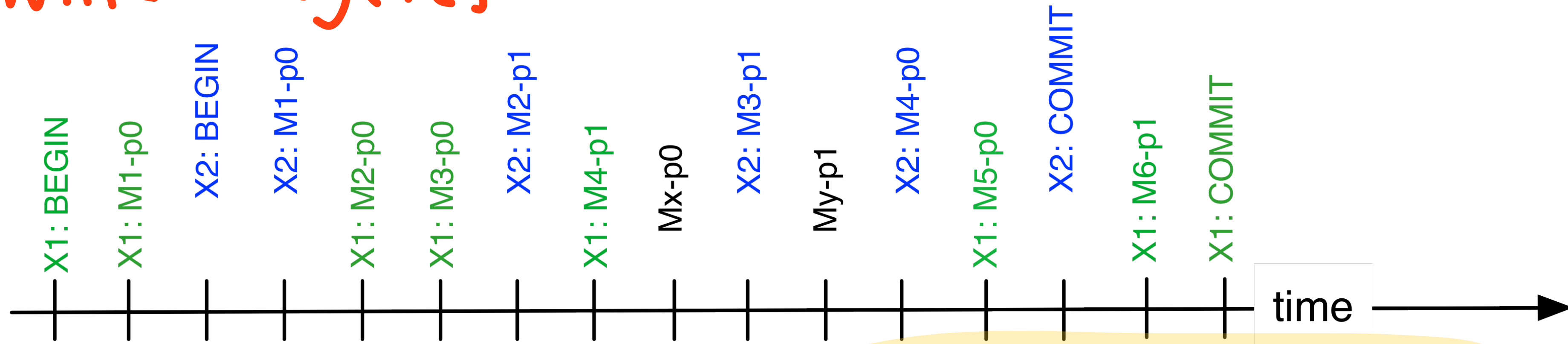
Commit order: $X2 < X0$

Since X2 is committed first, each partition will expose messages from X2 before X1.

Consumer processing order

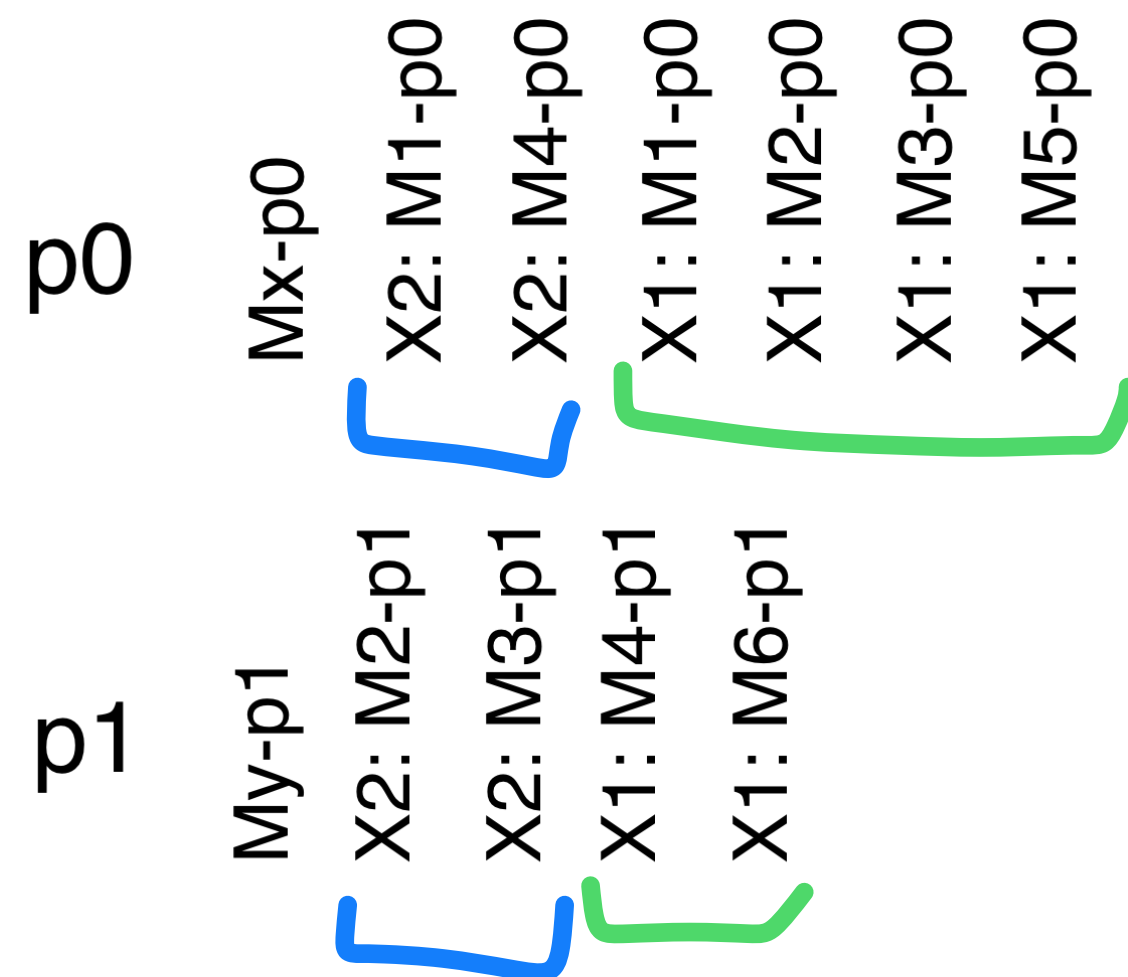
	p0	p1
	Mx-p0	My-p1
	X2: M1-p0	X2: M2-p1
	X2: M4-p0	X2: M3-p1
	X1: M1-p0	X1: M4-p1
	X1: M2-p0	X1: M6-p1
	X1: M3-p0	
	X1: M5-p0	

#8 Write Cycles



Commit order: $X2 < X0$

Consumer processing order



Since X2 is committed first, each partition will expose messages from X2 before X1.

#8 Write Cycles

Liiiiieees

Since X2 is committed first
each partition will expose
messages from X2 before X1

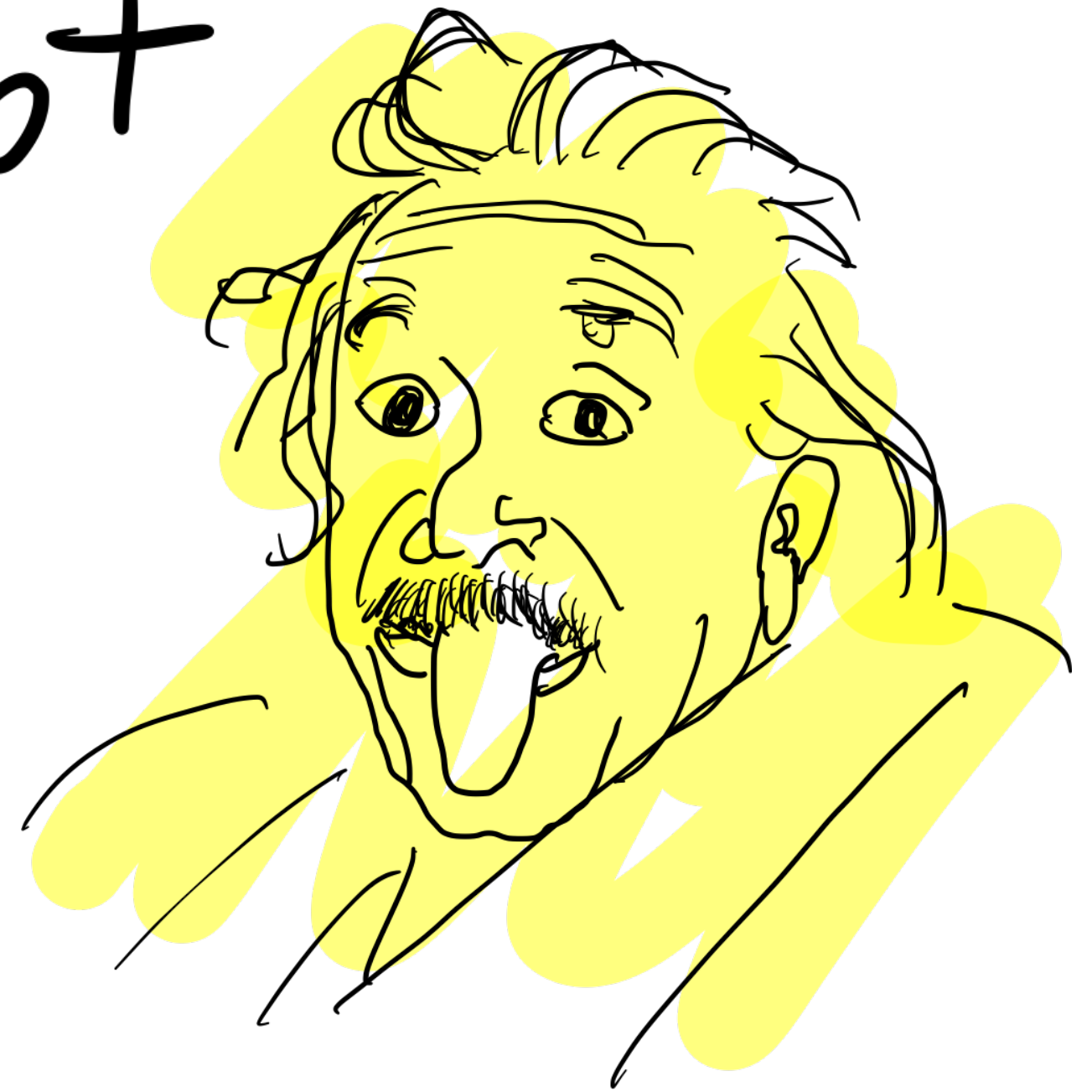
"Transactional Messaging

in Kafka"

we have no strong evidence that applications can benefit from the commit order option, we opted for not implementing it

"Exactly Once Delivery and Transactional Messaging in Kafka"

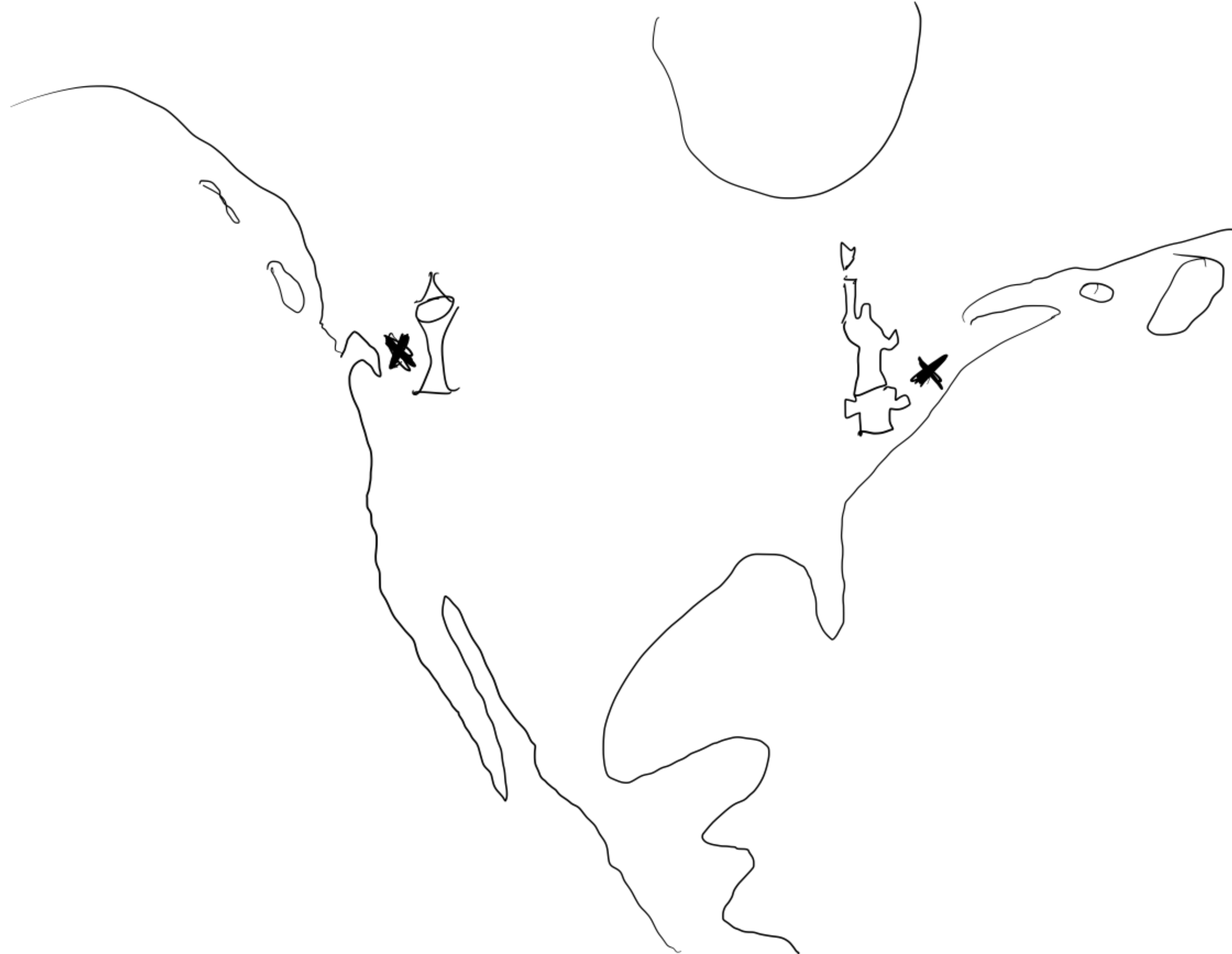
In theory, theory and
practice are the same. In
practice, they are not



Kafka protocol

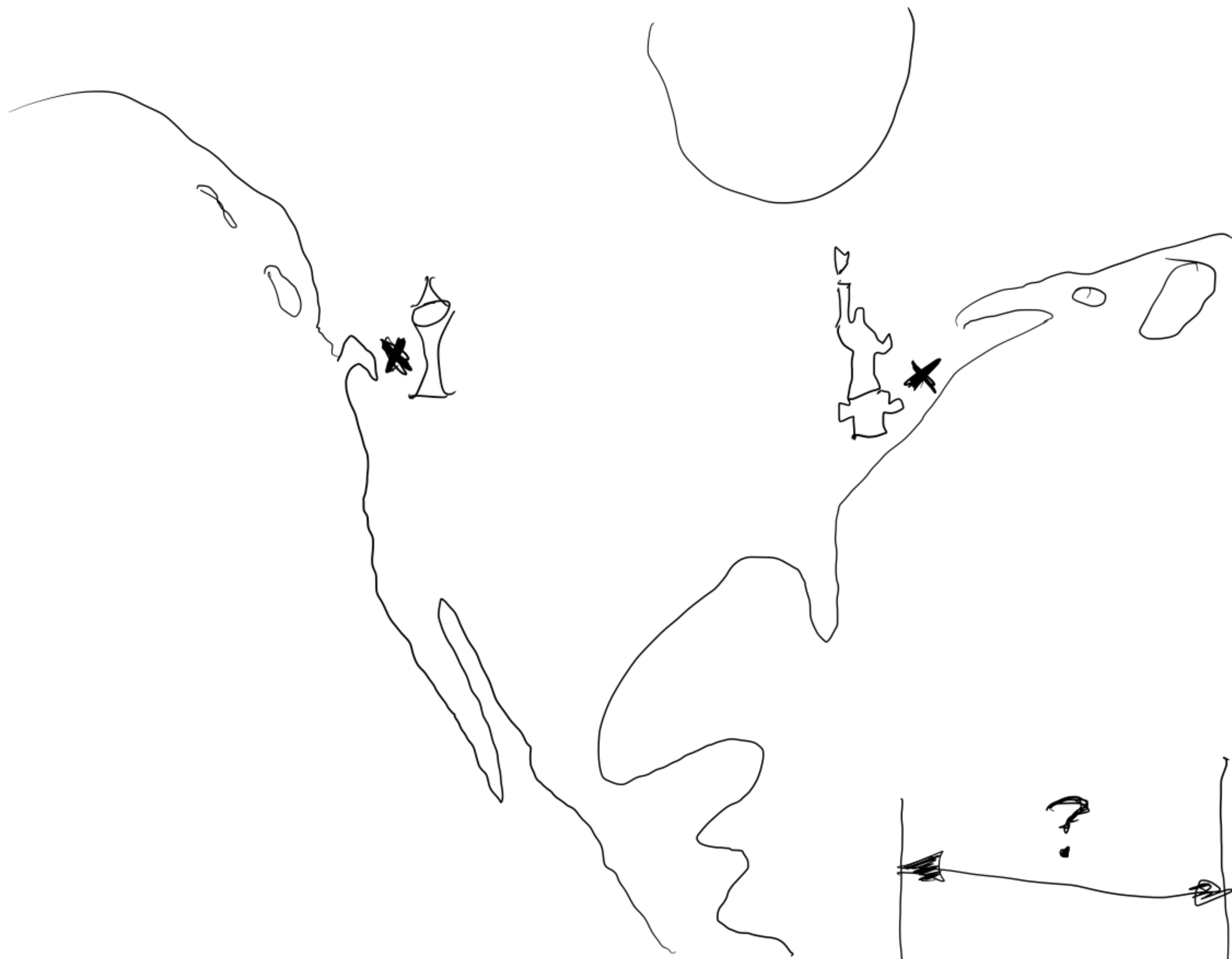
permits GO

(tx interleaving)

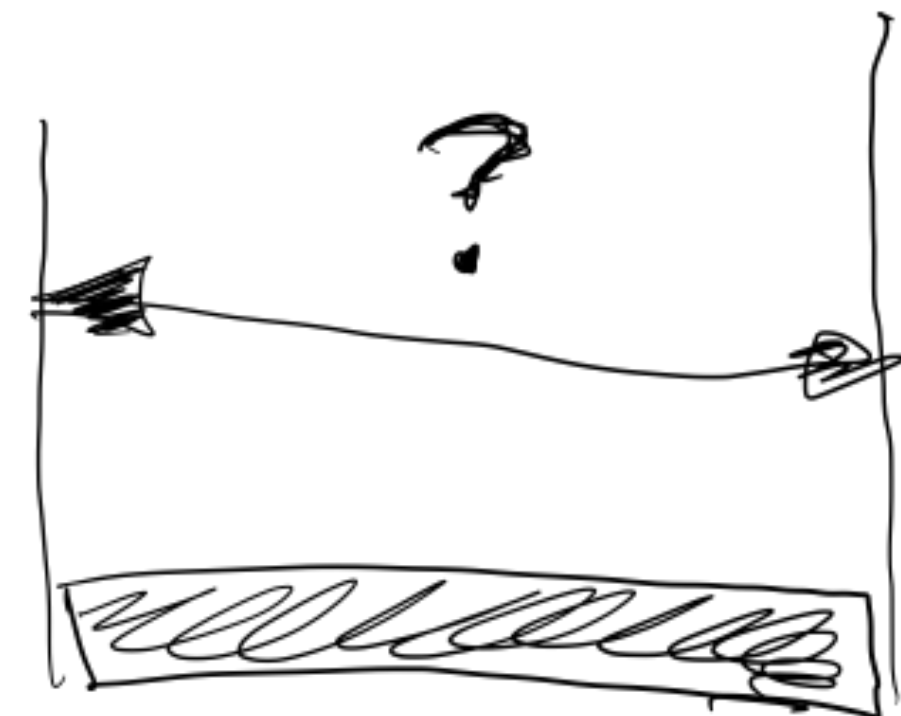


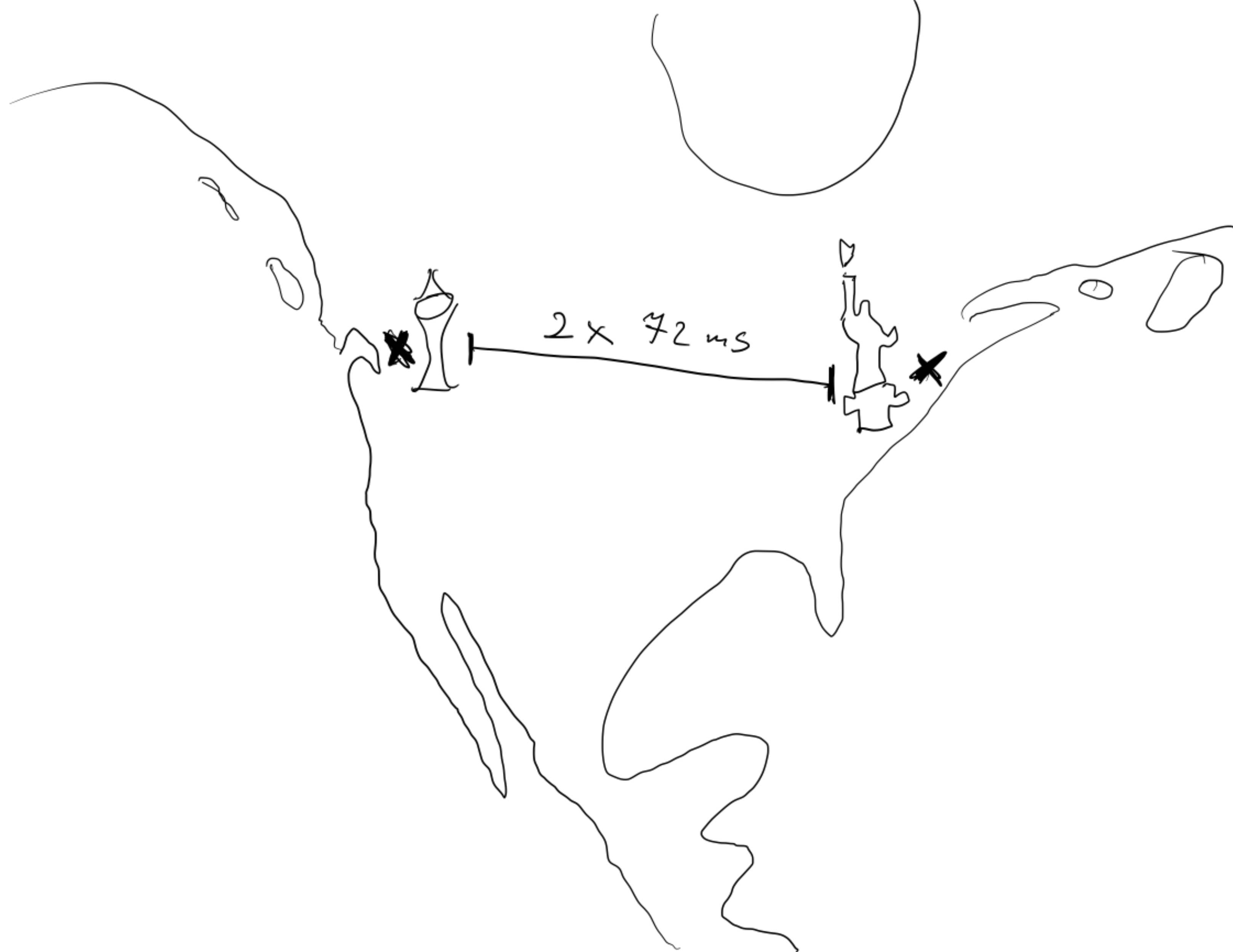
Seattle
NYC





Seattle
NYC





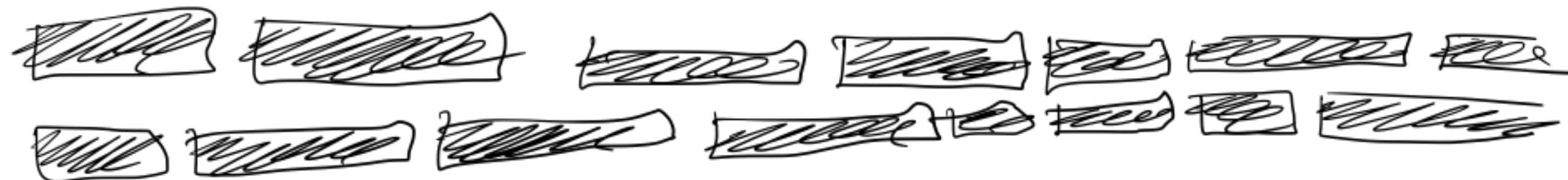
Seattle
NYC





Seattle

NYC



Non-Monotonic Snapshot Isolation: scalable and strong consistency for geo-replicated transactional systems

Masoud Saeida Ardekani

Université Pierre-et-Marie Curie (UPMC-LIP6)

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University of Neuchâtel

Neuchâtel, Switzerland

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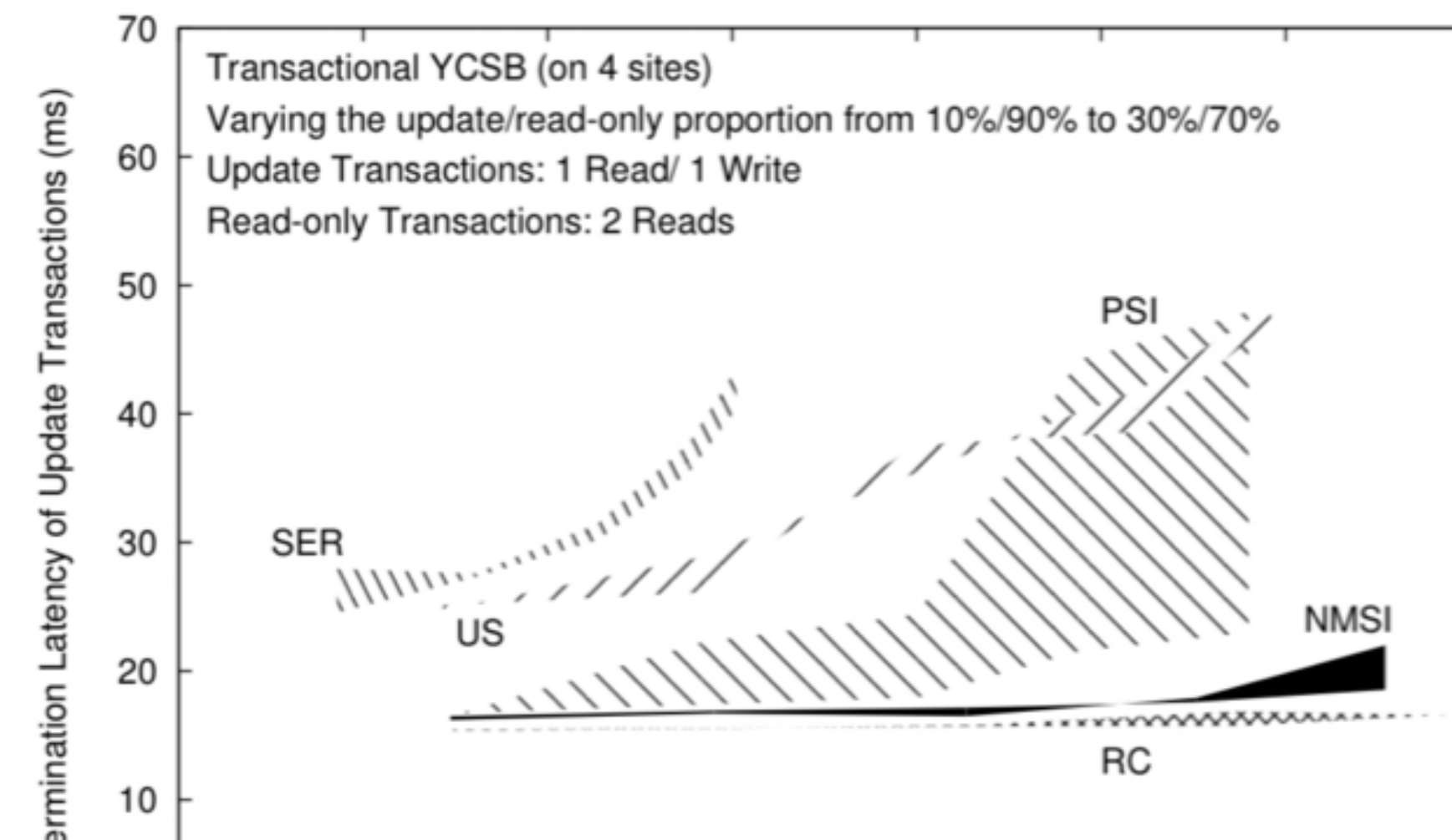
Marc Shapiro

INRIA & UPMC-LIP6

Paris, France

<http://lip6.fr/Marc.Shapiro/>

Abstract—Modern cloud systems are geo-replicated to improve application latency and availability. Transactional consistency is essential for application developers; however, the corresponding concurrency control and commitment protocols are costly in a geo-replicated setting. To minimize this cost, we identify the following essential scalability properties: (i) only replicas updated by a transaction T make steps to execute T ; (ii) a read-only transaction never waits for concurrent transactions and always commits; (iii) a transaction may read object versions committed after it started; and (iv) two transactions synchronize with each other only if their writes conflict. We present Non-Monotonic Snapshot Isolation (NMSI), the first strong consistency criterion to allow implementations with all four properties. We also present a practical implementation of NMSI called Jessy, which we



- 1) all-or-nothing for multi-partitions writes
- 2) atomic consume-transform-produce loop
- 3) ~~consistent snapshots~~

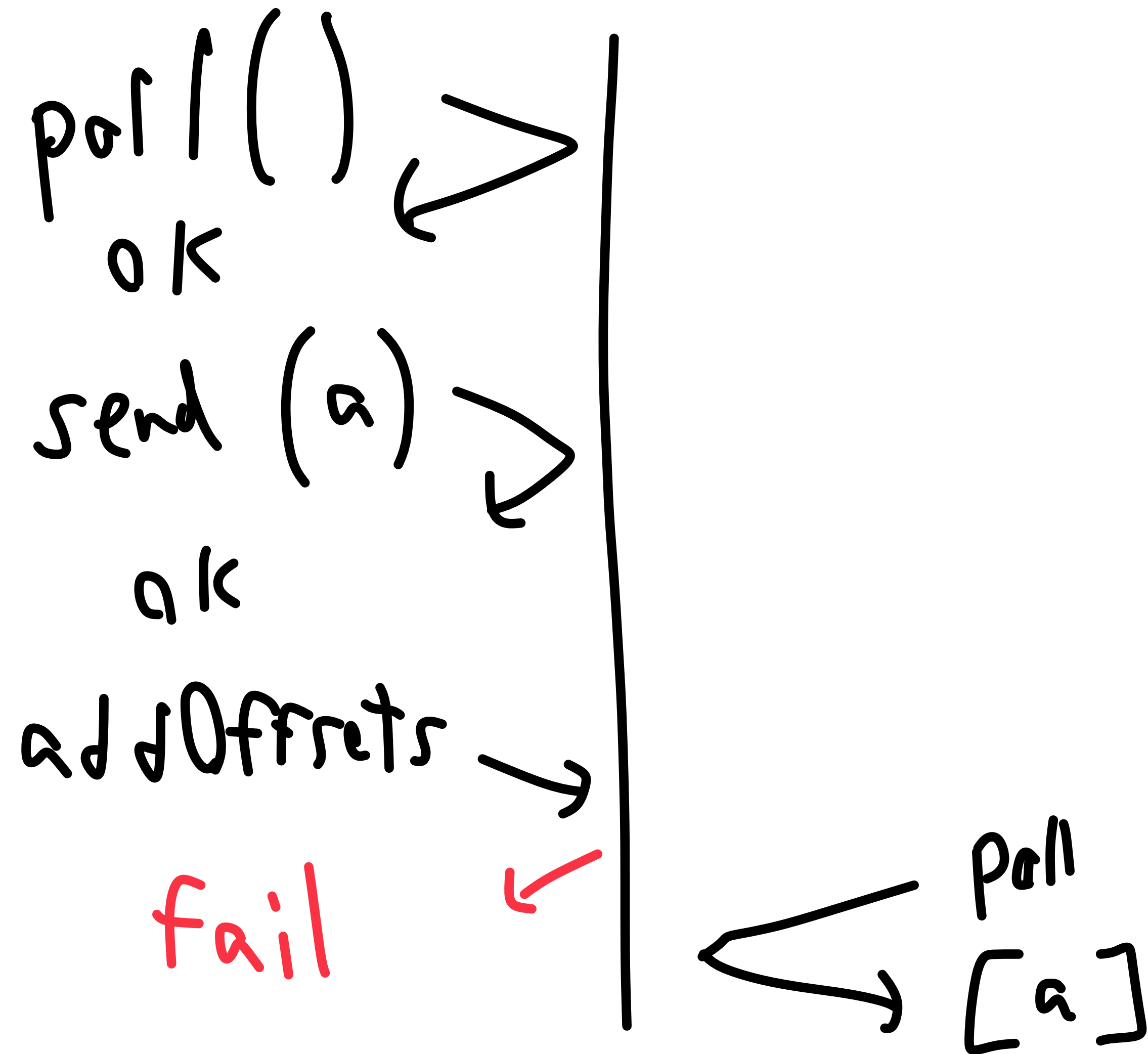
#8 Write Cycles

This is just how Kafka/

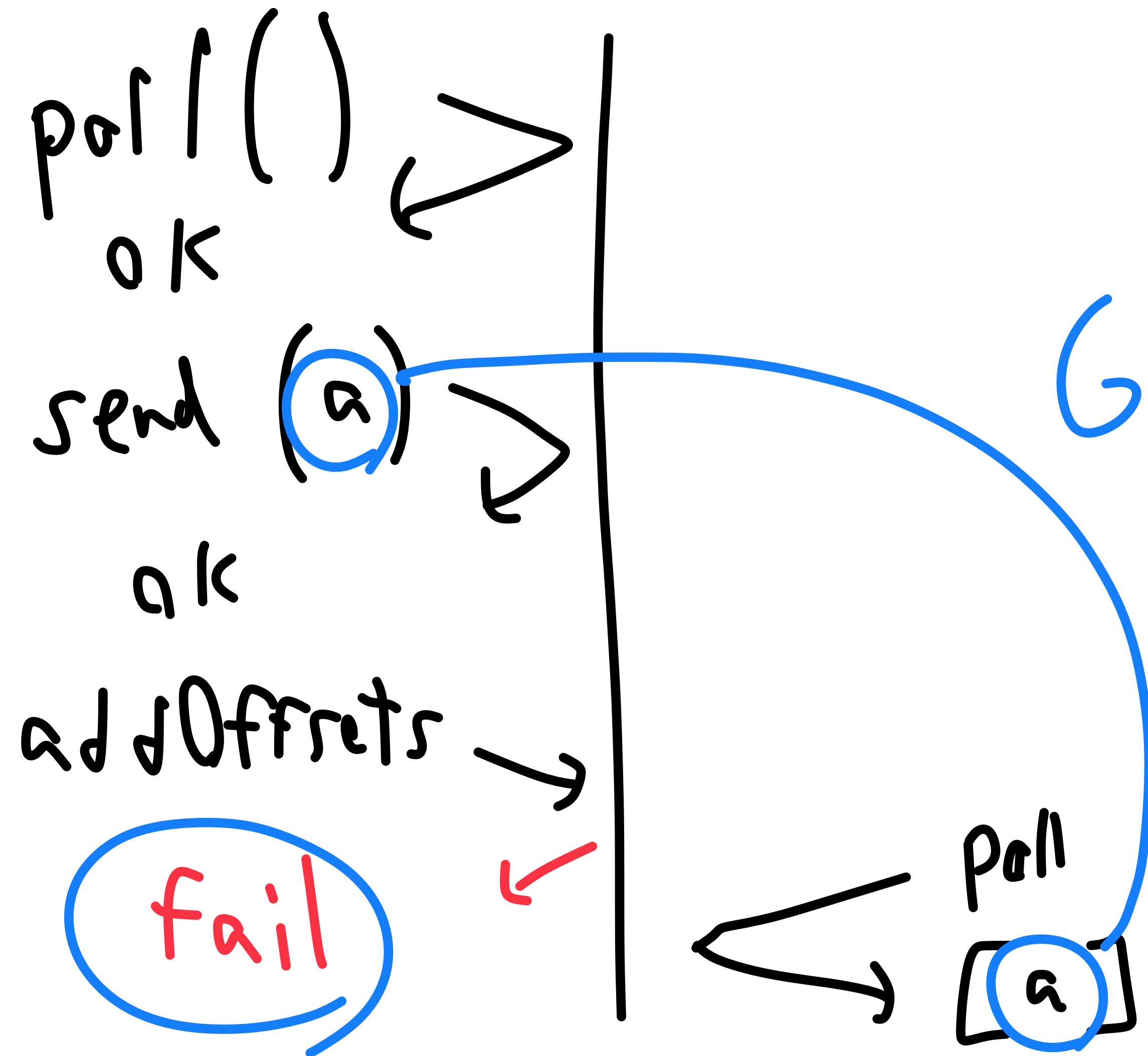
Redpanda txns work

→ Needs documentation

#3036 Aborted Reads & Circular Information Flow



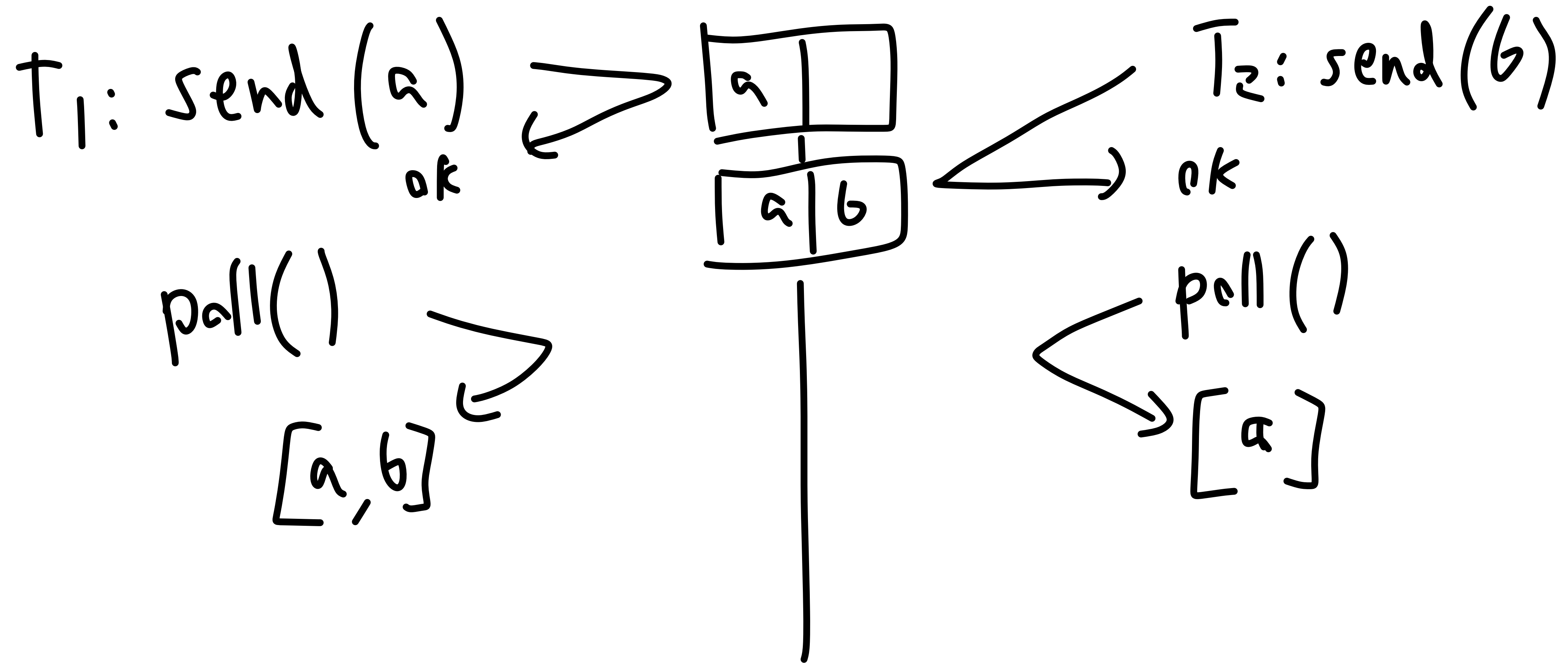
#3036 Aborted Reads & Circular Information Flow



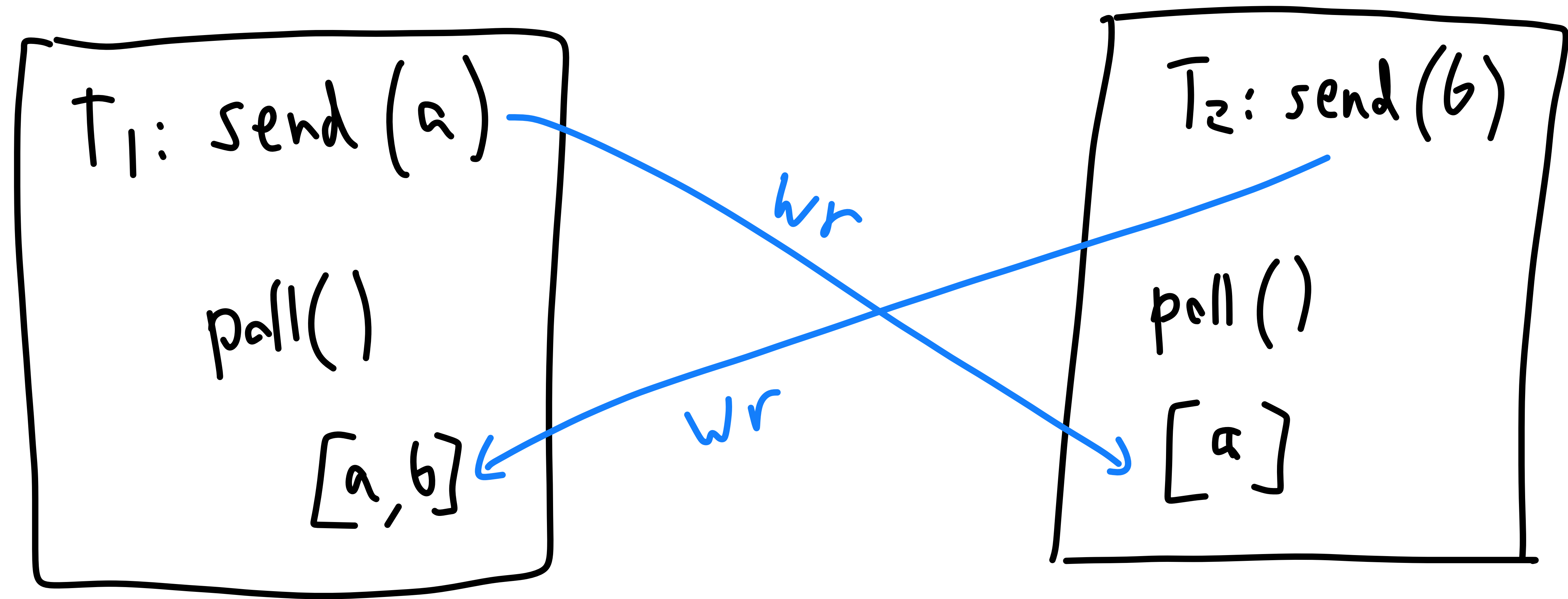
G1a: Aborted Read

Happened often in
healthy clusters!

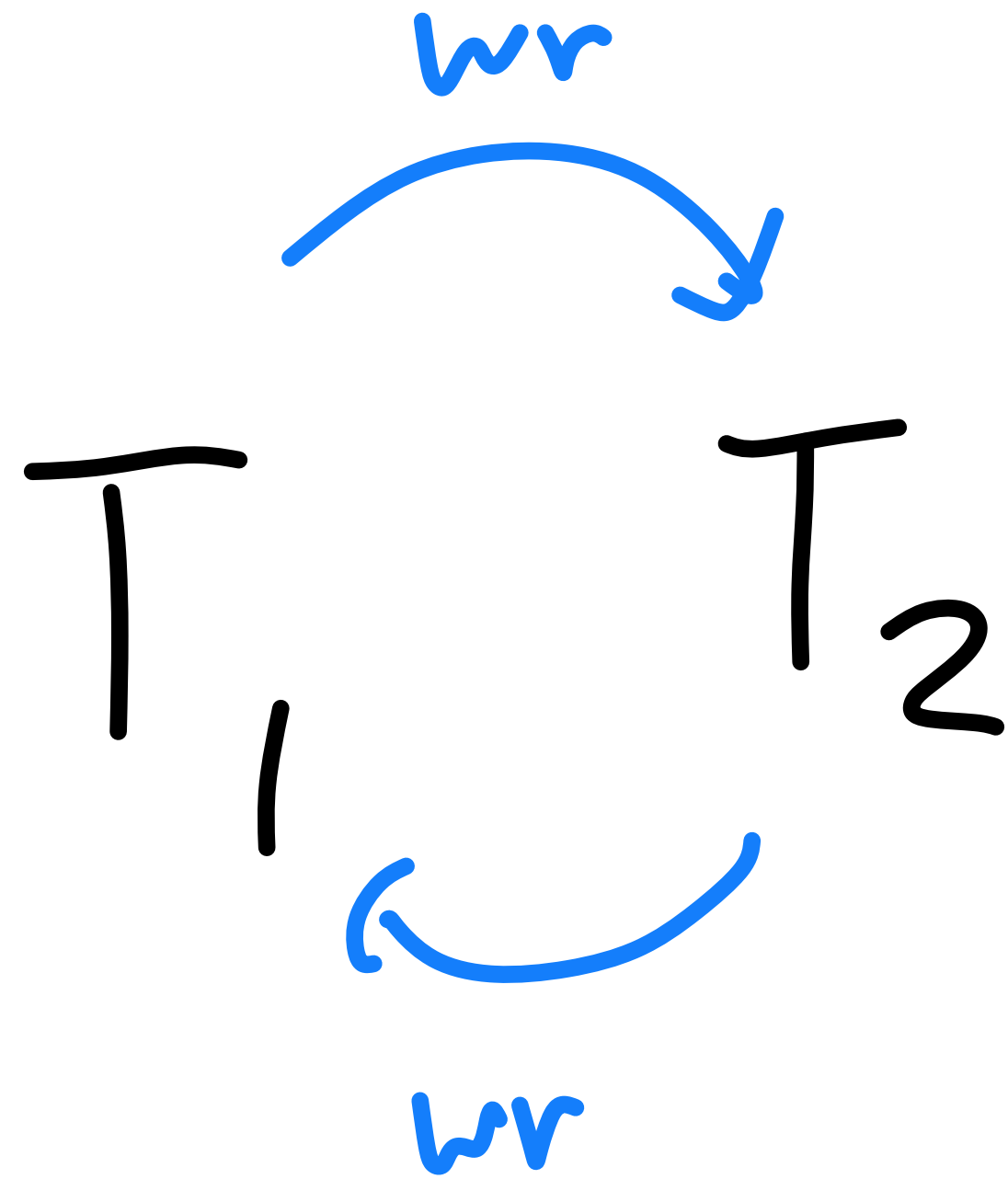
#3036 Aborted Reads & Circular Information Flow



#3036 Aborted Reads & Circular Information Flow



#3036 Aborted Reads & Circular Information Flow



Gluc: Circular
Information
Flow

Illegal in Adya Read Committed!

#3036 Aborted Reads & Circular Information Flow

- Off-by-one error allowed LSO to advance beyond committed offsets
 - #3232 accidentally aborted more txns than necessary
- Fixed in 21.10.2

#10 Internal Non-monotonic Polls

beginTxn()

poll()

poll()

addOffsetsToTxn(...)

commitTxn()

#10 Internal Non-monotonic Polls

beginTxn()

poll() → [a, b, c]

poll() → [b, c, d]

addOffsetsToTxn(...)

commitTxn()

#10 Internal Non-monotonic Polls

- Caused by consumer rebalance
- Clients must register rebalance listeners to detect potential non-monotonicity in transactions

#3616-a Aborted Read w/ Invalid Txn State Exception

beginTxn()

send(a)

commitTxn() → Invalid Txn State Exception

#3616-a Aborted Read w/ Invalid Txn State Exception

beginTxn()

send(a)

commitTxn() → Invalid Txn State Exception

⋮

poll() → [a]

#3616-a Aborted Read w/ Invalid Txn State Exception

Pause, crash, or partition

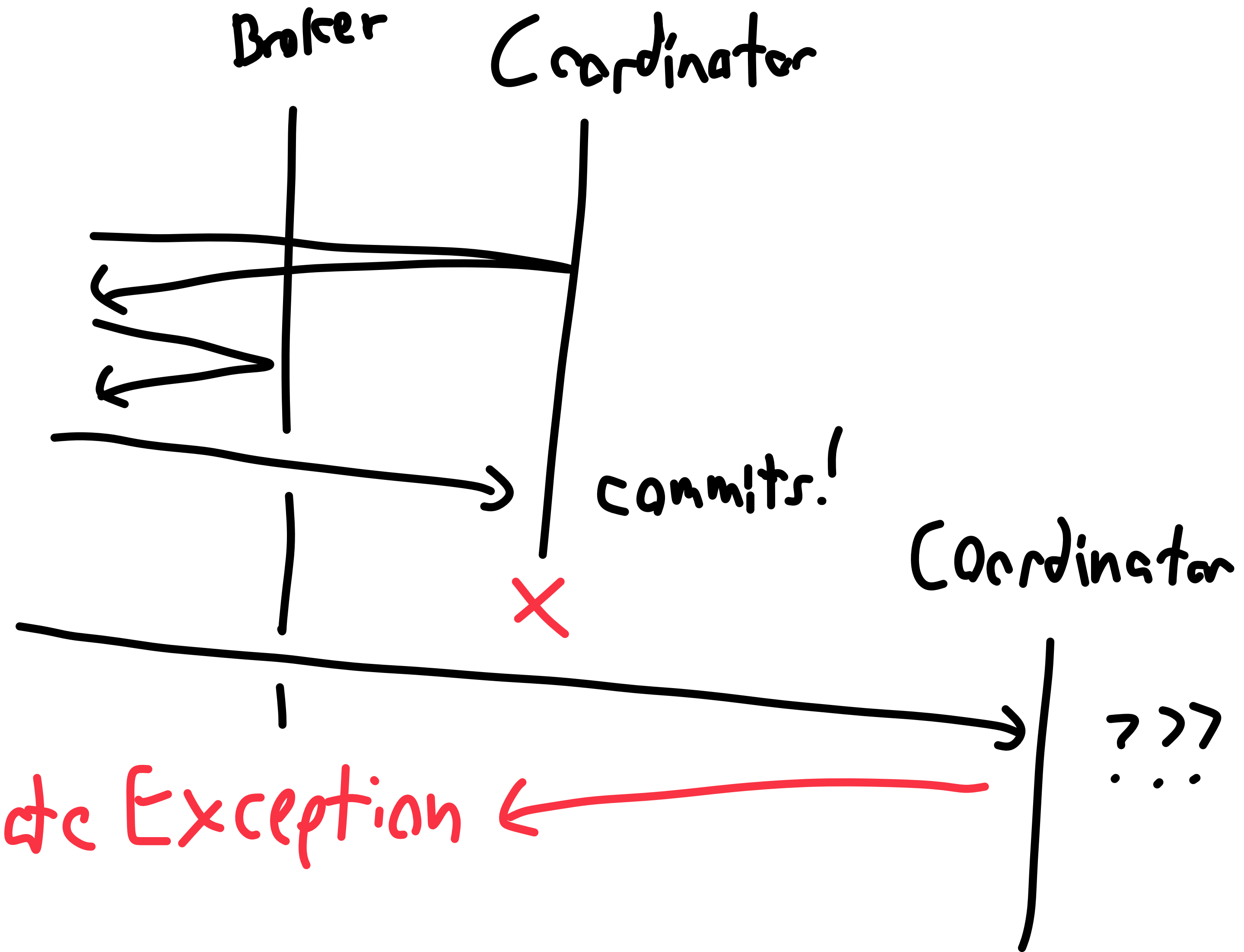
led to GLs

#3616-a Aborted Read w/ Invalid Txn State Exception

begin Txn()

send (a)

commit Txn()
(retry)



Invalid Txn State Exception

#3616-a Aborted Read w/ Invalid Txn State Exception

Fixed by returning UnknownServerError

January 21, 2022

↳ Released in 21.11.15

#3616-6 Last Transactional Writes

beginTxn()

send(a) → offset 2

commitTxn() → ok

#3616-6 Last Transactional Writes

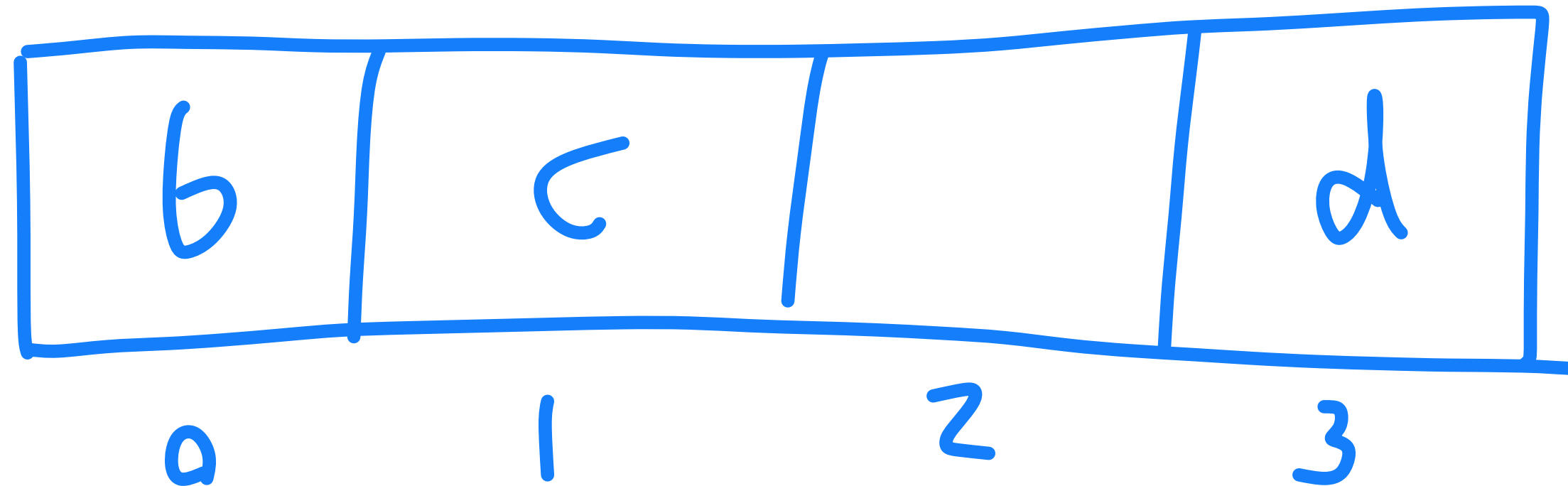
beginTxn()

send(a) → offset 2

commitTxn() → ok

poll()

→



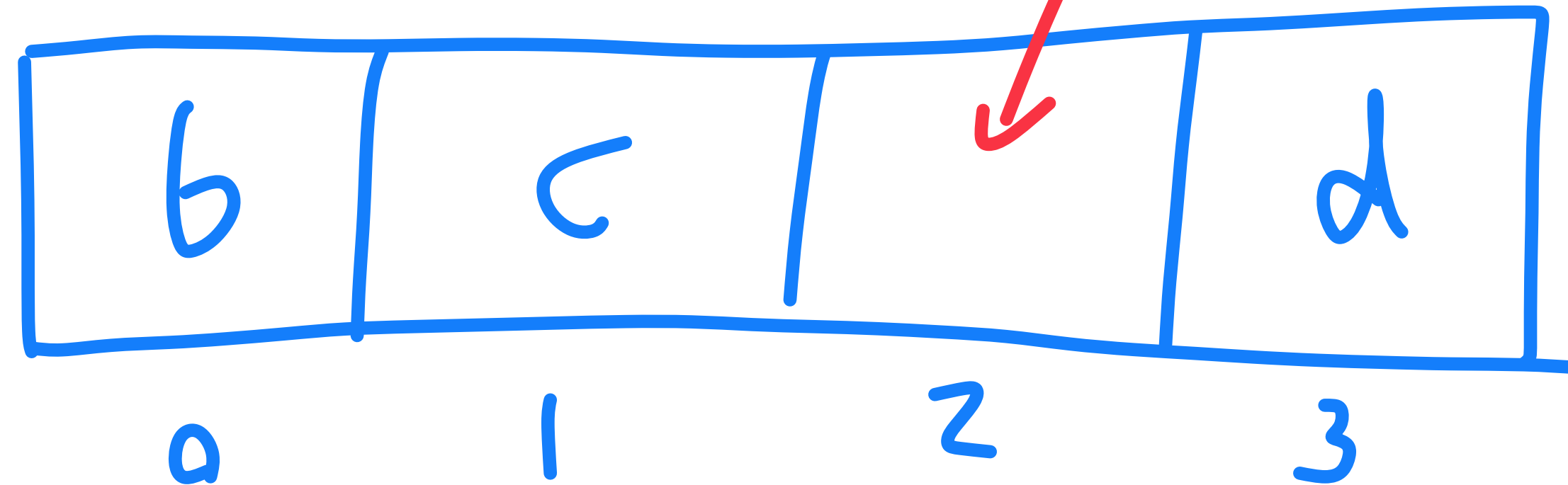
#3616-6 Last Transactional Writes

beginTxn()

send(a) → offset 2

commitTxn() → ok

poll() →



#3616-6 Last Transactional Writes

Happened in healthy clusters

on 21.11.2

#3616-6 Last Transactional Writes

```
lock.acquire()  
if (!log.sync()) {  
    return;  
}
```

Critical
section

```
[ do_x();  
  do_y();  
  lock.release()
```

#3616-6 Last Transactional Writes

```
lock.acquire()  
if (!log.sync()) {  
    return;  
}
```

Critical
section

```
[ do_x();  
  do_y();  
  lock.release()
```

Made loses, then regains
leadership — other leader did
things in our critical section!

#3616-6 Last Transactional Writes

```
lock.acquire()
```

```
if (!log.sync()) {
```

```
    return;
```

```
}
```

```
term = _insync_term
```

```
do_x(term);
```

```
do_y(term);
```

```
lock.release()
```

Critical
section

#3616-6 Last Transactional Writes

Fixed January 21, 2022

↳ Released in 21.11.15

Discussion

#	Issue	Fault	Resolved
1	Duplicate writer by default	P ∨ C ∨ Part	22.1.1
3039	Duplicate writer w/ idempotence	P ∨ C ∨ Part	21.10.3
3335	Assert crash deallocating partitions	M	
3336	Assert crash w/ partition IDs	C	22.1.1
3003	Inconsistent offsets	C ∨ Part	21.10.3
6	Last/Stale messages	P ∨ C ∨ Part	
7	Aborted read — NotLeader or Follower	M ∧ P	
8	Write cycles	∅	
3036	Aborted read/Circular information flow	∅	21.10.2
10	Internal non-monotonic polls	∅	
3616-a	Aborted read — Invalid Txn state	P ∨ C	21.11.15
3616-b	Last txn writer	∅	21.11.15

#	Issue	Fault	Resolved
1	Duplicate writer by default	P v C v Part	22.1.1
3039	Duplicate writer w/ '' '	P.. C v Part	21.10.3
3335	Asrent crash d		
3336	Asrent crash		22.1.1
3003	Inconsistent ori.	, ant	21.10.3
6	Last/Stale messageres	P v C v Part	
7	Aborted read - Not Leader or Follower	M ^ P	
8	Write cycles	x	
3036	Aborted read/Circular information flow	x	21.10.2
10	Internal non-monotonic polls	x	
3616-a	Aborted read - Invalid Txn state	P v C	21.11.15
3616-b	Last txn writers	∅	21.11.15

#	Issue	Fault	Resolved
1	Duplicate writer by node H	P v C v Part	22.1.1
3039	Duplicate write	r v Part	21.10.3
3335	Assert crash	Effective data loss	
3336	Assert crash		22.1.1
3003	Inconsistent offset.	r v Part	21.10.3
6	Last/stale messages	P v C v Part	
7	Aborted read - Not Leader or Follower	M \wedge P	
8	Write cycles	\emptyset	
3036	Aborted read/Circular information flow	\emptyset	21.10.2
10	Internal non-monotonic calls	\emptyset	
3616-a	Aborted read - Invalid Txn state	P v C	21.11.15
3616-b	Last txn writer	\emptyset	21.11.15

Redpanda found many
of these bugs independently

#3039 #3003 #3035

Most frequent issues

resolved in 21.10.3 -

but some serious problems

in 21.11.2

No clock skew
issues

There may be more...

Ease of operation

Nat fault tolerant by default



Idempotence

T_{XNS}

Default topics

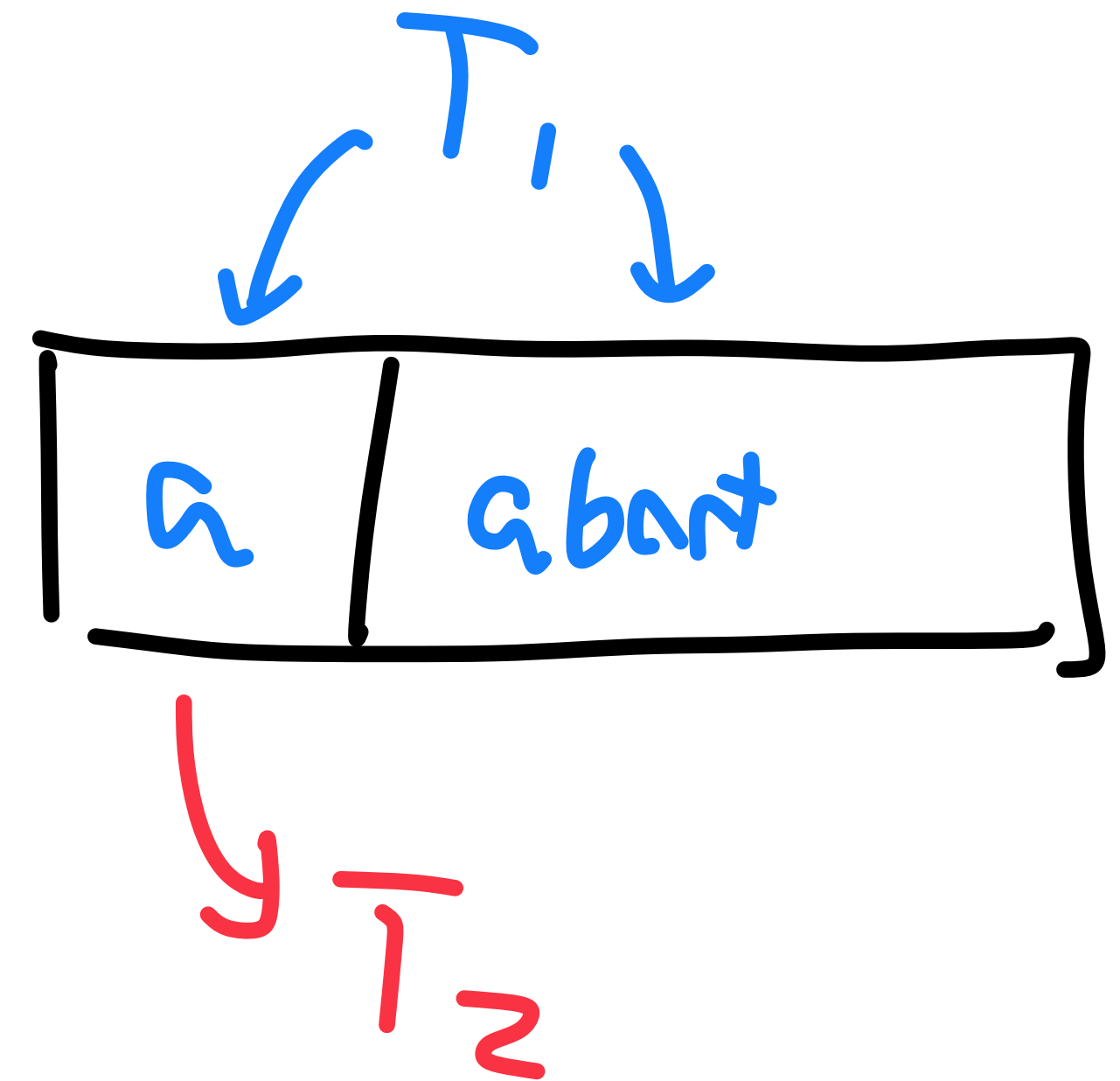
Transactions

G1a

T_1 (aborted)

↓ wr

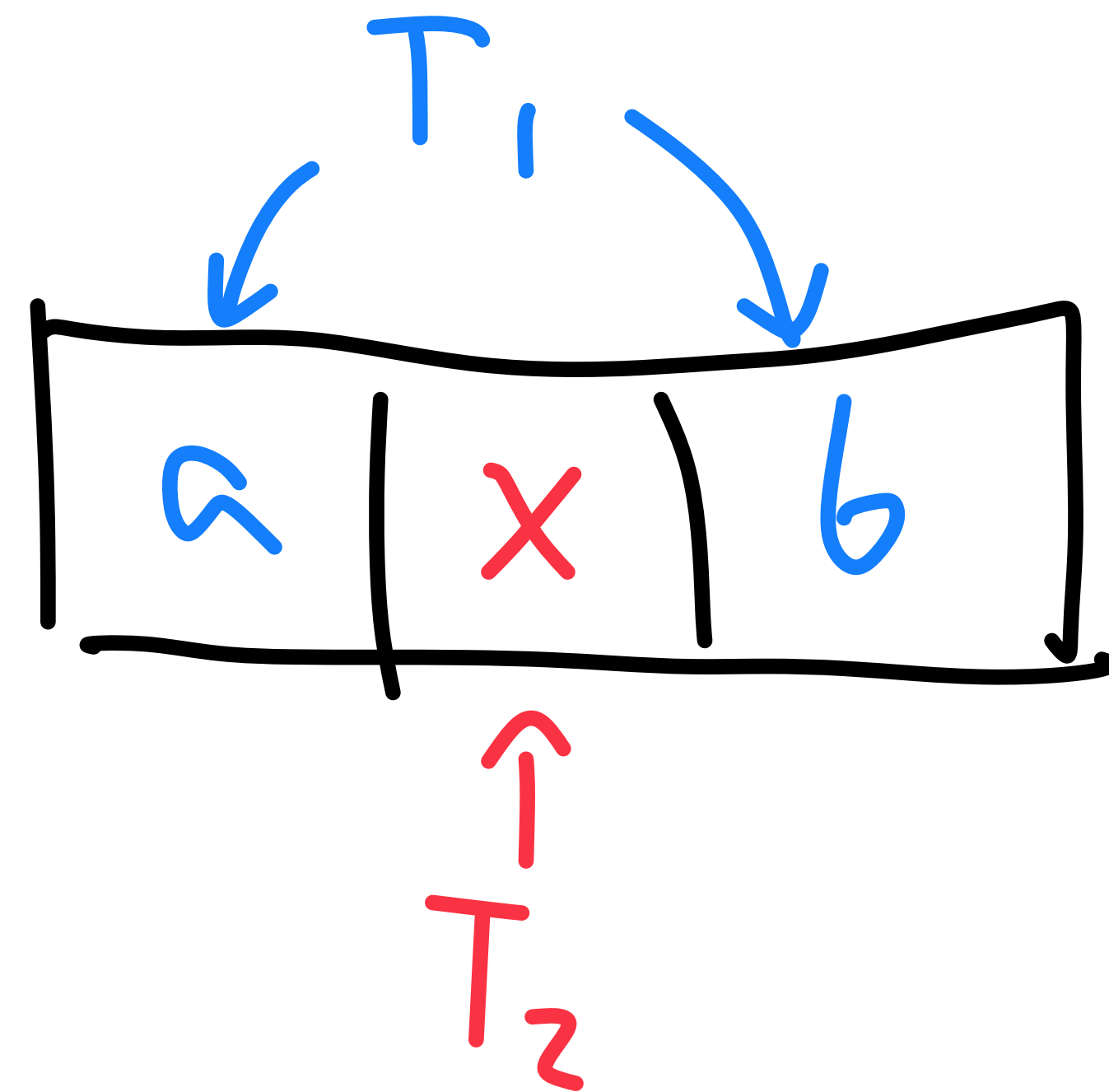
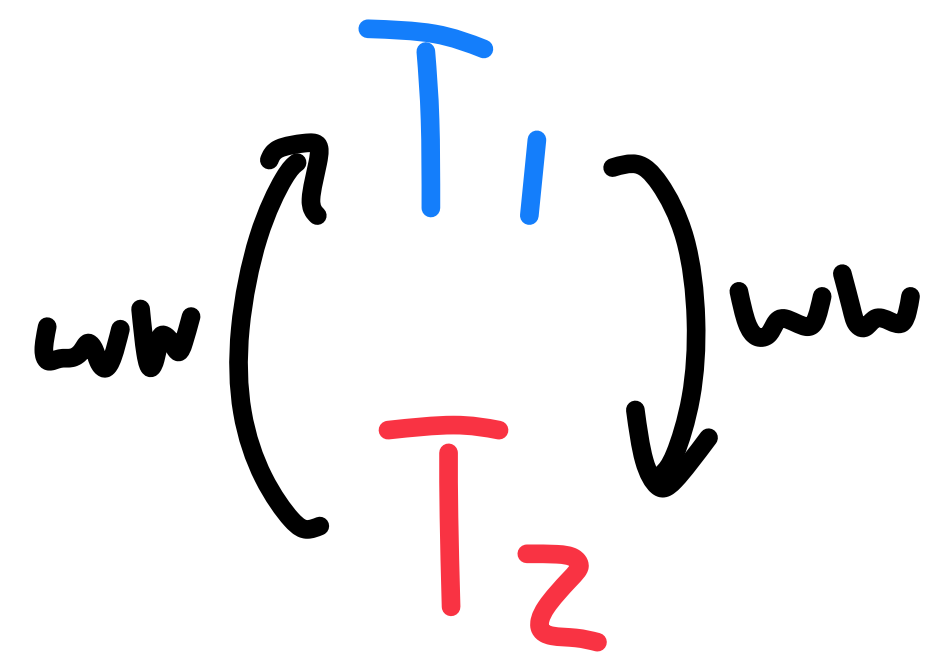
T_2



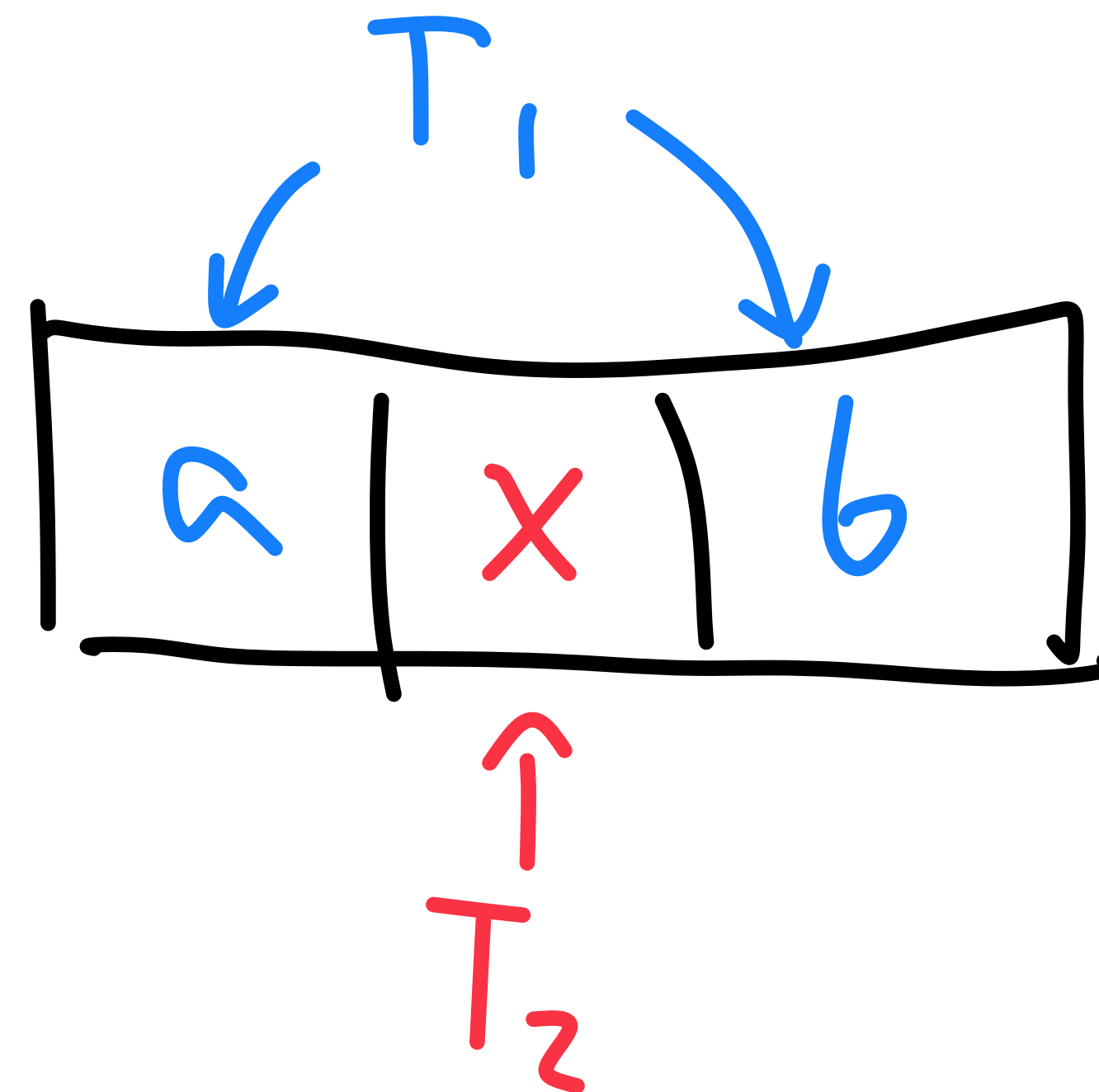
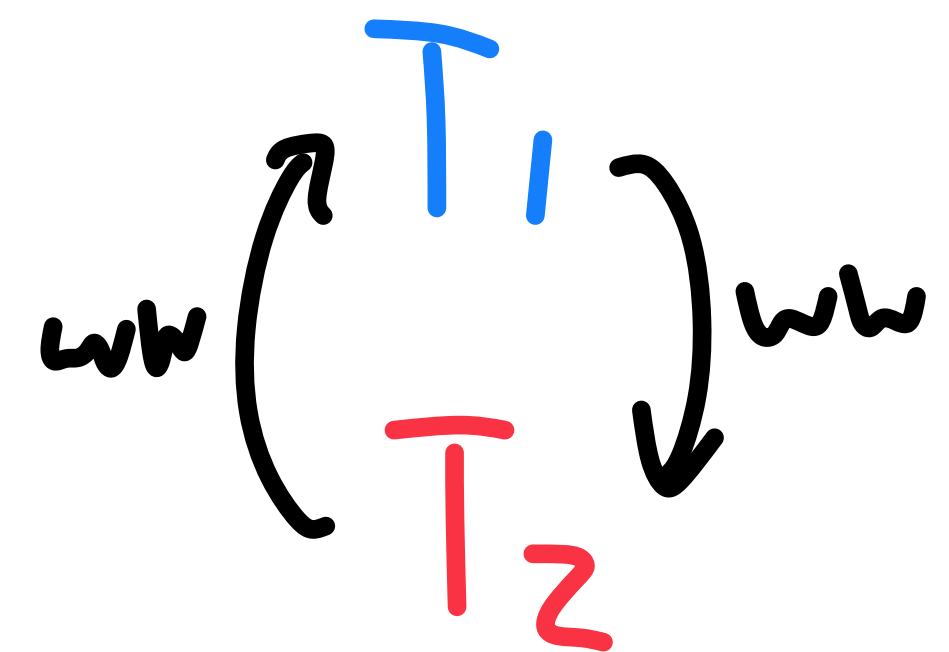
As of 21.11.15, forbidden

by read-committed

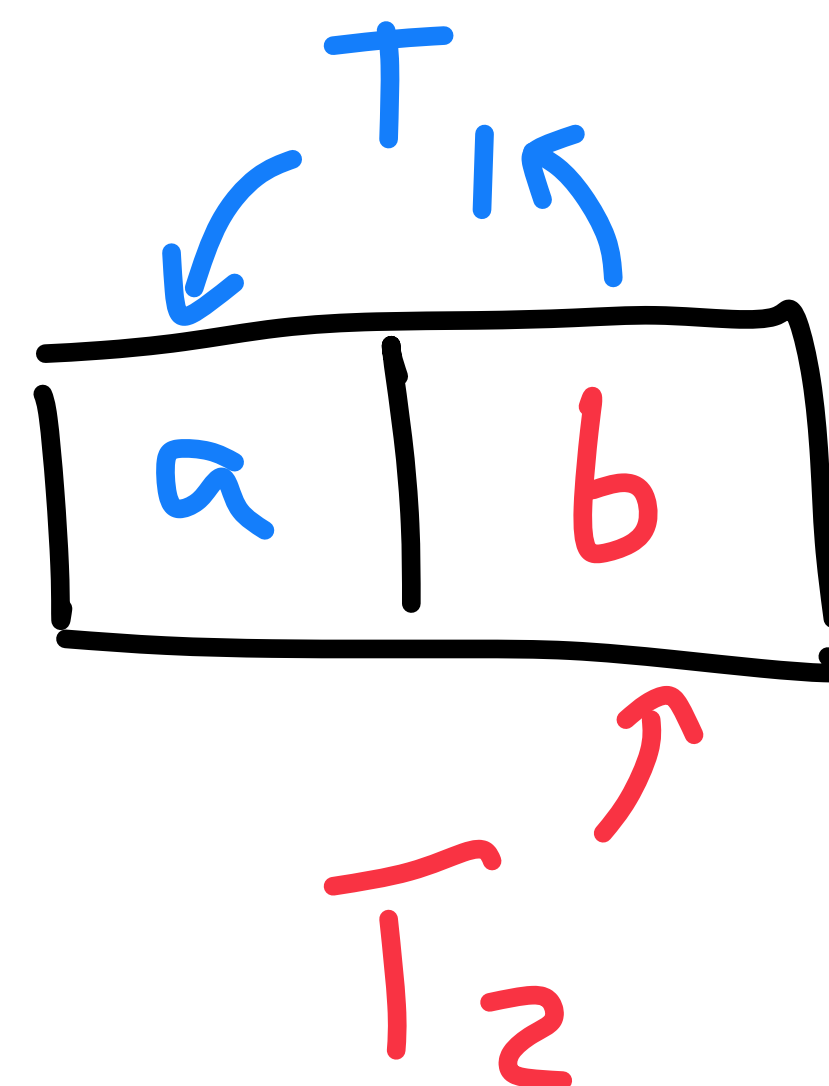
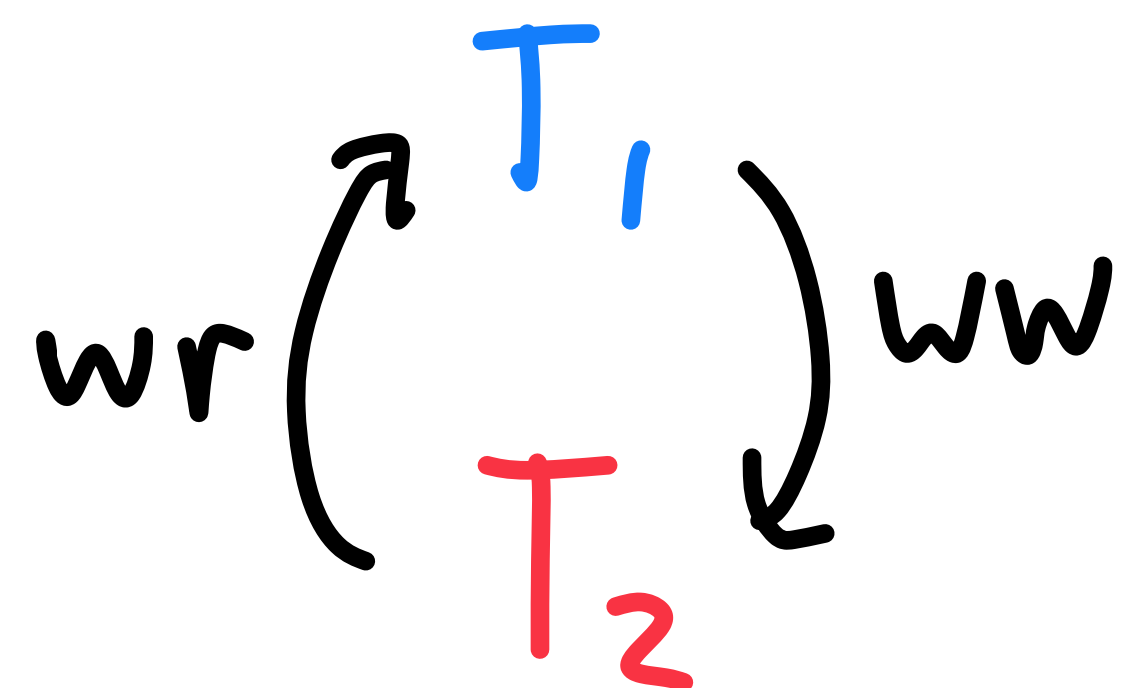
GO



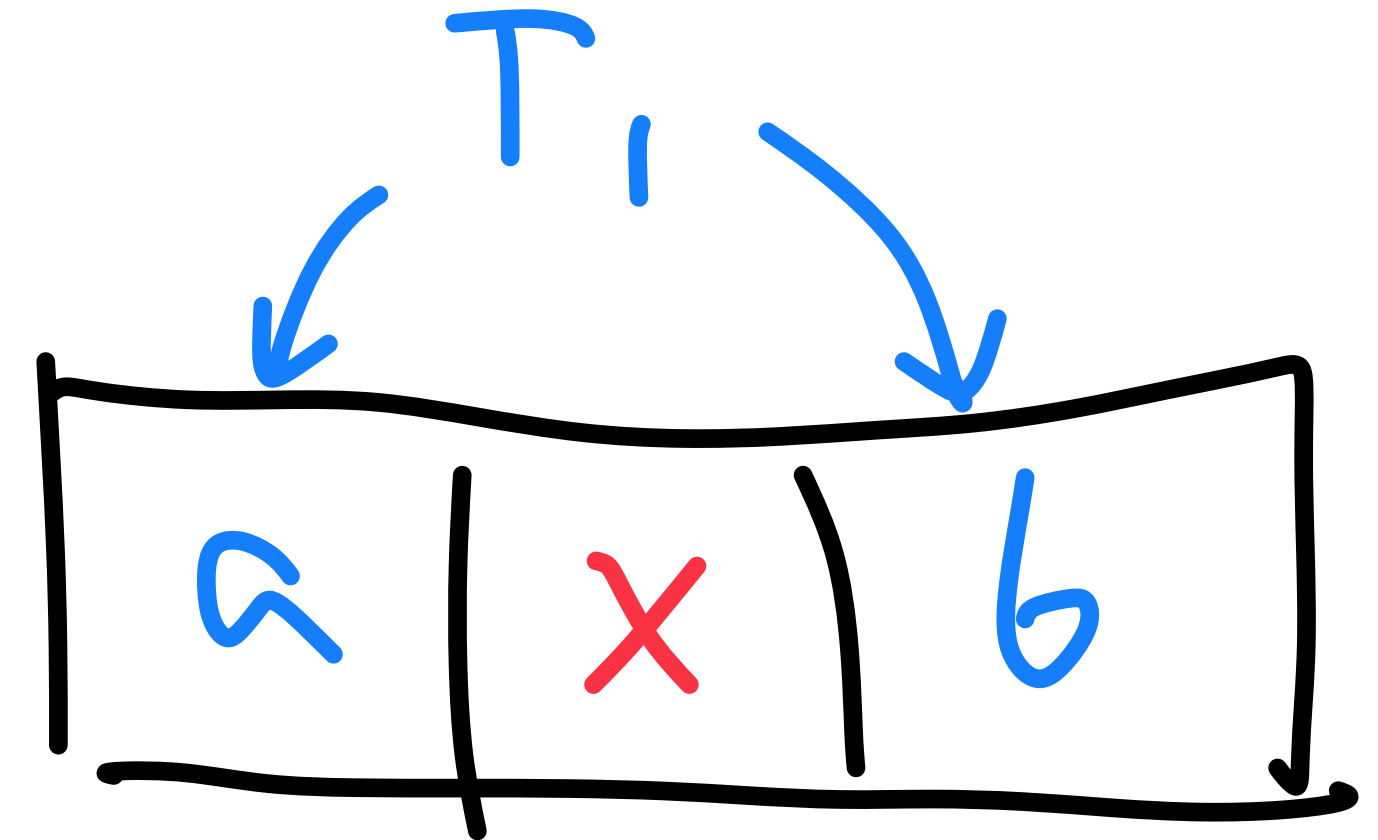
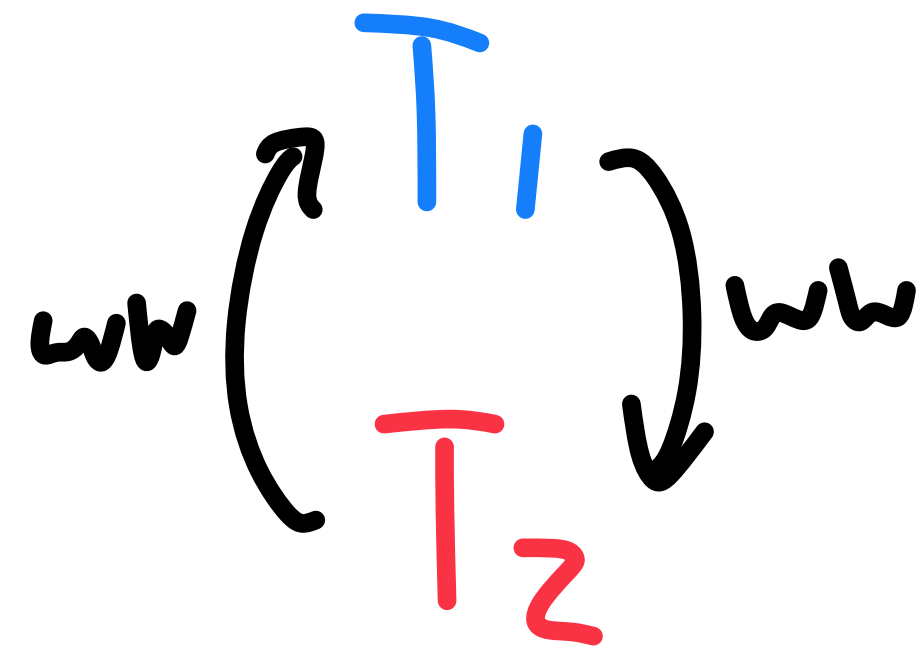
GO



G1c

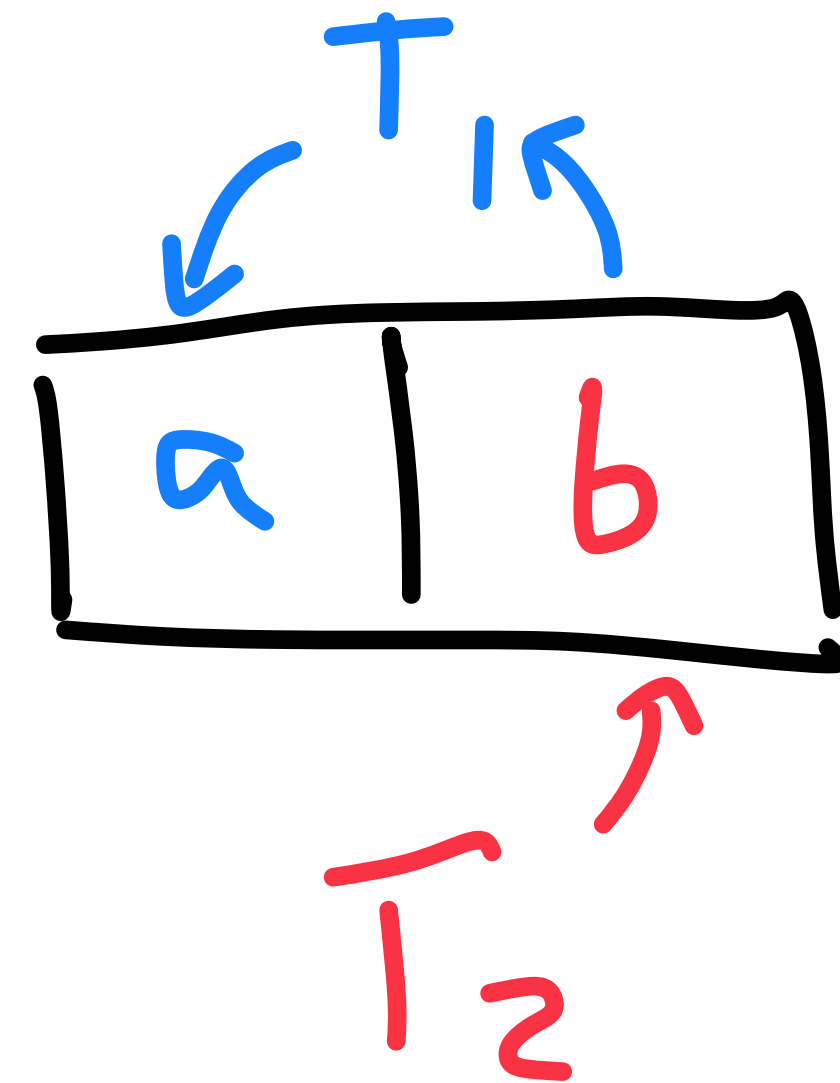
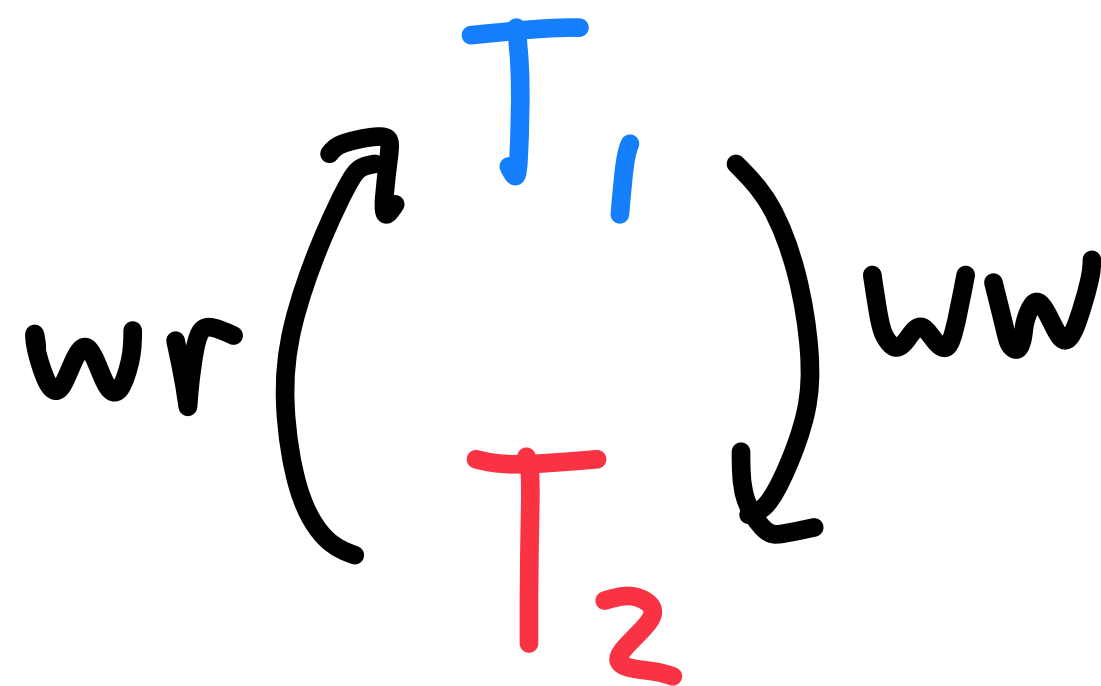


$G0$



(Ambiguous if this is OK)

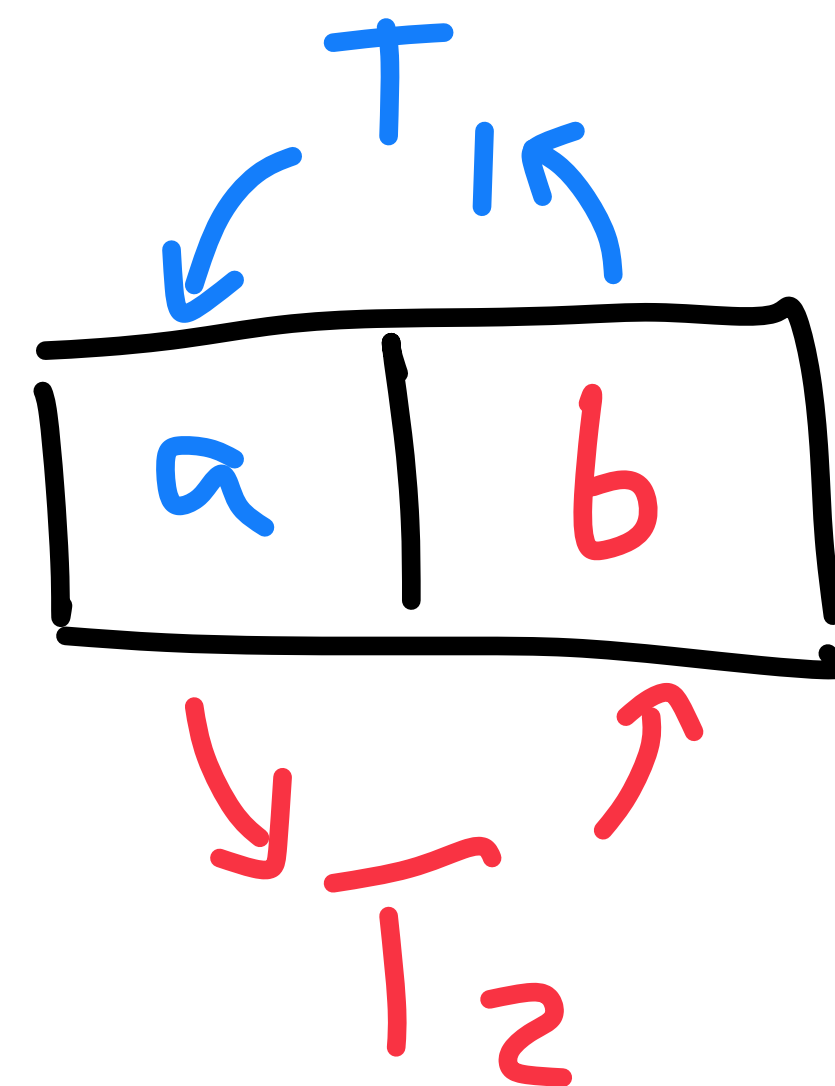
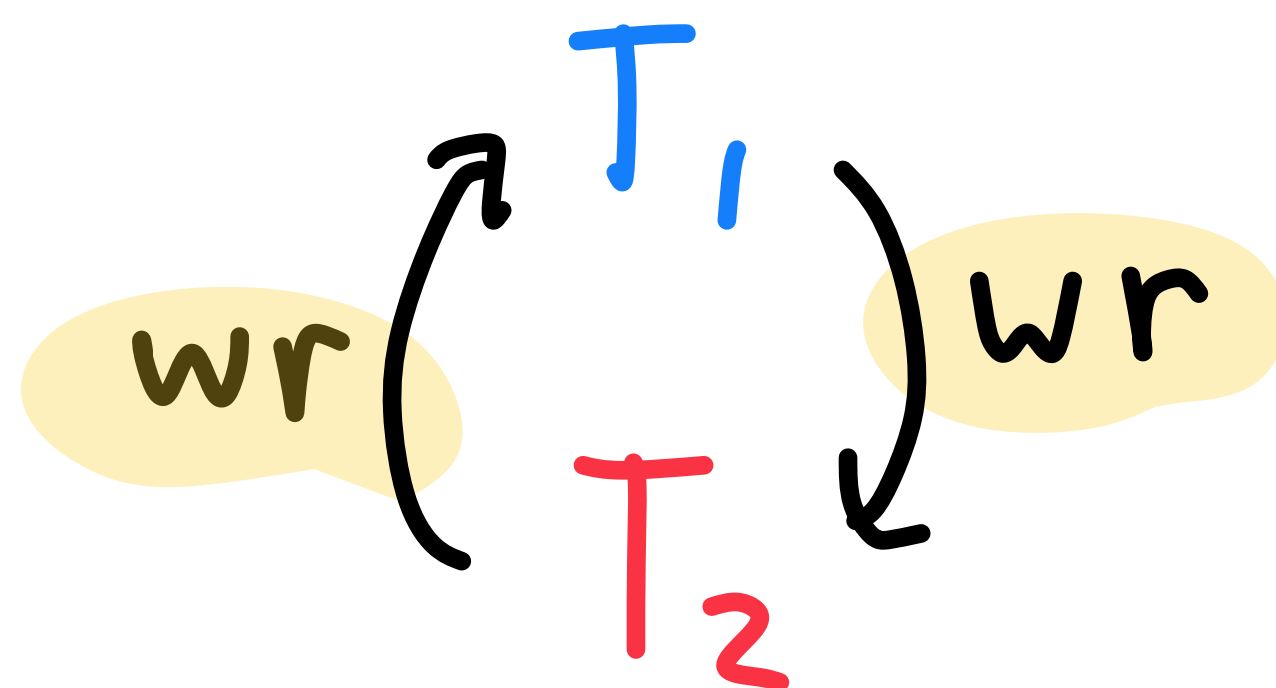
$G1c$



Should be fixed under
read-committed in 21.11.15!

G1c

w/all wr edges



Didn't check exactly-once semantics

1. Multiple polls per txn

2. Allowed more than one

transactional.id to consume a partition

All gets are off!?

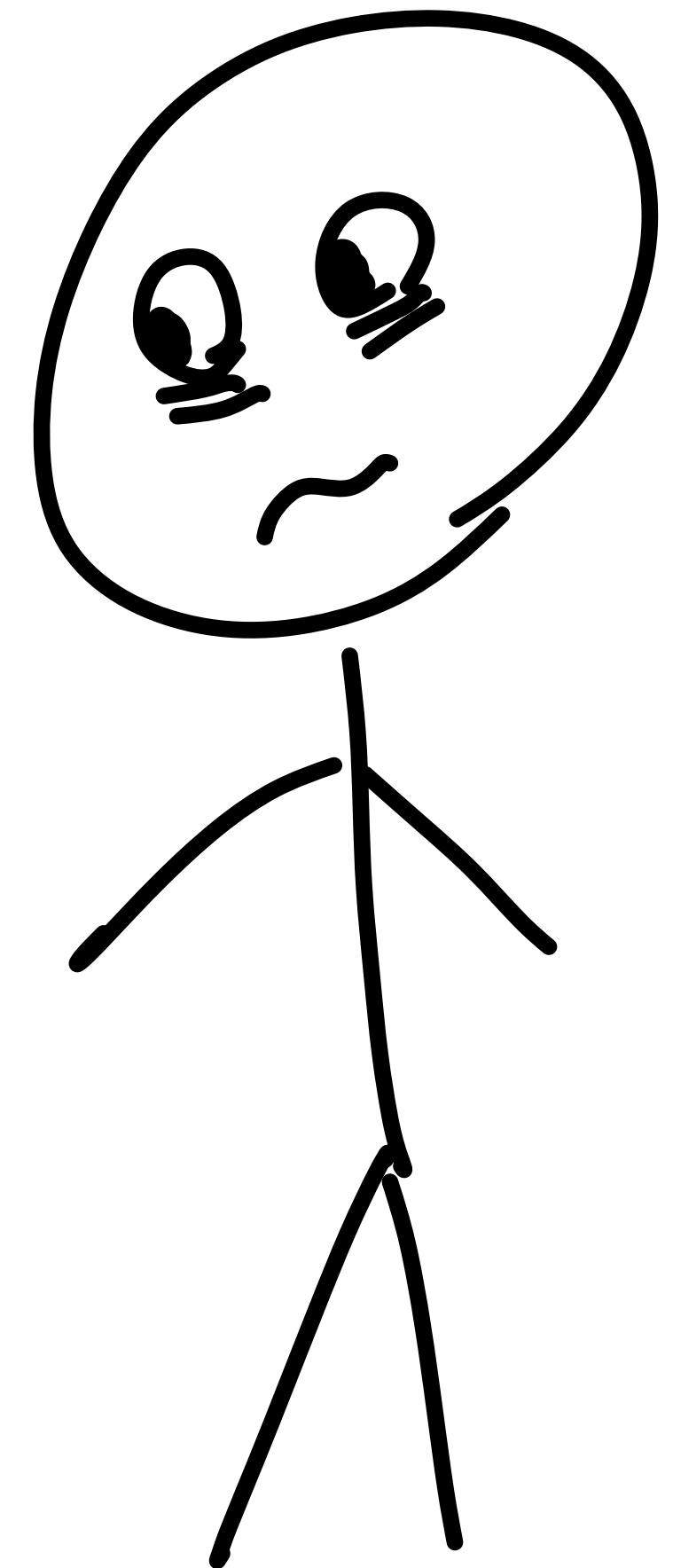
1. Multiple polls per txn

2. Allowed more than one

transactional.id to consume a partition

Txn Documentation

- Absent
- Incomplete
- Vague
- Misleading
- Wrong
- Largely explained in terms of implementation
- Go read 67 pages of design doc?



Membership

Changes

- No docs, but customers were learning how in Slack
- Improved APIs & some docs now
- Don't re-use node IPs!

FUTURE WORK!

- Binding clients to nodes
- Adding partitions
- Exactly-once semantics
- Streams

Thanks!

Funded by
Redpanda Data

- Camilo Aguilar

- Travis Bischel

- Bob Dever

- Juan Castilla

- Alexander Gallego

- Dhruv Gupta

- Michal Maslanka

- Denis Rystsov

- John Spray

- Coral Waters

- David Wang

- Noah Watkins

- Allison Daly



<https://jepsen.io>

<https://redpanda.com>